

STRING-NET MODELS FOR PIVOTAL BICATEGORIES

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ABSTRACT. We develop a string-net construction of a modular functor whose algebraic input is a pivotal bicategory; this extends the standard construction based on a spherical fusion category. An essential ingredient in our construction is a graphical calculus for pivotal bicategories, which we express in terms of a category of colored corollas. The globalization of this calculus to oriented surfaces yields the bicategorical string-net spaces as colimits. We show that every rigid separable Frobenius functor between strictly pivotal bicategories induces linear maps between the corresponding bicategorical string-net spaces that are compatible with the mapping class group actions and with sewing. Our results are inspired by and have applications to the description of correlators in two-dimensional conformal field theories.

Contents

| | | |
|---|---|-----|
| 1 | Introduction | 474 |
| 2 | Graphical calculus for pivotal bicategories | 478 |
| 3 | String-net models based on pivotal bicategories | 511 |
| 4 | Application: Universal correlators in RCFT | 532 |

1. Introduction

For several decades, skein theoretic constructions have played a prominent role in quantum topology, see e.g. [BHMV, Wa] for early work. Such constructions are appealing for several reasons: They exist, for suitable algebraic input data, in various dimensions, and they provide a rather direct relation between these algebraic data and geometry, whereby they afford in particular geometric actions of mapping class groups. Moreover, skein theoretic

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constructions can typically be extended to higher codimensions, so that they can provide extended topological field theories. For instance, they allow one to construct interesting categories by evaluating the algebraic input data on manifolds of codimension 2, see e.g. [Ho] for categories associated to one-manifolds (cylinder categories) and e.g. [BeBJ, KiT] for categories associated to two-manifolds (elliptic Drinfeld centers etc.). Finally, there are close links to factorization homology [Coo]. More recently, in the form of string-net models, skein theoretic constructions have turned out to provide a particularly direct description of modular functors which form the basis for the construction of correlators of two-dimensional conformal field theories [SY, Tr, FuSY].

In the present paper we study string-net constructions that associate vector spaces to oriented surfaces and categories to oriented one-manifolds. Traditionally, following [LeW], in such constructions a spherical fusion category \mathcal{C} is taken as the algebraic input – one considers finite embedded graphs whose edges and vertices are labeled by objects and morphisms of \mathcal{C} , respectively. Here we take instead a more general algebraic datum as input: a pivotal bicategory \mathcal{B} . This is a three-layered structure and, as a consequence, we now have labels not only for edges and vertices of an embedded graph, but also for the connected components of the complement of the graph. Given the categorical and geometric dimensions involved, the idea to use a pivotal bicategory is indeed most natural, involved, and it dates back at least to [MoW, App.C.2]. What we achieve in the present paper is a precise construction based on pivotal bicategories in the sense of e.g. [FuGJS]. In the assumptions about our input bicategories we are parsimonious, the essential requirement being that they are linear over a field k ; in particular, for most of our considerations we do not need to impose semisimplicity.

The basic idea of string nets is to globalize a graphical calculus that exists for some standard canvas to general manifolds with the same tangential structure as the standard canvas. Accordingly, we start in Section 2 by developing a graphical calculus for pivotal bicategories, with the oriented disk as a standard canvas. First, in Section 2.1, we summarize the ordinary string diagram calculus for general bicategories, for which the canvas is the standard square in the complex plane (regarded as 2-framed, with vector fields parallel to the two coordinate axes) and the coupons for 2-morphisms are rectangles. If the bicategory \mathcal{B} is endowed with the additional structure of a (strictly) pivotal bicategory, a different graphical calculus can be derived from the ordinary string diagram calculus. This calculus has the standard oriented disk as its canvas and treats in- and outputs of a 2-morphism on the same footing. Thereby the coupons can be chosen to be circles; as we explain in Section 2.6, the space of colors associated with such a circular coupon can be defined as a limit over a contractible groupoid.

To be suited for string-net models, we need a concise formalization of this graphical calculus. This is provided in Sections 2.6 and 2.15 in terms of a symmetric monoidal functor $\mathbf{GCal}_{\mathcal{B}}$ from a category $\mathbf{Crl}_{\mathcal{B}}$ of partially \mathcal{B} -colored corollas to vector spaces, as described in Summary 2.18. In Section 2.21 we prove structure theorems that are instrumental later on: we show that (non-trivial) morphisms in the category $\mathbf{Crl}_{\mathcal{B}}$ can be decomposed into a finite disjoint union of partial compositions of morphisms of a few

specific types (Proposition 2.22) and establish a similar presentation of the subcategory $\mathbf{Crl}_\mathcal{B}^{\text{conn}}$ of $\mathbf{Crl}_\mathcal{B}$ that only contains those morphisms of $\mathbf{Crl}_\mathcal{B}$ all of whose underlying graphs are connected in the disks they are embedded in (Corollary 2.23).

These presentations of $\mathbf{Crl}_\mathcal{B}$ and $\mathbf{Crl}_\mathcal{B}^{\text{conn}}$ allow us to investigate the functoriality of the graphical calculus encoded in the functor $\mathbf{GCal}_\mathcal{B}$ in its dependence on the bicategory \mathcal{B} ; this is developed in Section 2.25. It is worth stressing that considering pseudofunctors between bicategories \mathcal{B} and \mathcal{B}' does not provide sufficient insight; instead, we examine Frobenius monoidal functors, adapted to bicategories. Such functors first arose in the context of linearly distributive categories. These are monoidal categories with two tensor products, which makes it natural to consider functors with monoidal constraints that are lax for one of the tensor products and oplax for the other. The lax and oplax structures are not required to be inverse to each other (which in the linearly distributive context would not make sense). Instead, they have to satisfy a compatibility condition which implies in particular that the image of the terminal bicategory with one object, one 1-morphism and one 2-morphism under a Frobenius functor gives a 1-endomorphism with the structure of a Frobenius algebra. In the present paper, we are interested in general bicategories, rather than in the special case of monoidal categories; the notion of a Frobenius functor can be easily extended to these. We thus introduce the notion of conjugation of a 2-morphism by a functor equipped with lax and oplax structures (Definition 2.26). We show that conjugation by a rigid Frobenius functor between two strictly pivotal bicategories \mathcal{B} and \mathcal{B}' canonically induces a monoidal natural transformation between the graphical calculi for \mathcal{B} and \mathcal{B}' when restricted to $\mathbf{Crl}_\mathcal{B}^{\text{conn}}$ and $\mathbf{Crl}_{\mathcal{B}'}^{\text{conn}}$, respectively (Theorem 2.30), and that this canonically extends to a monoidal natural transformation between the full graphical calculi if F is even a rigid pseudofunctor (Corollary 2.33).

The two types of bicategories we are particularly interested in are both based on a pivotal category \mathcal{C} : its delooping \mathcal{BC} , and the bicategory $\mathcal{Fr}(\mathcal{C})$ which has simple special symmetric Frobenius algebras internal to \mathcal{C} as objects, bimodules over a pair of such algebras as 1-morphisms, and bimodule morphisms as 2-morphisms. The obvious forgetful functor \mathcal{U} from $\mathcal{Fr}(\mathcal{C})$ to \mathcal{BC} is canonically a rigid separable Frobenius functor. In applications a crucial fact, explained in Example 2.34, is that the lax and oplax structures of \mathcal{U} provide an idempotent which exhibits the tensor product over a special symmetric Frobenius algebra as a retract.

In Section 3 we globalize the graphical calculus and construct modular functors through a string-net construction that uses a strictly pivotal bicategory \mathcal{B} as an input. In Section 3.1 we define the (bare) string-net space $\text{SN}_\mathcal{B}^\circ(\Sigma, \mathbf{b})$ assigned to a surface Σ and a \mathcal{B} -boundary datum \mathbf{b} on Σ . Similarly as in the conventional string-net construction, this is based on the idea to utilize the graphical calculus on disks to impose local relations on a vector space freely generated by all \mathcal{B} -colored graphs on Σ whose boundary datum is given by \mathbf{b} . As usual, the space $\text{SN}_\mathcal{B}^\circ(\Sigma, \mathbf{b})$ carries a geometric action of the mapping class group of Σ . In Section 3.7 we characterize $\text{SN}_\mathcal{B}^\circ(\Sigma, \mathbf{b})$ as a colimit over a functor $\mathcal{E}_\mathcal{B}^{\Sigma, \mathbf{b}}$ from a category of partially \mathcal{B} -colored graphs to vector spaces (Theorem 3.9). This insight lends itself to future generalizations of the string-net construction, e.g. replacing linear categories by

dg-categories and colimits by homotopy colimits.

In Section 3.10 we show that the string-net construction is functorial with respect to rigid pseudofunctors. Recall that, by Corollary 2.33, this class of functors preserves the entire graphical calculus. Concretely, we show that for any surface Σ and \mathcal{B} -boundary datum \mathbf{b} on Σ , a rigid pseudofunctor F between strictly pivotal bicategories \mathcal{B} and \mathcal{B}' gives rise to a canonical $\text{Map}(\Sigma)$ -intertwiner between the string-net spaces $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ and $\text{SN}_{\mathcal{B}'}^{\circ}(\Sigma, \mathbf{b}')$, where the \mathcal{B}' -boundary datum \mathbf{b}' is obtained from \mathbf{b} by a change of coloring induced by F (Theorem 3.11). Having now at our disposal string-net spaces on non-trivial manifolds, we can introduce, in Sections 3.13 and 3.17, the cylinder categories over circles and intervals, respectively. For the latter we impose the additional requirement that the strictly pivotal bicategory \mathcal{B} must be pointed, i.e. be endowed with a distinguished object. This allows us to find the cylinder categories for intervals as the endomorphism category of the distinguished object (Proposition 3.22), and to obtain the central result of Section 3.21, which states that the assignment of cylinder categories is functorial under embeddings of 1-manifolds (Proposition 3.24).

As a standard feature of skein theoretic constructions, the cylinder categories have, in general, no reason to be idempotent complete or even abelian; we examine their idempotent completion in Section 3.27. The general bicategorical string-net construction shares characteristic features of the special case based on a spherical fusion category. In particular, as we show in Section 3.29 (Theorem 3.31), it obeys factorization, or excision, which we formulate in terms of coends. Finally, in Section 3.34 we can combine all these results to show, in Theorem 3.36 that the string-net construction based on a pointed strictly pivotal bicategory \mathcal{B} provides an open-closed modular functor $\mathcal{SN}_{\mathcal{B}}^{\circ}$, i.e. a symmetric monoidal pseudofunctor

$$\mathcal{SN}_{\mathcal{B}}^{\circ} : \text{Bord}_{2,o/c}^{\text{or}} \longrightarrow \mathcal{P}\text{rof}_{\mathbf{k}}$$

from the symmetric monoidal bicategory of open-closed bordisms to the symmetric monoidal bicategory of \mathbf{k} -linear profunctors. The same applies to the functor $\mathcal{SN}_{\mathcal{B}}$ furnished by idempotent-completed string nets.

In the final Section 4 we provide an application of bicategorical string nets to correlators of two-dimensional conformal field theories. This application makes use of the fact that if the input bicategory \mathcal{B} is the delooping of a spherical fusion category \mathcal{C} , then the open-closed modular functor $\mathcal{SN}_{\mathcal{B}}$ extends the standard Turaev-Viro functor. In [FuSY] a construction of correlators via \mathcal{C} -colored string nets has been achieved. This construction can be paraphrased as follows: A topological world sheet \mathcal{S} with physical boundaries and topological defect lines for a rational CFT whose chiral data are encoded in a modular fusion category \mathcal{C} is naturally an $\mathcal{F}r(\mathcal{C})$ -colored graph $\tilde{\mathcal{S}}$ on its underlying surface Σ ; via the canonical quotient map $q(\Sigma, \mathbf{b}) : \mathbf{kG}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b}) \twoheadrightarrow \text{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}(\Sigma, \mathbf{b})$ it thus tautologically determines a vector $[\tilde{\mathcal{S}}]$ in the string-net space based on the bicategory $\mathcal{F}r(\mathcal{C})$. On the other hand, in [FuSY] we associated to any world sheet \mathcal{S} an explicit string net and a corresponding vector

$$\text{Cor}_{\mathcal{C}}(\mathcal{S}) \in \text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma}(\mathbf{b}_{\mathcal{S}}))$$

and showed that this assignment gives a consistent system of correlators [FuSY, Thm. 3.28]. The assignment of correlators to world sheets with fixed underlying surface Σ and boundary datum \mathbf{b} is encoded in a linear map $\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b}) : \mathbb{k}\mathcal{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b}) \rightarrow \text{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b}))$. The main result of Section 4 is the proof of Theorem 4.8, which states that this linear map factorizes over the string-net spaces based on $\mathcal{F}r(\mathcal{C})$: For every compact oriented surface Σ and $\mathcal{F}r(\mathcal{C})$ -boundary datum \mathbf{b} there is a unique $\text{Map}(\Sigma)$ -intertwiner

$$\text{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b}) : \text{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}(\Sigma, \mathbf{b}) \rightarrow \text{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b}))$$

such that

$$\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b}) = \text{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b}) \circ \mathbf{q}(\Sigma, \mathbf{b}).$$

This establishes a new conceptual role of defect data: they constitute the input for a string-net construction that provides a systematic home for relations between different correlators.

This may be seen as a special case of the following general statement (Theorem 4.10): Every rigid separable Frobenius functor between strictly pivotal bicategories induces linear maps between the corresponding bicategorical string-net spaces that are compatible with the mapping class group actions and with sewing, albeit in a less straightforward way than rigid pseudofunctors do – in addition to conjugating by the functor, full Frobenius graphs are added to compensate for the incomplete preservation of the graphical calculi.

2. Graphical calculus for pivotal bicategories

2.1. RECTANGULAR STRING DIAGRAMS FOR BICATEGORIES. A bicategory \mathcal{B} can be described as a category weakly enriched in the symmetric monoidal 2-category \mathcal{Cat} of small categories, functors and natural transformations, with the Cartesian product as monoidal product. Thus in particular for any pair of objects $a, b \in \mathcal{B}$ there is a *hom-category* $\mathcal{B}(a, b)$. The only general requirements that we impose on \mathcal{B} is that all these hom-categories are themselves enriched in the category $\text{Vect}_{\mathbb{k}}$ of (not necessarily finite-dimensional) \mathbb{k} -vector spaces and linear maps, and that the endomorphisms of the identity 1-morphism id_a of any object $a \in \mathcal{B}$ are isomorphic as a \mathbb{k} -algebra to the ground field \mathbb{k} . Here \mathbb{k} is an algebraically closed field, which we fix once and for all.

For horizontal composition we use *diagrammatic order*. We denote the horizontal compositions generally by “ \star ”, but for brevity in some long formulas suppress this symbol and indicate the composition instead just by juxtaposition. The vertical composition of 2-morphisms is denoted by “ \circ ”. Composite 2-morphisms can be expressed through *pasting diagrams* or, alternatively, through *string diagrams* on the standard square $I \times I$ as a canvas. The two descriptions are Poincaré dual to each other. For instance, given objects $a, b, c \in \mathcal{B}$, 1-morphisms $f, f', f'' \in \mathcal{B}(a, b)$ and $g, g' \in \mathcal{B}(b, c)$, and 2-morphisms $\alpha : f \Rightarrow f'$,

$\beta: f' \Rightarrow f''$ and $\gamma: g \Rightarrow g'$, the pasting diagram for $(\beta \circ \alpha) \star \gamma: f \star g \Rightarrow f'' \star g'$ is

$$\begin{array}{ccc}
 & f'' & \\
 a & \xrightarrow{f'} & b \\
 & f & \\
 & \alpha & \\
 & \beta & \\
 & f'' &
 \end{array}
 \quad
 \begin{array}{ccc}
 & g' & \\
 & \gamma & \\
 & g & \\
 c & &
 \end{array}
 \quad (2.1)$$

We portray the vertices of a string diagram as rectangular coupons and shade the two-dimensional regions with different colors that indicate the different objects. The string diagram corresponding to the pasting diagram (2.1) is then

$$\begin{array}{c}
 f'' \\
 \uparrow \beta \\
 f' \\
 \uparrow \alpha \\
 f
 \end{array}
 \quad
 \begin{array}{c}
 g' \\
 \uparrow \gamma \\
 g
 \end{array}
 \quad (2.2)$$

Or, to give a somewhat more complicated example, the diagrams

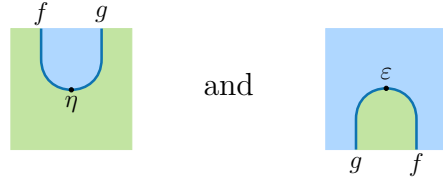
$$\begin{array}{c}
 f'' \\
 \uparrow \beta \\
 f' \\
 \uparrow \alpha \\
 f
 \end{array}
 \quad
 \begin{array}{c}
 b'' \\
 \uparrow g' \\
 b' \\
 \downarrow g \\
 b
 \end{array}
 \quad
 \begin{array}{c}
 h' \\
 \uparrow \gamma \\
 h
 \end{array}
 \quad
 c \xrightarrow{i} d
 \quad \text{and} \quad
 \begin{array}{c}
 f'' \\
 \uparrow \beta \\
 f' \\
 \uparrow \alpha \\
 f
 \end{array}
 \quad
 \begin{array}{c}
 g' \\
 \uparrow \gamma \\
 g
 \end{array}
 \quad
 \begin{array}{c}
 h' \\
 \uparrow \\
 h
 \end{array}
 \quad
 i
 \quad (2.3)$$

express the same 2-morphism in \mathcal{B} .

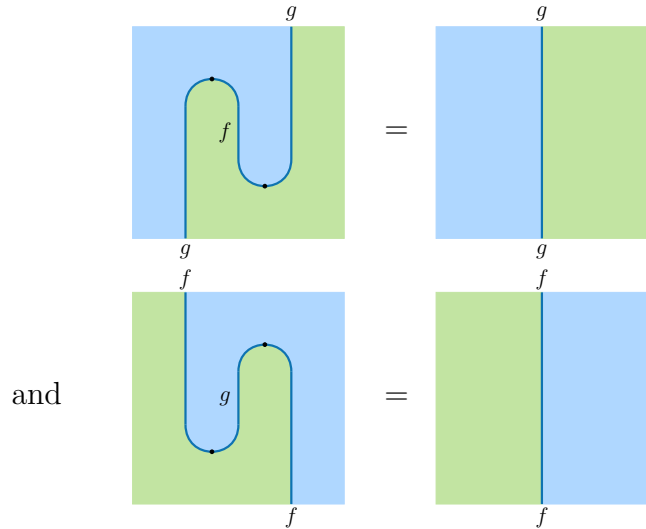
Actually, to make sense of either type of diagram, one first needs to select for each layer of horizontal composite of 1-morphisms a bracketing, which includes a choice of insertions of identity 1-morphisms. However, thanks to the *coherence theorem* for bicategories (see e.g. [JoY, Sect.3.6]), between each pair of bracketed horizontal composites of the same composable sequence of 1-morphisms there is a unique 2-isomorphism made up of associators and unitors that connects the pair. As a consequence, for any choice of bracketings for the *source* and the *target* of a (string or pasting) diagram, there is a unique 2-morphism assigned to the diagram, and the 2-morphisms for any two such choices are connected by a unique isomorphism of 2-hom spaces. Thus a pasting or string diagram uniquely determines a *contractible groupoid* in $\text{Vect}_{\mathbb{K}}$ whose vertices are the 2-hom spaces corresponding to the possible bracketings, as well as a coherent choice of elements in each of the 2-hom spaces. We refer to this coherent choice as the *value* of the diagram.

Put differently, every equality of string (or pasting) diagrams stands for an infinite family of equalities in different 2-hom spaces, one for each simultaneous choice of bracketings for both sides of the equality. Alternatively, one can invoke the *strictification theorem* for bicategories (see e.g. [Gur, Ch. 2]), according to which every bicategory is canonically biequivalent to a canonical strict 2-category associated with it, and treat any bicategory as if it were strict. We will freely use both of these perspectives.

The value of a string diagram is unaffected by any isotopy of the diagram that fixes the orientation of the rectangular coupons while keeping the diagram *progressive* and the end points of its legs fixed. To allow also for non-progressive string diagrams, one needs appropriate dualities. A *dual pair*, or *adjoint pair*, in a bicategory \mathcal{B} is a quadruple $(f, g, \eta, \varepsilon)$ consisting of two 1-morphisms $f \in \mathcal{B}(a, b)$ and $g \in \mathcal{B}(b, a)$ and two 2-morphisms $\eta: \text{id}_a \Rightarrow f \star g$ and $\varepsilon: g \star f \Rightarrow \text{id}_b$, called the unit and counit of the dual pair, that satisfy two *yanking equalities*. When the unit and counit are represented by the string diagrams


(2.4)

the yanking equalities read (after making the identity 1-morphisms invisible)


(2.5)

Duals are unique up to unique isomorphisms. We call f the *left dual* (or *left adjoint*) of g and write $f = {}^\vee g$, while g is called the *right dual* (or *right adjoint*) of f , written as $g = f^\vee$.

A *bicategory with duals* is a bicategory \mathcal{B} such that every 1-morphism in \mathcal{B} has both a left and a right dual. Fixing a right dual for each 1-morphism yields a pseudofunctor

$$(-)^\vee : \mathcal{B} \longrightarrow \mathcal{B}^{\text{co op}} \quad (2.6)$$

from \mathcal{B} to the bicategory $\mathcal{B}^{\text{co op}}$ that obtained by reversing both the 1- and the 2-morphisms in \mathcal{B} . This pseudofunctor is the identity on objects and sends every 1-morphism to its

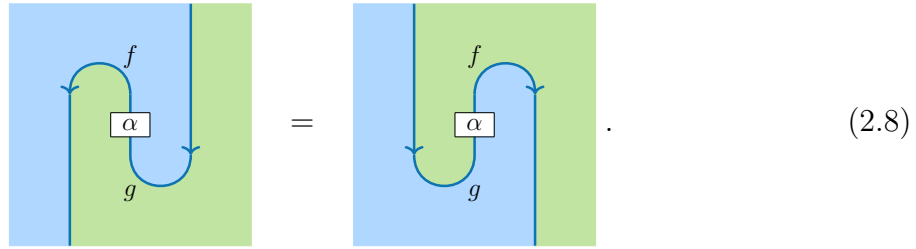
chosen right dual and every 2-morphism $\alpha: f \Rightarrow g$ to its *transpose* $\alpha^\vee: g^\vee \Rightarrow f^\vee$. A *pivotal* bicategory is one for which the double dual can be trivialized [FuGJS, Sect. 5.3]:

2.2. DEFINITION.

- (i) A *pivotal structure* on a bicategory \mathcal{B} with fixed left and right duals is an identity component pseudonatural transformation (i.e. every component 1-morphism is an identity) $\text{id}_{\mathcal{B}} \Rightarrow (-)^{\vee\vee}$. Equipped with a pivotal structure, \mathcal{B} is called a *pivotal bicategory*.
- (ii) A *strictly pivotal bicategory* is a pivotal bicategory for which the double dual is the identity,

$$\text{id}_{\mathcal{B}} = (-)^{\vee\vee}. \quad (2.7)$$

In a strictly pivotal bicategory \mathcal{B} we have ${}^\vee f = ({}^\vee f)^{\vee\vee} = f^\vee$, so that we can speak of *the dual* of a 1-morphism. In a string diagram we can then replace a string labeled by the dual of f by a string labeled by f but having opposite direction. As a consequence, for strictly pivotal bicategories, non-progressive diagrams make sense. Moreover, for any 2-morphism $\alpha: f \Rightarrow g$ in a strictly pivotal bicategory one has



$$= . \quad (2.8)$$

It follows that an isotopy resulting in a 2π -rotation of any coupon in a string diagram for a strictly pivotal bicategory does not affect the value of the diagram. Note that the canvas of the string diagrams – the standard square $I \times I$ – comes with a canonical 2-framing, i.e. a trivialization of its tangent bundle T given by two non-vanishing vector fields parallel to the x - and y -axis of $\mathbb{R}^2 = T(I \times I)$, respectively. In both diagrams in (2.8) the coupon is aligned with this canonical 2-framing, so that in particular its boundary segments are parallel to those of $I \times I$. We summarize these observations as

2.3. PROPOSITION. *Any (not necessarily progressive) string diagram on the standard square (with coupons not necessarily aligned with the frame) for a strictly pivotal bicategory \mathcal{B} has a well defined value. Any isotopy of the diagram that keeps the end points of its legs fixed (but may rotate the coupons) leaves this value unchanged.*

PROOF. Any string diagram Γ on $I \times I$ is isotopic, through an isotopy that fixes the end points of its legs, to a string diagram $\tilde{\Gamma}$ whose coupons are aligned with the frame. Moreover, any other choice $\tilde{\Gamma}'$ that yields such a diagram differs from $\tilde{\Gamma}$ by an isotopy that fixes the end points and rotates each coupon by some multiple of 2π , which because of the equality (2.8) does not change the value of the diagram. ■

There are two types – \mathcal{BC} and $\mathcal{Fr}(\mathcal{C})$ – of pivotal bicategories that are of particular interest to us:

2.4. **EXAMPLE.** Given a pivotal tensor category \mathcal{C} , its *delooping* \mathcal{BC} , i.e. \mathcal{C} viewed as a bicategory with a single object, is a pivotal bicategory, with the pivotal structure of \mathcal{C} viewed as a pivotal structure for \mathcal{BC} . Upon strictifying the pivotal structure of \mathcal{C} (which is always possible [NS, Thm. 2.2]), the bicategory \mathcal{BC} becomes strictly pivotal.

2.5. **EXAMPLE.** For any pivotal tensor category \mathcal{C} there is a bicategory $\mathcal{Fr}(\mathcal{C})$ which has simple special symmetric Frobenius algebras internal to \mathcal{C} as objects, bimodules over a pair of such algebras as 1-morphisms, and bimodule morphisms as 2-morphisms. The bicategory inherits from \mathcal{C} a canonical pivotal structure. If \mathcal{C} is strictly pivotal, then so is $\mathcal{Fr}(\mathcal{C})$. That we require the symmetric Frobenius algebras which are objects of this bicategory to be special is motivated by the application to conformal field theory that we will describe in Section 4. An essential ingredient in that application is the idempotent that realizes (see Example 2.34) the tensor product over the algebra.

2.6. **PARTIALLY COLORED COROLLAS AND POLARIZATIONS.** The string diagrams for bicategories described in Section 2.1 use the standard square $I \times I$ as the canvas. Asserting that a coupon is *aligned with the frame* is thus the same as saying that the coupon is aligned with the canonical 2-framing of the standard square. For a pivotal category, by Proposition 2.3, the 2-framing in the *interior* of the square does not affect the evaluation of any string diagram for a (strictly) pivotal bicategory: the value is unchanged under isotopies which do not necessarily preserve the alignment of the coupons with the 2-framing. In contrast, so far we still use the 2-framing at the *boundary* of the square: it tells us which part of the boundary is the bottom and which part is the top, namely the intervals in $\partial(I \times I)$ at which the framing is pointing inwards and outwards, respectively. The distinction between top and the bottom separates the *output ports* from the *input ports* of a string diagram. As we will show in this section, in the case of a *pivotal* bicategory, the distinction between input and output is immaterial to the graphical calculus, so that the 2-framing of the canvas can be completely disregarded. This result constitutes a crucial step towards the formulation of string-net models – in a sense, string-net models are generalizations of the graphical calculus for which string diagrams can have *any* compact oriented surface as their canvas.

We now fix a strictly pivotal bicategory \mathcal{B} . To proceed, let us introduce the notion of a *partially \mathcal{B} -colored graph* on an oriented surface. For the moment we restrict our attention to the case that the surface is the *standard disk*. By definition, the standard disk $D \subset \mathbb{C}$ is the closed unit disk of radius 1 centered at $0 \in \mathbb{C}$; the boundary $S^1 := \partial D = \{z \in \mathbb{C} \mid |z| = 1\}$ of the standard disk is called the *standard circle*. More general surfaces than D will be considered later on.

2.7. **DEFINITION.** A partially \mathcal{B} -colored graph Γ on a compact oriented surface Σ with possibly non-empty boundary consists of the following data:

- (i) An underlying directed finite graph, i.e. a diagram

$$\begin{array}{ccccc}
 & & \overset{i}{\curvearrowright} & & \\
 E(\Gamma) & \xleftarrow{\quad} & H(\Gamma) & \xleftarrow{s} & I(\Gamma) \xrightarrow{t} V(\Gamma) \\
 & & \underset{\delta}{\curvearrowleft} & &
 \end{array} \tag{2.9}$$

of finite sets, with $V(\Gamma)$, $H(\Gamma)$ and $I(\Gamma)$ the sets of internal vertices, of half-edges, and of half-edges that touch an internal vertex, respectively. The map t indicates the incidence of half-edges in $I(\Gamma)$ to the internal vertices, while s is the canonical inclusion. The map i is the fixed-point-free involution that indicates the juncture of pairs of half-edges; its set of orbits is $E(\Gamma)$, to be interpreted as the set of edges. An edge which consists of a pair of internal half-edges is called an internal edge, while an edge consisting of a single half-edge is called a leg. The map δ is a section of the canonical quotient map $H(\Gamma) \twoheadrightarrow E(\Gamma)$; for each edge it picks out its starting half-edge and thereby directs the edges.

- (ii) An embedding into Σ of the geometric realization $|\Gamma|$ of the underlying graph, i.e. of the topological space

$$\left(\bigsqcup_{v \in V(\Gamma)} \{v\} \right) \sqcup \left(\bigsqcup_{e \in H(\Gamma)} [0, 1]_e \right) / \sim, \tag{2.10}$$

where the equivalence relation is generated by $[0, 1]_e \ni 1 \sim 1 \in [0, 1]_{i(e)}$ for all $e \in H(\Gamma)$ and $[0, 1]_e \ni 0 \sim \{t(e)\}$ for all $e \in I(\Gamma)$. The embedding is subject to the requirement that the intersection of $\partial\Sigma$ with the image of $|\Gamma|$ is exactly the image of the end points of the legs, i.e. the image of $\{0 \in [0, 1]_l \hookrightarrow |\Gamma|\}_{l \in H(\Gamma) \setminus I(\Gamma)}$.

- (iii) A coloring of the patches, i.e. of the connected components of $\Sigma \setminus |\Gamma|$, with objects of \mathcal{B} , as well as a coloring of the directed edges with 1-morphisms of \mathcal{B} . More specifically, a directed edge e is labeled with a 1-morphism $f \in B(a_l, a_r)$, where a_l and a_r are the labels of the patches adjacent to the left and to the right of e , respectively.

Note that the vertices of the graph Γ are *not* colored – hence the name *partially* \mathcal{B} -colored graph – and that the empty graph $\Gamma = \emptyset$ is allowed. Also, the distinction between left and right in part (iii) of the definition is meant to be such that a vector pointing to the right and a vector pointing in the direction of e form a coordinate system whose orientation coincides with the orientation of Σ .

2.8. DEFINITION. A partially \mathcal{B} -colored corolla is a partially \mathcal{B} -colored graph on the standard disk D whose underlying directed finite graph has a single vertex – called the center of the corolla – which is mapped by the embedding to $0 \in D \subset \mathbb{C}$, and for which the image of each edge is a straight line connecting the center with a point on ∂D . (The number of edges is allowed to be zero.)

2.9. EXAMPLE. The following picture shows a partially \mathcal{B} -colored corolla K on the standard disk (which we equip with the counterclockwise orientation):

$$K = \begin{array}{c} \begin{array}{c} f \\ \uparrow \\ \text{green} \quad \text{blue} \\ \downarrow \\ \text{purple} \\ h \quad g \end{array} \\ \text{v} \\ c \end{array} \quad (2.11)$$

Here a, b, c are objects in \mathcal{B} and $f \in \mathcal{B}(a, b)$, $g \in \mathcal{B}(b, c)$ and $h \in \mathcal{B}(c, a)$ are 1-morphisms. We also indicate the objects by shading the regions with different colors,¹ In the pictures below, we will only keep the shadings and suppress the corresponding object labels.

2.10. DEFINITION. For \mathcal{B} a strictly pivotal bicategory, a \mathcal{B} -boundary datum \mathbf{b} on a compact oriented 1-manifold ℓ with possibly non-empty boundary consists of the following data:

- (i) A (possibly empty) finite set O_ℓ of points in the interior of ℓ .
- (ii) A coloring of each connected component $a \in \pi_0(\ell \setminus O_\ell)$ with an object $\mathbf{C}_\mathbf{b}(a) \in \mathcal{B}$.
- (iii) A coloring of each element of O_ℓ with a 1-morphism in \mathcal{B} .

Here the convention is that the color of $p \in O_\ell$ is a 1-morphism in $\mathcal{B}(\mathbf{C}_\mathbf{b}(a), \mathbf{C}_\mathbf{b}(a'))$ if the connected components a and a' of $\ell \setminus O_\ell$ are located to the left and right of p , respectively, with the orientation of ℓ pointing from the right to the left.

In view of the lack of coloring for the center of a corolla, there is a canonical bijection between the set of partially \mathcal{B} -colored corollas and the set of \mathcal{B} -boundary data on S^1 , as illustrated in the following picture:

$$\begin{array}{c} \begin{array}{c} f \\ \uparrow \\ \text{green} \quad \text{blue} \\ \downarrow \\ \text{purple} \\ h \quad g \end{array} \\ \text{v} \\ c \end{array} \longleftrightarrow \begin{array}{c} f \\ \bullet \\ \text{green} \quad \text{blue} \\ \bullet \\ \text{purple} \\ h^\vee \end{array} \quad (2.12)$$

We are now going to associate to the center v of a partially \mathcal{B} -colored corolla K a vector space $H_v^\mathcal{B}$, to be called the *space of colors* for the vertex v . We specify this space as a certain 2-hom space in \mathcal{B} . To define this space, we make use of the auxiliary datum of a linear order on the set $H(v)$ of half-edges incident to the vertex v . The orientation of D naturally induces a cyclic order on $H(v)$; with the chosen counterclockwise orientation of D , this cyclic order is clockwise. A linear order on $H(v)$ that is compatible with the induced cyclic order is uniquely determined by the choice of a *root*, or *starting half-edge*,

¹ In the color version: green, blue, and purple.

$e \in H(v)$. The choice of e determines a particular 2-hom space $h_v^{\mathcal{B}}(e)$ in \mathcal{B} as a space of invariants. For instance, if for the corolla K shown in Example 2.9 we choose the root e to be the half-edge e_h that is labeled by $h \in \mathcal{B}(c, a)$, then we associate to v the 2-hom space

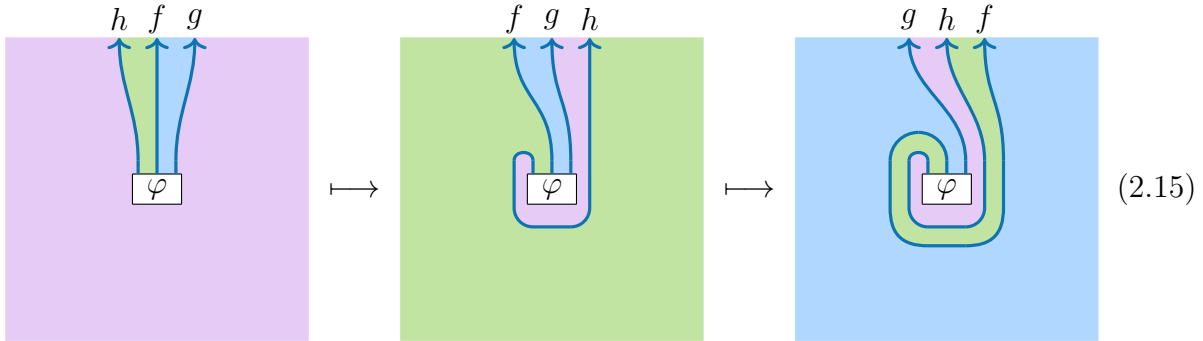
$$h_v^{\mathcal{B}}(e_h) := \text{End}_{\mathcal{B}}(c)(\text{id}_c, h \star f \star g). \quad (2.13)$$

Similarly, for the choices $e = e_f$ and $e = e_g$ we get the spaces $h_v^{\mathcal{B}}(e_f) := \text{End}_{\mathcal{B}}(a)(\text{id}_a, f \star g \star h)$ and $h_v^{\mathcal{B}}(e_g) := \text{End}_{\mathcal{B}}(b)(\text{id}_b, g \star h \star f)$, respectively.

By using the units and counits of dual pairs to turn edges around, we can construct canonically an isomorphism between any two of the so assigned 2-hom spaces. For instance, in the case of the corolla K from Example 2.9, changing the root from h to f and then to g gives rise to the chain

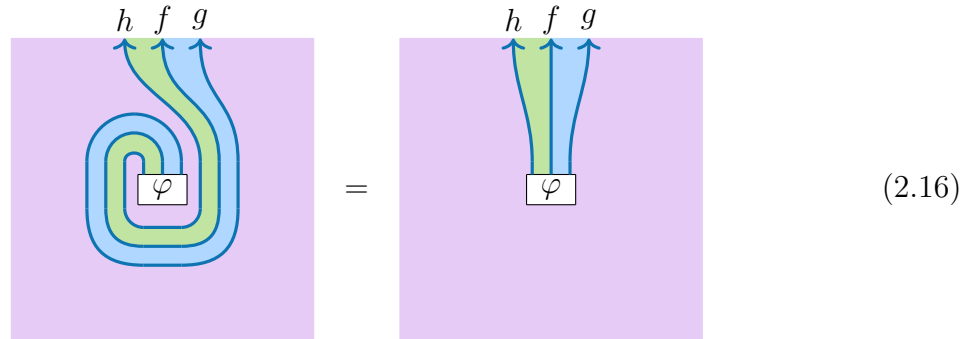
$$h_v^{\mathcal{B}}(e_h) \xrightarrow{\cong} h_v^{\mathcal{B}}(e_f) \xrightarrow{\cong} h_v^{\mathcal{B}}(e_g) \quad (2.14)$$

of isomorphisms of vector spaces, whose action on an element $\varphi \in h_v^{\mathcal{B}}(e_h) = \text{End}_{\mathcal{B}}(c)(\text{id}_c, hfg)$ is given by



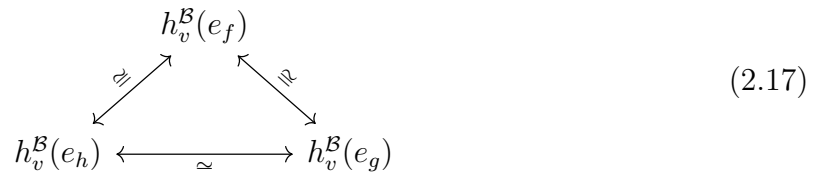
$$(2.15)$$

Further, owing to the yanking equations for duals together with (2.8), after one further step we get back the original element φ :



$$(2.16)$$

Altogether this yields a groupoid in $\text{Vect}_{\mathbf{k}}$ that is generated by the diagram



$$(2.17)$$

This groupoid is contractible, i.e. between any ordered pair of objects there is a single morphism, which in particular implies that any two composites of morphisms with the same source and same target are equal. Thus by selecting an element in *any* of the three 2-hom spaces, we simultaneously select an element in *each* of the spaces.

This leads us to the following prescription: Let K be a partially \mathcal{B} -colored corolla with center v and $n = |H(v)| > 0$. Then to K we associate:

1. A contractible groupoid $\mathcal{G}_v^{\mathcal{B}}$. The objects of $\mathcal{G}_v^{\mathcal{B}}$ are the elements of the set $H(v) = \{e_1, e_2, \dots, e_n\}$ of half-edges incident to v , which is conveniently indexed according to an arbitrary choice of compatible linear order. The morphisms of $\mathcal{G}_v^{\mathcal{B}}$ are generated by the diagram

$$\begin{array}{ccccc} e_1 & \longleftrightarrow & e_2 & \longleftrightarrow & e_3 \\ \updownarrow & & & & \updownarrow \\ e_n & \longleftrightarrow & \cdots & \longleftrightarrow & e_4 \end{array} \quad (2.18)$$

with relations uniquely determined by requiring that $\mathcal{G}_v^{\mathcal{B}}$ is contractible, i.e. that each hom-set $\mathcal{G}_v^{\mathcal{B}}(e_i, e_j)$ has a single element.

2. A functor $h_v^{\mathcal{B}}: \mathcal{G}_v^{\mathcal{B}} \rightarrow \text{Vect}_{\mathbb{k}}$ that acts on objects by

$$e_i \mapsto h_v^{\mathcal{B}}(e_i) := \text{End}_{\mathcal{B}}(a_i)(\text{id}_{a_i}, f_i^{\epsilon_i} \star f_{i+1}^{\epsilon_{i+1}} \star \cdots \star f_{i+n-1}^{\epsilon_{i+n-1}}), \quad (2.19)$$

where the subscripts are taken mod n , f_j is the color of the half-edge e_j with $f_j^{\epsilon_j} \in \mathcal{B}(a_j, a_{j+1})$, and where $\epsilon_j = +$ if e_j is directed away from v and $\epsilon_j = -$ otherwise. The action on the generating morphisms is given by “dragging the leg around”.

We can now assign a space of colors to a partially colored corolla:

2.11. DEFINITION. *The vector space of colors $H_v^{\mathcal{B}}$ for a partially colored corolla with at least one leg is the limit*

$$H_v^{\mathcal{B}} := \lim h_v^{\mathcal{B}} \in \text{Vect}_{\mathbb{k}}. \quad (2.20)$$

For a corolla K_a without legs and its single patch $D \setminus v_a$ colored with an object $a \in \mathcal{B}$, the space of colors is

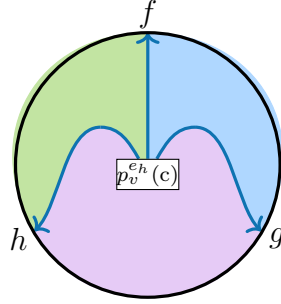
$$H_{v_a}^{\mathcal{B}} := \text{End}_{\mathcal{B}}(a)(\text{id}_a, \text{id}_a). \quad (2.21)$$

Being the limit of a contractible groupoid, the space $H_v^{\mathcal{B}}$ is determined by an isomorphism $p_v^e: H_v^{\mathcal{B}} \xrightarrow{\cong} h_v^{\mathcal{B}}(e)$ for any choice of $e \in H(v)$, which uniquely extends to a limit cone over $h_v^{\mathcal{B}}$ with every leg being an isomorphism of vector spaces. Therefore, by choosing a *color* for the vertex v , i.e. an element $c \in H_v^{\mathcal{B}}$, we produce a family

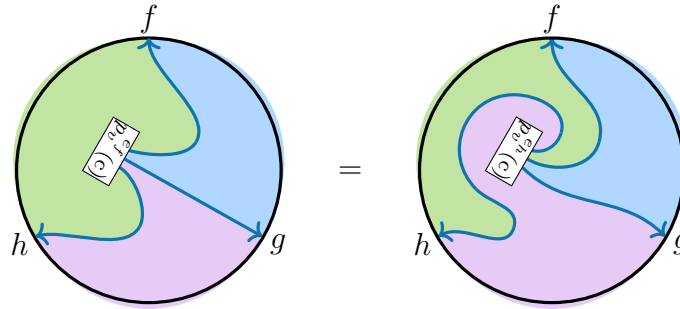
$$\{p_v^e(c) \in h_v^{\mathcal{B}}(e)\}_{e \in H(v)} \quad (2.22)$$

of 2-morphisms in \mathcal{B} that is coherent in the sense that for any two members $p_v^{e_i}(c)$ and $p_v^{e_j}(c)$ of the family there is a unique isomorphism $h_v^{\mathcal{B}}(e_i) \xrightarrow{\cong} h_v^{\mathcal{B}}(e_j)$. This isomorphism is obtained by evaluating the functor $h_v^{\mathcal{B}}$ on the unique morphism $e_i \xrightarrow{\cong} e_j$ in the groupoid $\mathcal{G}_v^{\mathcal{B}}$, such that $p_v^{e_i}(c)$ is mapped to $p_v^{e_j}(c)$.

2.12. REMARK. In still more detail, the relation between the elements in the different spaces of 2-morphisms that are determined by the same color $c \in H_v^{\mathcal{B}}$ can be expressed as follows: By selecting a root $e \in H(v)$ we produce, up to isotopies that fix the boundary $\partial D = S^1$, a string diagram with a rectangular coupon on D , which is isotopic relative to ∂D to the string diagram produced by choosing any other root. Let us illustrate this with the corolla (2.11). The choice $e = e_h$ results in the string diagram


(2.23)

If we choose instead $e = e_f$, then the diagram is


(2.24)

which is isotopic to (2.23).

So far we only considered 2-morphisms that have all non-trivial 1-morphisms in their outputs. This requirement is too restrictive for applications; to be able to remove it, we introduce the following notion:

2.13. DEFINITION. A polarization on the vertex v of a partially \mathcal{B} -colored corolla is a partition

$$H(v) = H^{\text{in}}(v) \sqcup H^{\text{out}}(v) \quad (2.25)$$

of the cyclically ordered set of half-edges into two linearly ordered sets $H^{\text{in}}(v)$ and $H^{\text{out}}(v)$ of input and output half-edges, in such a way that any two half-edges of the same type (i.e., either in- or output) are consecutive with respect to the cyclic order on $H(v)$ if they are consecutive with respect to the linear order on $H^{\text{in}}(v)$ or $H^{\text{out}}(v)$ that is induced by the cyclic order on $H(v)$.

If $H^{\text{in}}(v) = \emptyset$ or $H^{\text{out}}(v) = \emptyset$, a polarization reduces to a compatible linear order on $H(v)$. Note that whether a half-edge of a corolla is an input or output half-edge with

respect to a chosen polarization is not correlated with the direction of the edge to which it belongs.

With this notion we obtain a contractible groupoid $\widehat{\mathcal{G}}_v^{\mathcal{B}}$ that is generated by the set of all polarizations on v . We have a canonical extension

$$\mathcal{G}_v^{\mathcal{B}} \hookrightarrow \widehat{\mathcal{G}}_v^{\mathcal{B}} \xrightarrow{\simeq} 1 \quad (2.26)$$

of groupoids, as well as an extension

$$\begin{array}{ccc} \widehat{\mathcal{G}}_v^{\mathcal{B}} & \xrightarrow{\widehat{h}_v^{\mathcal{B}}} & \mathbf{Vect}_{\mathbb{k}} \\ \wr \downarrow & \nearrow h_v^{\mathcal{B}} & \\ \mathcal{G}_v^{\mathcal{B}} & & \end{array} \quad (2.27)$$

of the functor $h_v^{\mathcal{B}}$, which is defined by setting

$$\widehat{h}_v^{\mathcal{B}}(k) := \mathrm{Hom}_{\mathcal{B}(a_i, a_{j+1})}(f_{i+n-1}^{-\epsilon_{i+n-1}} \star \cdots \star f_{j+2}^{-\epsilon_{j+2}} \star f_{j+1}^{-\epsilon_{j+1}}, f_i^{\epsilon_i} \star f_{i+1}^{\epsilon_{i+1}} \star \cdots \star f_j^{\epsilon_j}). \quad (2.28)$$

for a polarization k with $H^{\mathrm{out}}(v) = \{e_i, e_{i+1}, \dots, c, e_j\}$ and $H^{\mathrm{in}}(v) = \{e_{j+1}, e_{j+2}, \dots, c, e_{i+n-1}\}$ (with indices counted mod n).

In the graphical description of a polarization we draw ingoing edges as dashed lines. (In the color version, we use in addition two different colors for the 1-morphisms that label the in- and outgoing edges.)

2.14. EXAMPLE. For the corolla (2.11) described in Example 2.9, we denote by



$$(2.29)$$

the polarization k that has $H^{\mathrm{in}}(v) = \{e_h\}$ and $H^{\mathrm{out}}(v) = \{e_f, e_g\}$. For this polarization we have $\widehat{h}_v^{\mathcal{B}}(k) = \mathrm{Hom}_{\mathcal{B}(a, c)}(h^{\vee}, f \star g)$.

It follows in particular that the space $H_v^{\mathcal{B}}$ is also equipped with a unique limit cone over $\widehat{h}_v^{\mathcal{B}}: \widehat{\mathcal{G}}_v^{\mathcal{B}} \rightarrow \mathbf{Vect}_{\mathbb{k}}$ that restricts to the limit cone over $h_v^{\mathcal{B}}$. We denote the legs of this cone (which are isomorphisms) by

$$\{p_v^k: H_v^{\mathcal{B}} \xrightarrow{\cong} \widehat{h}_v^{\mathcal{B}}(k)\}_{k \in \mathrm{obj} \widehat{\mathcal{G}}_v^{\mathcal{B}}}. \quad (2.30)$$

Now from a color $c \in H_v^{\mathcal{B}}$ we obtain an entire family of 2-morphisms $\{p_v^k(c) \in \widehat{h}_v^{\mathcal{B}}(k)\}_{k \in \mathrm{obj} \widehat{\mathcal{G}}_v^{\mathcal{B}}}$ that are related by dualities. For instance, with the polarization k selected in (2.29), the

2-morphism $p_v^k(c) \in \text{Hom}_{\mathcal{B}(a,c)}(h^\vee, f \star g)$ can be expressed in terms of $p_v^{e_h}(c) \in \text{End}_{\mathcal{B}}(c)(\text{id}_c, h \star f \star g)$ as

$$p_v^k(c) = \text{Diagram} \quad (2.31)$$

Together with a choice of polarization, a color $c \in H_v^{\mathcal{B}}$ again determines an isotopy class of string diagrams on D with a rectangular coupon, and the isotopy class is independent of the choice of polarization.

2.15. GRAPHICAL CALCULUS ON DISKS FOR PIVOTAL BICATEGORIES. To represent the isotopy class of the string diagrams which a color $c \in H_v^{\mathcal{B}}$ on the vertex v of a corolla produces, without choosing a specific polarization, we place a *circular coupon* labeled by c at the vertex v . For instance,

$$\text{Diagram} \quad (2.32)$$

represents the isotopy class of string diagrams that contains both the diagram (2.23) and the diagram (2.24). Conversely, a diagram like (2.32) can be represented by a string diagram on D with rectangular coupon that corresponds to any choice of polarization and a 2-morphism in the associated 2-hom vector space.

We allot to a partially \mathcal{B} -colored corolla K the space of colors for its center v . We denote this assignment by

$$\text{GCal}_{\mathcal{B}}(K) := H_v^{\mathcal{B}}. \quad (2.33)$$

(The notation $\text{GCal}_{\mathcal{B}}$ is supposed to remind of the ‘graphical calculus’.) The assignment (2.33) is functorial with respect to orientation preserving embeddings of the standard disk to itself that induce a local isomorphism of the partially colored embedded corollas, in such a way that the colors of the patches match while a half-edge may be mapped to either a half-edge with the same orientation and the same color or a half-edge with the

opposite orientation and the dual color. For instance, the picture

$$K = \text{disk with patches } f, g, h \hookrightarrow h^v \text{ disk with patches } f, g = K' \quad (2.34)$$

shows an embedding of the corolla K into K' , with the image of K indicated by the shaded region in K' , for which the half-edge colored by the 1-morphism h is mapped to the one colored by the dual h^\vee . The action on the assignment (2.33) is in this case given by

$$\begin{aligned} \mathbf{GCal}_{\mathcal{B}}(K) &= H_v^{\mathcal{B}} \xrightarrow{\cong} h_v^{\mathcal{B}}(e_h) = \text{End}_{\mathcal{B}}(c)(\text{id}_c, hfg) \\ &= h_{v'}^{\mathcal{B}}(e_{h^\vee}) \\ &\xrightarrow{\cong} H_{v'}^{\mathcal{B}} = \mathbf{GCal}_{\mathcal{B}}(K'). \end{aligned} \quad (2.35)$$

Note that due to this functoriality, for a *monochromatic* corolla K_{mnc} , i.e. a corolla each of whose half-edges is colored with the same 1-morphism and oriented in the same way, the vector space $\mathbf{GCal}_{\mathcal{B}}(K_{\text{mnc}})$ carries an action of the cyclic group of appropriate order.

Next we extend the assignment (2.33) to a symmetric monoidal functor

$$\mathbf{GCal}_{\mathcal{B}} : \mathbf{Crl}_{\mathcal{B}} \longrightarrow \mathbf{Vect}_{\mathbb{k}} \quad (2.36)$$

from a category of partially \mathcal{B} -colored corollas and graphs to $\mathbf{Vect}_{\mathbb{k}}$. We first introduce the source category of $\mathbf{GCal}_{\mathcal{B}}$. For \mathcal{B} a pivotal bicategory, a *boundary datum* of a partially \mathcal{B} -colored corolla is the \mathcal{B} -boundary datum, in the sense of Definition 2.10, for the boundary circle of the underlying disk that is determined by the colors of the patches of the disk.

2.16. DEFINITION. *Let \mathcal{B} be a pivotal bicategory. $\mathbf{Crl}_{\mathcal{B}}$ is the following symmetric monoidal category:*

- *Objects in $\mathbf{Crl}_{\mathcal{B}}$ are finite disjoint unions of partially \mathcal{B} -colored corollas. (This includes the empty disjoint union \emptyset .)*
- *A morphism of type $K_1 \sqcup \cdots \sqcup K_n \rightarrow K_{n+1}$ in $\mathbf{Crl}_{\mathcal{B}}$ is a partially \mathcal{B} -colored graph G on the standard disk whose boundary datum coincides with that of K_{n+1} , together with an orientation preserving embedding $D_1 \sqcup \cdots \sqcup D_n \hookrightarrow D_{n+1}$ of the underlying disks of the source (the number of which is allowed to be zero) to the underlying disk of the target. This embedding is required to induce a local isomorphism of the graphs that respects the colorings of the patches, while the orientation or_e and color f_e of a half-edge e are required to be respected either on the nose or up to simultaneous reversal of orientation and dualizing of the color. Moreover, each internal vertex of G must be covered by the image of exactly one of the disks D_1, \dots, D_n .*

- General morphisms in $\text{Crll}_{\mathcal{B}}$ are obtained by taking disjoint unions of morphisms of the type just described. The composition $G_2 \circ G_1$ of morphisms is given by blowing up the internal vertices of G_2 by G_1 , using the embeddings of disks.
- The symmetric monoidal product on $\text{Crll}_{\mathcal{B}}$ is given by disjoint union, with monoidal unit \emptyset .

2.17. EXAMPLE. An example for a morphism in $\text{Crll}_{\mathcal{B}}$ of type

$$G : \begin{array}{c} \text{Disk with } v_1 \text{ and } f_1, f_2, f_6 \\ \sqcup \\ \text{Disk with } v_2 \text{ and } f_3, f_4, f_5, f_6 \end{array} \longrightarrow \begin{array}{c} \text{Disk with } v_3 \text{ and } f_1, f_2, f_3, f_4 \end{array} \quad (2.37)$$

is determined by the following disk in which, as we already did in the picture (2.34) above, the images of the source disks in the target are indicated by shaded regions:

$$G = \begin{array}{c} \text{Disk with } v_3 \text{ and } f_1, f_2, f_3, f_4, f_5, f_6 \end{array} \quad (2.38)$$

An illustration of the blowing-up procedure that defines the composition of morphisms in $\text{Crll}_{\mathcal{B}}$ is given in the following picture:

$$\begin{array}{c} \text{Disk with } v_1 \text{ and } f_1, f_2, f_4 \end{array} \circ \begin{array}{c} \text{Disk with } v_2 \text{ and } f_1, f_2, f_3, f_4 \end{array} = \begin{array}{c} \text{Disk with } v_3 \text{ and } f_1, f_2, f_3, f_4 \end{array} \quad (2.39)$$

In words, the shaded region within the first disk D_1 specifies the embedding of the second disk D_2 into D_1 , and in particular the boundary datum for the shaded region in D_1 matches the boundary datum of the disk D_2 .

We are now in a position to define the functor (2.36). On objects, $\mathbf{GCal}_{\mathcal{B}}$ acts as

$$\mathbf{GCal}_{\mathcal{B}}(K_1 \sqcup \cdots \sqcup K_n) := \mathbf{GCal}_{\mathcal{B}}(K_1) \otimes_{\mathbb{k}} \cdots \otimes_{\mathbb{k}} \mathbf{GCal}_{\mathcal{B}}(K_n) = H_{v_1}^{\mathcal{B}} \otimes_{\mathbb{k}} \cdots \otimes_{\mathbb{k}} H_{v_n}^{\mathcal{B}}. \quad (2.40)$$

As the case $n = 0$, this includes $\mathbf{GCal}_{\mathcal{B}}(\emptyset) := \mathbb{k}$. To define the action of $\mathbf{GCal}_{\mathcal{B}}$ on morphisms, a more extensive description is needed. We demonstrate it for the case of a morphism $G: K_1 \sqcup K_2 \rightarrow K_3$ of the type shown in (2.38): To an arbitrarily chosen element $c_1 \otimes_{\mathbb{k}} c_2 \in \mathbf{GCal}_{\mathcal{B}}(K_1) \otimes_{\mathbb{k}} \mathbf{GCal}_{\mathcal{B}}(K_2) = H_{v_1}^{\mathcal{B}} \otimes_{\mathbb{k}} H_{v_2}^{\mathcal{B}}$ we assign an element $c_3 \in \mathbf{GCal}_{\mathcal{B}}(K_3) = H_{v_3}^{\mathcal{B}}$ and then obtain a linear map by linear extension. To specify c_3 , we first choose polarizations for the vertices $v_1 \in K_1$ and $v_2 \in K_2$, say

$$k_1 : \quad \text{and} \quad k_2 : \quad (2.41)$$

with k_1 defined by selecting the edge labeled with f_6 (drawn as a dashed line, and with label displayed in red in the color version) as its single input edge and k_2 defined by selecting the edge labeled with f_4 as its single input edge. Now for $i = 1, 2$ set $\tilde{c}_i := p_{v_i}^{k_i}(c_i) \in \widehat{h_{v_i}^{\mathcal{B}}}(k_i)$, with p_v^k given by (2.30). According to the discussion in Section 2.6 we obtain, up to isotopies relative to the boundary of D , two string diagrams with rectangular coupons on D that correspond to the choice of colors and polarizations. These string diagrams are then pushed forward along the embedding $D_1 \sqcup D_2 \hookrightarrow D_3$, replacing the images of the corollas K_1 and K_2 . Thereby we obtain, up to isotopies fixing the boundary, a string diagram G' with rectangular coupons on D that has the same boundary datum as K_3 :

$$G' = \quad (2.42)$$

We now choose a polarization for the vertex v_3 of the corolla that is the target of the

morphism, e.g.

$$k_3 : \quad \begin{array}{c} \text{Diagram of a circle with a central vertex } v_3. \text{ The circle is divided into three regions: blue (top), green (bottom-left), and purple (bottom-right).} \\ \text{Edges are labeled } f_1, f_2, f_3, f_4. \text{ } f_1 \text{ and } f_2 \text{ are blue, } f_3 \text{ is purple, and } f_4 \text{ is red.} \end{array} \quad (2.43)$$

with the edge labeled with f_4 as single input edge. With this choice we produce a string diagram G'' on the standard square $I \times I$, uniquely up to isotopies fixing the top and the bottom of the square setwise:

$$G'' = \quad \begin{array}{c} \text{Diagram of a square with two vertices } \tilde{c}_1 \text{ and } \tilde{c}_2. \text{ The square is divided into three regions: blue (top), green (bottom-left), and purple (bottom-right).} \\ \text{Edges are labeled } f_1, f_2, f_3, f_4, f_5, f_6. \text{ } f_1, f_2, f_3, f_5 \text{ are blue, } f_4 \text{ is purple, and } f_6 \text{ is green.} \end{array} \quad (2.44)$$

(Recall that, by construction, no 1-morphisms are ending on the left and right boundary segments.) As explained in Section 2.1, such a rectified string diagram uniquely determines a 2-morphism $\tilde{c}_3 \in \widehat{h}_{v_3}^{\mathcal{B}}(k_3)$. This 2-morphism \tilde{c}_3 determines, in turn, a unique element $c_3 = (p_{v_3}^{k_3})^{-1}(\tilde{c}_3) \in H_{v_3}^{\mathcal{B}} = \mathbf{GCal}_{\mathcal{B}}(K_3)$ which, besides on the colors c_1 and c_2 , only depends on the isotopy class (relative to the boundary) of G ; in particular, it does not depend on the auxiliary choice of polarizations. We write

$$c_3 = \left\langle \begin{array}{c} \text{Diagram of a circle with two vertices } c_1 \text{ and } c_2. \text{ The circle is divided into three regions: blue (top), green (bottom-left), and purple (bottom-right).} \\ \text{Edges are labeled } f_1, f_2, f_3, f_4, f_5, f_6. \text{ } f_1, f_2, f_3, f_5 \text{ are blue, } f_4 \text{ is purple, and } f_6 \text{ is green.} \end{array} \right\rangle \quad (2.45)$$

and refer to c_3 as the *value* of the string diagram with circular coupon. Altogether we thus obtain a linear map

$$\mathbf{GCal}_{\mathcal{B}}(G) : \quad \mathbf{GCal}_{\mathcal{B}}(K_1 \sqcup K_2) = \mathbf{GCal}_{\mathcal{B}}(K_1) \otimes_{\mathbb{K}} \mathbf{GCal}_{\mathcal{B}}(K_2) \longrightarrow \mathbf{GCal}_{\mathcal{B}}(K_3). \quad (2.46)$$

We declare this linear map to be what the functor $\mathbf{GCal}_{\mathcal{B}}$ maps the morphism $G: K_1 \sqcup K_2 \rightarrow K_3$ to. Moreover, we set $\mathbf{GCal}_{\mathcal{B}}(G_1 \sqcup G_2) := \mathbf{GCal}_{\mathcal{B}}(G_1) \otimes_{\mathbb{k}} \mathbf{GCal}_{\mathcal{B}}(G_2)$. This concludes the definition of $\mathbf{GCal}_{\mathcal{B}}$.

2.18. SUMMARY. *The symmetric monoidal functor $\mathbf{GCal}_{\mathcal{B}}: \mathbf{Crl}_{\mathcal{B}} \rightarrow \mathbf{Vect}_{\mathbb{k}}$ from the symmetric monoidal category $\mathbf{Crl}_{\mathcal{B}}$ of partially \mathcal{B} -colored corollas and graphs to $\mathbf{Vect}_{\mathbb{k}}$ maps objects as in (2.40) and morphisms as in (2.46). The functor $\mathbf{GCal}_{\mathcal{B}}$ completely captures the graphical calculus on disks for pivotal bicategories.*

In the analysis above, as local models for internal vertices of embedded graphs we have only treated corollas on standard disks. But in fact we can make sense of the evaluation of an arbitrary string diagram with circular coupons on D as in (2.45) without reference to embedded disks. In more detail, denote by $V(G)$ the set of internal vertices of a partially \mathcal{B} -colored graph G on D . By choosing for each internal vertex $v \in V(G)$ an embedding of some partially colored corolla K_v , we can turn G into a morphism $\widehat{G}: \bigsqcup_{v \in V(G)} K_v \rightarrow K_G$ in the category $\mathbf{Crl}_{\mathcal{B}}$, where K_G is the unique corolla on D that has the same boundary datum as G . The embeddings of corollas canonically induces an isomorphism $\bigotimes_{v \in V(G)} H_v^{\mathcal{B}} \xrightarrow{\cong} \bigotimes_{v \in V(G)} \mathbf{GCal}_{\mathcal{B}}(K_v)$. We thus obtain a linear map

$$\widehat{\mathbf{GCal}}_{\mathcal{B}}(G): \bigotimes_{v \in V(G)} H_v^{\mathcal{B}} \xrightarrow{\cong} \bigotimes_{v \in V(G)} \mathbf{GCal}_{\mathcal{B}}(K_v) \xrightarrow{\mathbf{GCal}_{\mathcal{B}}(\widehat{G})} \mathbf{GCal}_{\mathcal{B}}(K_G) = H_{v_G}^{\mathcal{B}}, \quad (2.47)$$

with v_G the single vertex of the corolla K_G . Owing to the existence of coherent isomorphisms between the vector spaces $\bigotimes_{v \in V(G)} \mathbf{GCal}_{\mathcal{B}}(K_v)$ arising from different choices of embeddings of corollas, the so obtained map $\widehat{\mathbf{GCal}}_{\mathcal{B}}(G)$ does not depend on that choice. Moreover, $\widehat{\mathbf{GCal}}_{\mathcal{B}}(G)$ is unaffected by any isotopy of G that fixes its boundary. We can therefore give

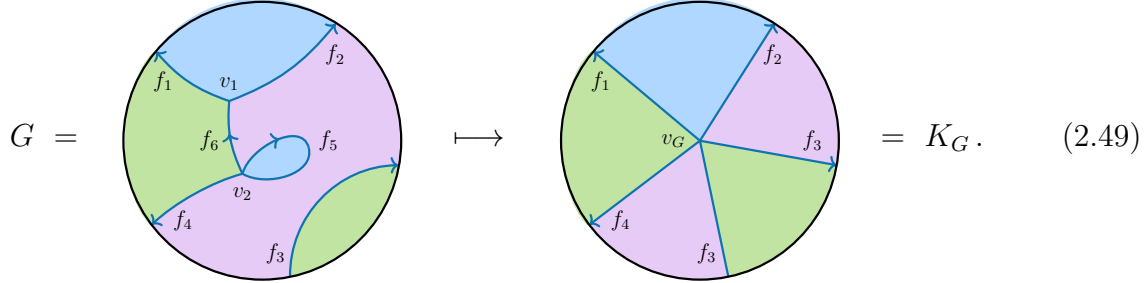
2.19. DEFINITION. *Let G_c be a fully colored graph, with the coloring of its internal vertices given by $c \in \bigotimes_{v \in V(G)} H_v^{\mathcal{B}}$. The value $\langle G_c \rangle \in H_{v_G}^{\mathcal{B}}$ of G_c is defined to be the element*

$$\langle G_c \rangle := \widehat{\mathbf{GCal}}_{\mathcal{B}}(G)(c) \quad (2.48)$$

of the vector space $H_{v_G}^{\mathcal{B}}$ assigned to the vertex of the corolla that corresponds to the boundary datum of G_c .

We think of the value $\langle G \rangle$ of a fully colored graph G as the color for the vertex $v_G \in K_G$ that is obtained by replacing G with the corolla K_G .

2.20. EXAMPLE. The following picture shows a graph G and the corresponding corolla K_G :



$$G = \text{[Diagram of } G \text{]} \mapsto \text{[Diagram of } K_G \text{]} = K_G. \quad (2.49)$$

In this case we have

$$\widehat{\text{GCal}}_{\mathcal{B}}(G) : H_{v_1}^{\mathcal{B}} \otimes_{\mathbb{k}} H_{v_2}^{\mathcal{B}} \longrightarrow H_{v_G}^{\mathcal{B}}. \quad (2.50)$$

$$c = c_1 \otimes_{\mathbb{k}} c_2 \mapsto \langle G_c \rangle,$$

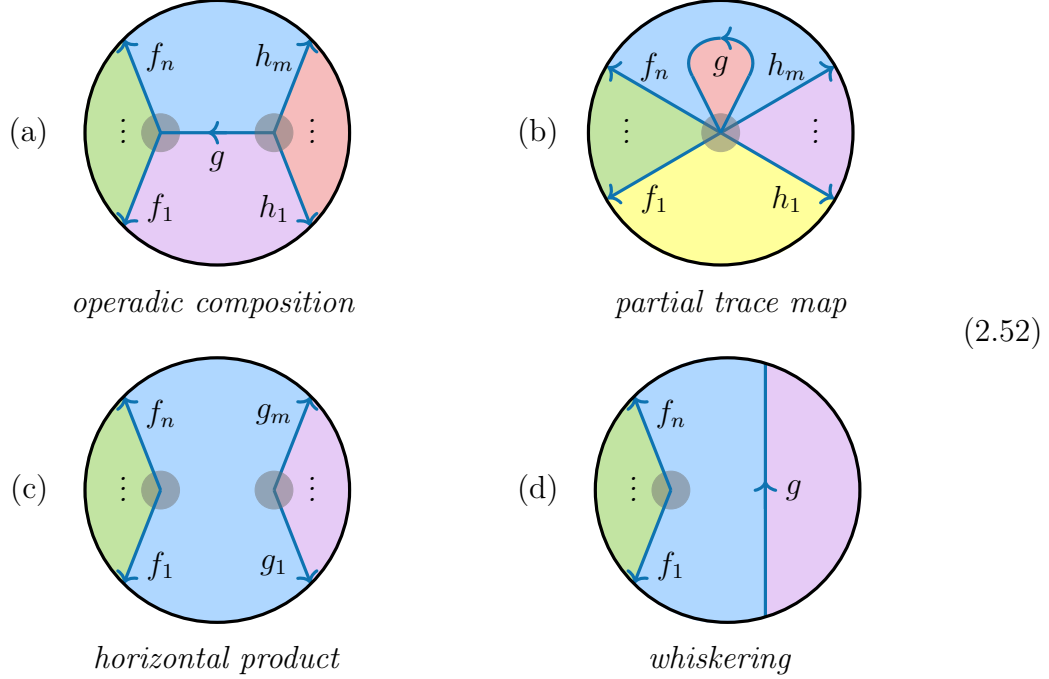
and the value $\langle G_c \rangle \in H_{v_G}^{\mathcal{B}} = \text{GCal}_{\mathcal{B}}(K_G)$ is the same as the one in (2.45), except that now the color $c = c_1 \otimes_{\mathbb{k}} c_2$ is to be understood as living in the vector space $H_{v_1}^{\mathcal{B}} \otimes_{\mathbb{k}} H_{v_2}^{\mathcal{B}}$ associated to the internal vertices of G .

2.21. PRESENTATIONS OF TWO CATEGORIES OF COROLLAS. We can think of the functor $\text{GCal}_{\mathcal{B}} : \text{Crl}_{\mathcal{B}} \rightarrow \text{Vect}_{\mathbb{k}}$ as a rule for assigning a space of morphisms to every partially colored corolla and a *composition map* to every partially colored graph. An important observation is that the composition map obtained from any partially colored graph can be decomposed into a finite sequence of maps each of which is of one of four types, to be introduced in Proposition 2.22. We call an internal edge that connects a pair of distinct internal vertices a *regular* edge, and any other internal edge a *loop*. A partially colored embedded graph in the standard disk is called *trivial* if it is isotopic to a partially colored corolla or does not contain any internal vertex. A morphism in $\text{Crl}_{\mathcal{B}}$ is called trivial if its underlying partially colored embedded graph is trivial. A composition of the type

$$G_2 \diamond G_1 := (K_1 \sqcup K_2 \xrightarrow{G_1 \sqcup \text{id}} K_3 \sqcup K_2 \xrightarrow{G_2} K_4) \quad (2.51)$$

is called a *partial composition* of the morphisms G_1 and G_2 .

2.22. PROPOSITION. *Any non-trivial morphism in \mathbf{Crl}_B can be decomposed into a finite disjoint union of partial compositions of morphisms of the following types:*



PROOF. (1) If the morphism G does not have any regular edge, jump to step (2) (with $G'_n = G$). Otherwise, pick a regular edge e of G and embed the standard disk with image a small disk-shaped neighborhood containing e . This embedding φ pulls back a partially colored graph on the standard disk, which after parametrizing each internal vertex by a corolla gives rise to a morphism G_1 of type (a). We thus have $G = G'_1 \diamond G_1^{(1)}$, where G'_1 is obtained by replacing the part of G that is contained in the image of φ by a corolla. We repeat this procedure until we end up with $G = G'_n \diamond G_n^{(1)} \diamond \cdots \diamond G_1^{(1)}$, where each $G_i^{(1)}$ is of type (a) and G'_n does not contain any regular edge. This expression is reached after finitely many steps, because all graphs are finite and in each step the number of internal vertices decreases. If G'_n is isotopic to a corolla, then the composite $G'_n \diamond G_n^{(1)}$ is of type (a) and we are done. Otherwise we proceed to step (2).

(2) If the morphism G'_n does not have any loop, jump to step (3) (with $G''_m = G'_n$). Otherwise, pick a loop of G'_n , and embed a standard disk to a small neighborhood that contains that loop. This gives $G = G''_1 \diamond G_1^{(2)} \diamond G'_n \diamond \cdots \diamond G_1^{(1)}$ with $G_1^{(2)}$ of type (b). We repeat this procedure until we end up with $G = G''_m \diamond G_m^{(2)} \diamond \cdots \diamond G_1^{(2)} \diamond G'_n \diamond \cdots \diamond G_1^{(1)}$ with each $G_j^{(2)}$ of type (b) and G''_m not containing any loop. If G''_m is isotopic to a corolla, then $G''_m \diamond G_m^{(2)}$ is of type (b) and we are done. Otherwise, G''_m is a union of corollas and edges that do not contain internal vertices, and we proceed to step (3).

(3) If the morphism G''_m contains at most one internal vertex, jump to step (4) (with $G'''_l = G''_m$). Otherwise, pick a pair of internal vertices of G''_m and embed a standard disk to a small neighborhood containing them. This gives $G = G'''_1 \diamond G_1^{(3)} \diamond G''_m \diamond \cdots \diamond G_1^{(2)} \diamond G'_n \diamond$

$\cdots \diamond G_1^{(1)}$ with $G_1^{(3)}$ of type (c). We repeat this procedure until

$$G = G_l''' \diamond G_l^{(3)} \diamond \cdots \diamond G_1^{(3)} \diamond G_m^{(2)} \diamond \cdots \diamond G_1^{(2)} \diamond G_n^{(1)} \diamond \cdots \diamond G_1^{(1)}, \quad (2.53)$$

where each $G_k^{(3)}$ is of type (c) and G_l''' contains a single internal vertex (and neither regular edges nor loops). If G_l''' is isotopic to a corolla, then $G_l''' \diamond G_l^{(3)}$ is of type (c) and we are done. Otherwise, G_l''' consists of a single corolla and multiple edges without internal vertices, and we proceed to step (4).

(4) After consecutively replacing each of the (finitely many) disk-shaped neighborhoods in G_l''' that are of type (d) by a corolla we arrive at the desired decomposition

$$G = G_k^{(4)} \diamond \cdots \diamond G_1^{(4)} \diamond G_l^{(3)} \diamond \cdots \diamond G_1^{(3)} \diamond G_m^{(2)} \diamond \cdots \diamond G_2^{(2)} \diamond G_n^{(1)} \diamond \cdots \diamond G_1^{(1)}, \quad (2.54)$$

with all $G_p^{(4)}$'s of type (d). ■

As indicated in the list (2.52), we refer to the linear maps that are obtained by evaluating the functor $\mathbf{GCal}_{\mathcal{B}}: \mathbf{Crl}_{\mathcal{B}} \rightarrow \mathbf{Vect}_{\mathbb{k}}$ at morphisms of the types (a), (b), (c), and (d) as operadic compositions, partial trace maps, horizontal products, and whiskerings, respectively. Note that in our formulation of the graphical calculus, the horizontal product is *not* a special case of operadic composition, and whiskering is *not* a special case of horizontal product. Owing to Proposition 2.22, the linear map obtained from evaluating $\mathbf{GCal}_{\mathcal{B}}$ at any non-trivial morphism can be written either as a partial composite of such elementary maps or as a tensor product of such partial composites.

Next we consider the subcategory of $\mathbf{Crl}_{\mathcal{B}}$ that contains all objects of $\mathbf{Crl}_{\mathcal{B}}$, but only those morphisms all of whose underlying graphs are connected in the disks they are embedded in. We denote this category by $\mathbf{Crl}_{\mathcal{B}}^{\text{conn}}$. For instance, a finite disjoint union of morphisms of type (a) belongs to $\mathbf{Crl}_{\mathcal{B}}^{\text{conn}}$, whereas a morphism of type (c) does not. We have

2.23. COROLLARY. *The subcategory $\mathbf{Crl}_{\mathcal{B}}^{\text{conn}}$ is generated, under monoidal product (i.e., disjoint union) and partial composition by the trivial morphisms with connected embedded graphs and by the morphisms of the types (a) and (b).*

We denote by

$$\mathbf{GCal}_{\mathcal{B}}^{\text{conn}}: \mathbf{Crl}_{\mathcal{B}}^{\text{conn}} \hookrightarrow \mathbf{Crl}_{\mathcal{B}} \xrightarrow{\mathbf{GCal}_{\mathcal{B}}} \mathbf{Vect}_{\mathbb{k}} \quad (2.55)$$

the restriction of the functor $\mathbf{GCal}_{\mathcal{B}}$ to the subcategory $\mathbf{Crl}_{\mathcal{B}}^{\text{conn}}$. According to Corollary 2.23, evaluating $\mathbf{GCal}_{\mathcal{B}}^{\text{conn}}$ at any non-trivial morphism in its domain produces a partial composite of operadic compositions and partial trace maps, or a tensor product of such partial composites.

2.24. REMARK. Inspiration for the constructions in this section comes from the study [Cos] of (symmetric) operads, cyclic operads and modular operads as symmetric monoidal functors defined on suitable categories of graphs. Here we work in a somewhat different setting: we deal with categories of graphs embedded in disks instead of graphs, and there is an additional coloring on the operads. When the bicategory \mathcal{B} has a single object, i.e. \mathcal{B}

describe the evaluation of the functor F on a 2-morphism, we use (similarly as in e.g. [Me]) the symbol F followed by a *window* that encloses the string diagram that expresses the 2-morphism. The defining properties of lax functoriality then amount to the following equalities of string diagrams:

1. *Naturality* is expressed as the family of equalities

$$\begin{array}{c}
 F(f'g') \\
 \uparrow \\
 F \left[\begin{array}{|c|c|} \hline \alpha & \beta \\ \hline \end{array} \right] \\
 \uparrow \\
 F(fg) \\
 \uparrow \\
 \begin{array}{cc} Ff & Fg \end{array}
 \end{array}
 =
 \begin{array}{c}
 F(f'g') \\
 \uparrow \\
 \begin{array}{cc} Ff' & Fg' \end{array} \\
 \uparrow \quad \uparrow \\
 \begin{array}{|c|c|} \hline F\alpha & F\beta \\ \hline \end{array} \\
 \uparrow \quad \uparrow \\
 \begin{array}{cc} Ff & Fg \end{array}
 \end{array}
 \quad (2.59)$$

for all objects $a, b, c \in \mathcal{B}$, 1-morphisms $f, f' \in \mathcal{B}(a, b)$ and $g, g' \in \mathcal{B}(b, c)$, and 2-morphisms $\alpha: f \Rightarrow f'$ and $\beta: g \Rightarrow g'$.

2. *Lax associativity* is expressed as the equalities

$$\begin{array}{c}
 F(fgh) \\
 \uparrow \\
 F(fg) \quad \uparrow \\
 \begin{array}{cc} Ff & Fg \end{array} \quad \uparrow \\
 \begin{array}{ccc} Ff & Fg & Fh \end{array}
 \end{array}
 =
 \begin{array}{c}
 F(fgh) \\
 \uparrow \\
 \begin{array}{cc} \uparrow & F(gh) \end{array} \\
 \begin{array}{ccc} Ff & Fg & Fh \end{array}
 \end{array}
 \quad (2.60)$$

for all composable triples f, g, h of 1-morphisms in \mathcal{B} .

3. *Lax left and right unity* are expressed as

$$\begin{array}{c}
 Ff \\
 \uparrow \\
 F\text{id}_a \quad \uparrow \\
 \begin{array}{cc} \circ & Ff \end{array}
 \end{array}
 =
 \begin{array}{c}
 Ff \\
 \uparrow \\
 Ff
 \end{array}
 =
 \begin{array}{c}
 Ff \\
 \uparrow \\
 \begin{array}{cc} Ff & \circ \end{array}
 \end{array}
 \quad F\text{id}_b \quad (2.61)$$

for all pairs of objects $a, b \in \mathcal{B}$ and all 1-morphisms $f \in \mathcal{B}(a, b)$.

For every composable string $a_1 \xrightarrow{f_1} a_2 \xrightarrow{f_2} \dots \xrightarrow{f_n} a_{n+1}$ of 1-morphisms in \mathcal{B} , due to the lax associativity there is a unique 2-morphism

$$F_{f_1, \dots, f_n}^{(n)} : Ff_1 \dots Ff_n \Rightarrow F(f_1 \dots f_n), \quad (2.62)$$

If any of the 1-morphisms is trivial, we need to insert in addition the lax and/or oplax unity constraints accordingly. For instance, if the 1-morphisms f_1 and g_n are trivial, then α^F is defined to be

$$\begin{array}{c}
 Fg_1 \cdots \\
 \cdots \\
 \text{---} \circ \text{---} \\
 \text{---} Fg_n \\
 \uparrow \\
 \boxed{F\alpha} \\
 \uparrow \\
 Ff_1 \quad \cdots \quad F(f_1 \cdots f_m) \\
 \text{---} \circ \text{---} \\
 \cdots \quad \cdots \quad Ff_m
 \end{array} \tag{2.66}$$

2.27. REMARK. According to Definition 2.26, the precise form of the conjugation depends on the notion of a trivial 1-morphism. Conventionally, by the trivial 1-endomorphism of an object a in a bicategory \mathcal{B} one means the identity morphism $\text{id}_a \in \mathcal{B}(a, a)$. However, in order to correctly describe the functoriality of graphical calculi in terms of a monoidal natural transformation (see Theorem 2.30 below), we need to sharpen this notion of triviality: in our context, a 1-morphism is trivial iff it is an identity 1-morphism *and* it is not an edge color.

For a generic functor with both lax and oplax structures there is no reason for the conjugation to preserve compositions or partial traces, nor to respect the coherent isomorphisms between 2-hom spaces that are related by dualities. To deal with a functor that does satisfy such relations we need to impose appropriate properties on the lax and oplax structures. Specifically, we need the following notions. (In the 1-object case, i.e. for monoidal categories, these have been studied in e.g. [McS] and have found applications in constructions of topological field theories [Mu].)

2.28. DEFINITION. Let $F \equiv (F, F^{(2)}, F^{(0)}, F_{(2)}, F_{(0)}): \mathcal{B} \rightarrow \mathcal{B}'$ be a functor between two strictly pivotal bicategories that is equipped with lax and oplax structures.

- (i) F is called rigid if F -conjugation preserves the units and counits of the duals, i.e. if for all $a, b \in \mathcal{B}$ and all $f \in \mathcal{B}(a, b)$ we have $F(f^\vee) = (Ff)^\vee$ and

$$\begin{array}{c}
 \circ \text{---} \text{Fid}_b \\
 \uparrow \\
 \boxed{F} \\
 \text{---} f \text{---} \\
 \uparrow \\
 F(f^\vee f) \\
 \uparrow \\
 F(f^\vee) \quad Ff
 \end{array} = \begin{array}{c}
 \text{---} \\
 \uparrow \\
 Ff \quad Ff
 \end{array} \tag{2.67}$$

(the ‘antenna’ on the left hand side stands for the oplax unit constraint for the object b) as well as similar conditions for the coevaluations.

- (ii) F is called separable if for every pair $a \xrightarrow{f} b \xrightarrow{g} c$ of composable 1-morphisms the lax and oplax structures are related by

$$\begin{array}{c}
 F(fg) \\
 \uparrow \\
 \text{F(}f\text{) } \text{---} \text{F(}g\text{)} \\
 \uparrow \\
 F(fg)
 \end{array}
 =
 \begin{array}{c}
 F(fg) \\
 \uparrow \\
 F(fg)
 \end{array}
 \quad (2.68)$$

- (iii) F is called Frobenius if for every triple $a \xrightarrow{f} b \xrightarrow{g} c \xrightarrow{h} d$ of composable 1-morphisms the two compatibility relations

$$\begin{array}{c}
 Ff \quad F(gh) \\
 \uparrow \quad \uparrow \\
 \text{F(}g\text{)} \\
 \uparrow \\
 F(fg) \quad Fh
 \end{array}
 =
 \begin{array}{c}
 Ff \quad F(gh) \\
 \uparrow \quad \uparrow \\
 \text{F(}fgh\text{)} \\
 \uparrow \\
 F(fg) \quad Fh
 \end{array}
 \quad (2.69)$$

and

$$\begin{array}{c}
 F(fg) \quad Fh \\
 \uparrow \quad \uparrow \\
 \text{F(}g\text{)} \\
 \uparrow \\
 Ff \quad F(gh)
 \end{array}
 =
 \begin{array}{c}
 F(fg) \quad Fh \\
 \uparrow \quad \uparrow \\
 \text{F(}fgh\text{)} \\
 \uparrow \\
 Ff \quad F(gh)
 \end{array}
 \quad (2.70)$$

between the lax and oplax structure are fulfilled.

2.29. REMARK. It is worth stressing that the ‘Frobenius functors’ considered here are 2-functors between bicategories. This notion should not be confused with the one of a Frobenius functor between categories, defined as a functor admitting a two-sided adjoint. A rigid separable Frobenius functor equips the image of the identity 1-morphism of any object in the domain bicategory with the structure of a Δ -separable symmetric Frobenius algebra, equips the image of every 1-morphism in the domain bicategory with the structure of a bimodule, and equips the conjugate of every 2-morphism in the domain bicategory with the structure of a bimodule morphism. (Here the notions of bimodule and bimodule morphism are generalized to the bicategorical setting, compare e.g. [CaR, Sect. 2.2].)

We will now show that a rigid separable Frobenius functor *almost* preserves the graphical calculus on disks. Let K be a partially \mathcal{B} -colored corolla on the standard disk and let $F \equiv (F, F^{(2)}, F^{(0)}, F_{(2)}, F_{(0)}): \mathcal{B} \rightarrow \mathcal{B}'$ be a rigid separable Frobenius functor. The map of objects and the local functors that are entailed by F give rise to a symmetric monoidal functor $F_*: \mathbf{Crl}_{\mathcal{B}}^{\text{conn}} \rightarrow \mathbf{Crl}_{\mathcal{B}'}$ by accordingly changing the colors of the patches and of the edges of the partially colored graphs. Moreover, since the change of colors does not affect the connectedness of the embedded graphs, this functor F_* restricts to a functor $F_*: \mathbf{Crl}_{\mathcal{B}}^{\text{conn}} \rightarrow \mathbf{Crl}_{\mathcal{B}'}^{\text{conn}}$. Restricting to the subcategory $\mathbf{Crl}_{\mathcal{B}}^{\text{conn}}$ turns out to have an important consequence:

2.30. THEOREM. *Let $F: \mathcal{B} \rightarrow \mathcal{B}'$ be a rigid Frobenius functor between two strictly pivotal bicategories. The F -conjugation introduced in Definition 2.26 canonically induces a monoidal natural transformation*

$$\begin{array}{ccc}
 \mathbf{Crl}_{\mathcal{B}}^{\text{conn}} & & \\
 \downarrow F_* & \searrow \text{GCal}_{\mathcal{B}}^{\text{conn}} & \\
 & (-)^F & \text{Vect}_{\mathbf{k}} \\
 & \nearrow \text{GCal}_{\mathcal{B}'}^{\text{conn}} & \\
 \mathbf{Crl}_{\mathcal{B}'}^{\text{conn}} & &
 \end{array} \quad (2.71)$$

(to which we still refer as F -conjugation). Its component at a corolla $K \in \text{GCal}_{\mathcal{B}}^{\text{conn}}$ is given by

$$\begin{aligned}
 (-)^F_K: \text{GCal}_{\mathcal{B}}^{\text{conn}}(K) &\xrightarrow{p_v^k} \widehat{h}_v^{\mathcal{B}}(k) \\
 &\xrightarrow{(-)^F} \widehat{h}_{v'}^{\mathcal{B}'}(k) \\
 &\xrightarrow{(p_{v'}^k)^{-1}} \text{GCal}_{\mathcal{B}'}^{\text{conn}}(F_* K)
 \end{aligned} \quad (2.72)$$

for any choice of polarization k on the vertex $v \in K$, where p_v^k is the isomorphism defined in formula (2.30) and where the labels of all edges are treated as non-trivial 1-morphisms, while $(-)^F_{K_1 \sqcup K_2} := (-)^F_{K_1} \otimes_{\mathbf{k}} (-)^F_{K_2}$.

PROOF. We first show that (2.72) is well defined, i.e. does not depend on the choice of polarization. This amounts to showing that the F -conjugation commutes with all canonical isomorphisms between 2-hom spaces that are related by duality. We present the argument for one specific case – its generalization to other cases is straightforward: we claim that the diagram

$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{B}(a,a)}(\text{id}_a, f \star g) & \xrightarrow{\cong} & \text{Hom}_{\mathcal{B}(a,b)}(g^\vee, f) \\
 \downarrow (-)^F & & \downarrow (-)^F \\
 \text{Hom}_{\mathcal{B}'(Fa, Fa)}(\text{id}_{Fa}, Ff \star Fg) & \xrightarrow{\cong} & \text{Hom}_{\mathcal{B}'(Fa, Fb)}((Fg)^\vee, Ff)
 \end{array} \quad (2.73)$$

commutes for all objects $a, b \in \mathcal{B}$ and 1-morphisms $f \in \mathcal{B}(a, b)$ and $g \in \mathcal{B}(b, a)$ (note that by rigidity we have $F(g^\vee) = (Fg)^\vee$). To prove this claim we pick a 2-morphism $\alpha \in \text{Hom}_{\mathcal{B}(a,a)}(\text{id}_a, fg)$. Since the 1-morphisms f and g come from the coloring of edges, they are regarded as non-trivial, and they will be treated as such by the F -conjugation. Now when first going right and then downwards in the square (2.73) we get

$$\begin{array}{c} f \quad g \\ \uparrow \quad \uparrow \\ \boxed{\alpha} \end{array} \mapsto \begin{array}{c} f \\ \uparrow \\ \boxed{\alpha} \end{array} \begin{array}{c} \text{curved arrow} \\ \downarrow \\ g \end{array} \mapsto \begin{array}{c} Ff \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ Fg \end{array} \quad (2.74)$$

The graph on the right hand side of (2.74) can be rewritten as

$$\begin{array}{c} Ff \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ Fg \end{array} = \begin{array}{c} Ff \quad \circ F\text{id}_b \\ \uparrow \quad \uparrow \\ F \boxed{\alpha} \\ \uparrow \\ F(fgg^\vee) \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ F\text{id}_a \quad \circ Fg \end{array} = \begin{array}{c} Ff \quad \circ F\text{id}_b \\ \uparrow \quad \uparrow \\ F \boxed{\alpha} \\ \uparrow \\ F(gg^\vee) \\ \uparrow \\ F(fg) \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ F\text{id}_a \quad \circ Fg \end{array} \quad (2.75)$$

$$\begin{array}{c} Ff \quad \circ F\text{id}_b \\ \uparrow \quad \uparrow \\ F \boxed{\alpha} \\ \uparrow \\ F(gg^\vee) \\ \uparrow \\ F(fg) \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ F\text{id}_a \quad \circ Fg \end{array} \stackrel{(2.69)}{=} \begin{array}{c} Ff \quad \circ F\text{id}_b \\ \uparrow \quad \uparrow \\ F \boxed{\alpha} \\ \uparrow \\ F(gg^\vee) \\ \uparrow \\ F(fg) \\ \uparrow \\ F \boxed{\alpha} \\ \downarrow \\ F\text{id}_a \quad \circ Fg \end{array}$$

The final graph in (2.75) is the same as the one obtained when tracing the diagram (2.73) first downwards and then right.

Next we address naturality. The naturality squares for the trivial morphisms in $\mathbf{CrlI}_{\mathcal{B}}^{\text{conn}}$ commute trivially. In view of Corollary 2.23 we need to verify naturality only for morphisms of the types (a) and (b), i.e. we must show that F -conjugation commutes with operadic compositions and partial trace maps. First consider operadic compositions. To keep the discussion short we pick a specific morphism of type (a), say $G: K_1 \sqcup K_2 \rightarrow K_3$

given by

$$G = \begin{array}{c} \text{A circle divided horizontally into a light blue upper half and a light green lower half.} \\ \text{A horizontal blue line with arrows at both ends passes through the center.} \\ \text{Two grey dots are on this line, one in the blue half and one in the green half.} \\ \text{Below the line, the segments are labeled } f, g, h \text{ from left to right.} \end{array} \quad (2.76)$$

with $f, g, h \in \mathcal{B}(a, b)$. The naturality square for (2.76) reads

$$\begin{array}{ccc} \mathrm{GCal}_{\mathcal{B}}^{\mathrm{conn}}(K_1) \otimes_{\mathbb{k}} \mathrm{GCal}_{\mathcal{B}}^{\mathrm{conn}}(K_2) & \xrightarrow{\mathrm{GCal}_{\mathcal{B}}^{\mathrm{conn}}(G)} & \mathrm{GCal}_{\mathcal{B}}^{\mathrm{conn}}(K_3) \\ \downarrow (-)_{K_1}^F \otimes_{\mathbb{k}} (-)_{K_2}^F & & \downarrow (-)_{K_3}^F \\ \mathrm{GCal}_{\mathcal{B}'}^{\mathrm{conn}}(F_*K_1) \otimes_{\mathbb{k}} \mathrm{GCal}_{\mathcal{B}'}^{\mathrm{conn}}(F_*K_2) & \xrightarrow{\mathrm{GCal}_{\mathcal{B}'}^{\mathrm{conn}}(F_*G)} & \mathrm{GCal}_{\mathcal{B}'}^{\mathrm{conn}}(F_*K_3) \end{array} \quad (2.77)$$

We choose a polarization for each corolla K_i , as well as elements $\alpha \in \widehat{h}_{v_1}^{\mathcal{B}}(k_1) = \mathrm{Hom}_{\mathcal{B}(a,a)}(\mathrm{id}_a, fg^\vee)$ and $\beta \in \widehat{h}_{v_2}^{\mathcal{B}}(k_2) = \mathrm{Hom}_{\mathcal{B}(a,a)}(\mathrm{id}_a, gh)$ in the respective 2-hom spaces for K_1 and K_2 . The following calculation shows that the two paths in the square (2.77) give the same element in $\widehat{h}_{v_3}^{\mathcal{B}}(k_3) = \mathrm{Hom}_{\mathcal{B}(a,a)}(\mathrm{id}_a, fh)$:

$$(\alpha \star_g \beta)^F = \begin{array}{c} \text{Diagram 1: A box containing } \alpha \text{ and } \beta. \text{ An arrow } F\mathrm{id}_a \text{ enters from the bottom.} \\ \text{Two arrows } F\alpha \text{ and } F\beta \text{ exit from the top.} \\ \text{An arrow } F(fh) \text{ enters from the top.} \\ \text{Two arrows } Ff \text{ and } Fh \text{ exit from the top.} \end{array} = \begin{array}{c} \text{Diagram 2: A box containing } \alpha \text{ and } \beta. \text{ An arrow } F\mathrm{id}_a \text{ enters from the bottom.} \\ \text{Two arrows } F\alpha \text{ and } F\beta \text{ exit from the top.} \\ \text{An arrow } F(fg^\vee gh) \text{ enters from the top.} \\ \text{An arrow } F(fh) \text{ enters from the top.} \\ \text{Two arrows } Ff \text{ and } Fh \text{ exit from the top.} \end{array}$$

$$\begin{aligned}
&= \text{Diagram 1} \stackrel{(2.70)}{=} \text{Diagram 2} \\
&\stackrel{(2.69)}{=} \text{Diagram 3} \stackrel{(2.67)}{=} \text{Diagram 4} = \alpha^F \star_{Fg} \beta^F.
\end{aligned} \tag{2.78}$$

The diagrams are string diagrams representing morphisms in a 2-representation. They consist of blue strands with various labels: Ff , Fh , $F(fg^\vee g)$, $F(fg^\vee gh)$, $F(fg^\vee)$, $F(gh)$, Fid_a , $F\alpha$, $F\beta$, and Fg . Some strands are enclosed in boxes. The equations (2.70), (2.69), and (2.67) represent specific 2-morphisms or identities between these diagrams.

(Here $(-)^F$ denotes F -conjugation (see Definition 2.26), while \star_g stands for horizontal composition along the 1-morphism g .)

Next consider partial trace maps. Let $G: K_1 \rightarrow K_2$ be the morphism of type (b) in $\mathbf{Crl}^{\text{conn}}_{\mathcal{B}}$ given by

$$G = \text{Diagram of } G \tag{2.79}$$

The diagram for G is a green circle containing a blue teardrop shape labeled g . A blue arrow labeled f points from the center of the circle to the left.

with $f \in \mathcal{B}(a, a)$ and $g \in \mathcal{B}(b, a)$, for which the naturality square is

$$\begin{array}{ccc}
\text{GCal}_{\mathcal{B}}^{\text{conn}}(K_1) & \xrightarrow{\text{GCal}_{\mathcal{B}}^{\text{conn}}(G)} & \text{GCal}_{\mathcal{B}}^{\text{conn}}(K_2) \\
(-)^F_{K_1} \downarrow & & \downarrow (-)^F_{K_2} \\
\text{GCal}_{\mathcal{B}'}^{\text{conn}}(F_* K_1) & \xrightarrow{\text{GCal}_{\mathcal{B}'}^{\text{conn}}(F_* G)} & \text{GCal}_{\mathcal{B}'}^{\text{conn}}(F_* K_2)
\end{array} \tag{2.80}$$

We select polarizations k_1 and k_2 for the corollas K_1 and K_2 and choose a 2-morphism $\alpha \in \widehat{h}_{v_1}^{\mathcal{B}}(k_1) = \text{Hom}_{\mathcal{B}(a,a)}(\text{id}_a, fg^\vee g)$, and we consider the case that $\widehat{h}_{v_2}^{\mathcal{B}}(k_2) = \text{Hom}_{\mathcal{B}(a,a)}(\text{id}_a, f)$. Then the following calculation shows that taking the partial trace commutes with the F -

conjugation:

$$\begin{aligned}
 (\mathrm{tr}_g \alpha)^F &= \text{Diagram 1} = \text{Diagram 2} \stackrel{(2.68)}{=} \text{Diagram 3} \\
 &= \text{Diagram 4} = \text{Diagram 5} = \mathrm{tr}_{Fg} \alpha^F.
 \end{aligned} \tag{2.81}$$

The diagrams in (2.81) are as follows:

- Diagram 1:** A vertical blue line starts from a small circle at the bottom, goes up through a box labeled α , and continues upwards. The top part of the line is labeled Ff . The bottom part is labeled Fid_a .
- Diagram 2:** A vertical blue line starts from a small circle at the bottom, goes up through a box labeled α , then through a box labeled $F(fg^\vee g)$, and continues upwards. The top part is labeled Ff . The bottom part is labeled Fid_a .
- Diagram 3:** A vertical blue line starts from a small circle at the bottom, goes up through a box labeled $F\alpha$, then through a circle labeled $F(fg^\vee)$, and continues upwards. The top part is labeled Ff . The bottom part is labeled Fid_a .
- Diagram 4:** A vertical blue line starts from a small circle at the bottom, goes up through a box labeled $F\alpha$, then through a box labeled $F(fg^\vee)$, and continues upwards. The top part is labeled Ff . The bottom part is labeled Fid_a .
- Diagram 5:** A vertical blue line starts from a small circle at the bottom, goes up through a box labeled $F\alpha$, then through a box labeled Fg , and continues upwards. The top part is labeled Ff . The bottom part is labeled Fid_a .

In conclusion, we did construct a natural transformation $(-)^F : \mathrm{GCal}_{\mathcal{B}}^{\mathrm{conn}} \Rightarrow \mathrm{GCal}_{\mathcal{B}'}^{\mathrm{conn}} \circ F_*$. Moreover, this natural transformation is evidently monoidal. \blacksquare

It is worth stressing that a rigid separable Frobenius functor does *not*, in general, preserve the *entire* graphical calculus on disks. To understand what goes wrong, consider a morphism $G : K_1 \sqcup K_2 \rightarrow K_3$ of horizontal-product type (c), say

$$G = \text{Diagram 6} \tag{2.82}$$

The diagram in (2.82) is a circle divided into three vertical regions: green on the left, blue in the middle, and purple on the right. In the green region, there are two blue arrows pointing upwards, labeled f_1 and f_2 . In the blue region, there are two blue arrows pointing upwards, labeled g_1 and g_2 . In the purple region, there are two blue arrows pointing upwards, labeled g_1 and g_2 . There are two small black dots in the blue region, one near each arrow.

with $f_1, f_2 \in \mathcal{B}(a, b)$ and $g_1, g_2 \in \mathcal{B}(b, c)$. Upon choosing appropriate polarizations and 2-morphisms, G yields a horizontal product

$$\text{Diagram 7} \tag{2.83}$$

The diagram in (2.83) is a rectangle divided into three vertical regions: green on the left, blue in the middle, and purple on the right. In the green region, there are two blue arrows pointing upwards, labeled f_1 and f_2 . In the blue region, there are two blue arrows pointing upwards, labeled g_1 and g_2 . In the purple region, there are two blue arrows pointing upwards, labeled g_1 and g_2 . There are two small black dots in the blue region, one near each arrow.

The F -conjugate of this horizontal product is

$$\begin{array}{c}
 Ff_2 \uparrow \quad \uparrow Fg_2 \\
 \quad \quad \quad \cup \\
 (\alpha \star \beta)^F = F \left[\begin{array}{|c|c|} \hline \alpha & \beta \\ \hline \end{array} \right] = \begin{array}{c} Ff_2 \uparrow \quad \uparrow Fg_2 \\ \quad \quad \quad \cup \\ F(f_2 g_2) \\ \quad \quad \quad \cup \\ Ff_2 \uparrow \quad \uparrow Fg_2 \\ \quad \quad \quad \cup \\ \begin{array}{|c|c|} \hline F\alpha & F\beta \\ \hline \end{array} \\ \quad \quad \quad \cup \\ Ff_1 \uparrow \quad \uparrow Fg_1 \end{array} \\
 Ff_1 \uparrow \quad \uparrow Fg_1
 \end{array} \quad (2.84)$$

Thus the so obtained 2-morphism differs from $\alpha^F \star \beta^F = F\alpha \star F\beta$ by

$$\begin{array}{c}
 Ff_2 \uparrow \quad \uparrow Fg_2 \\
 \quad \quad \quad \cup \\
 \quad \quad \quad \uparrow \\
 Ff_2 \uparrow \quad \uparrow Fg_2
 \end{array} \quad (2.85)$$

which owing to the separability of F is an idempotent on the space of 2-endomorphism of $Ff_2 \star Fg_2$. This observation motivates

2.31. DEFINITION. A functor $F: \mathcal{B} \rightarrow \mathcal{B}'$ equipped with both lax and oplax structures for which the lax functoriality constraint is the inverse of its oplax functoriality constraint is called *strongly separable*.

From the previous observations it follows immediately that a *rigid strongly separable Frobenius functor* preserves *horizontal products* as well as *whiskerings*, and thus preserves the complete graphical calculus on disks of its domain bicategory. However, there is in fact no need for a separate notion of a strongly separable Frobenius functor. Recall that a *pseudofunctor* can be regarded as a functor with lax and oplax structures that are mutually inverse. We have

2.32. LEMMA. For F a functor between bicategories that is equipped with lax and oplax structures, we have:

- (i) F is strongly separable if and only if it is a pseudofunctor.
- (ii) If F is a pseudofunctor, then it is Frobenius.

PROOF. (i) A pseudofunctor is strongly separable by definition. To see the converse, assume that F is strongly separable. Then what remains to be shown is that $F^{(0)}$ and

$F_{(0)}$ are mutually inverse. Now for any object $a \in \mathcal{B}$ we have

$$\begin{array}{c} \text{Fid}_a \\ \downarrow \\ \circ \\ \downarrow \\ \circ \\ \text{Fid}_a \end{array} = \begin{array}{c} \text{Fid}_a \\ \downarrow \\ \circ \\ \text{Fid}_a \end{array} \quad \text{Fid}_a \quad \downarrow \quad \circ \\
 = \begin{array}{c} \text{Fid}_a \\ \downarrow \\ \circ \\ \text{Fid}_a \end{array} = \begin{array}{c} \text{Fid}_a \\ \downarrow \\ \circ \\ \text{Fid}_a \end{array} = \begin{array}{c} \text{Fid}_a \\ \downarrow \\ \circ \\ \text{Fid}_a \end{array} \quad (2.86)$$

where the premise that F is strongly separable is used in the second equality, showing that $F_a^{(0)} \circ F_{(0)a} = \text{id}_{F_a}$. Moreover, for a 1-morphism $f \in \mathcal{B}(a, b)$ with non-zero 2-endomorphism space we have

$$\begin{array}{c} Ff \\ \uparrow \\ Ff \end{array} = \begin{array}{c} Ff \\ \uparrow \\ \circ \\ \uparrow \\ \circ \\ Ff \end{array} = \begin{array}{c} Ff \\ \uparrow \\ \circ \\ \uparrow \\ \circ \\ Ff \end{array} = \begin{array}{c} Ff \\ \uparrow \\ \circ \\ \uparrow \\ \circ \\ Ff \end{array} \quad (2.87)$$

so that also $F_{(0)a} \circ F_a^{(0)} = \text{id}_{F_a}$. (If there does not exist such a 1-morphism f for any $b \in \mathcal{B}$, then the desired statement holds true trivially.)

(ii) We have

$$\begin{array}{c} F(fg) \\ \uparrow \\ Ff \end{array} \quad \begin{array}{c} Fh \\ \uparrow \\ F(gh) \end{array} = \begin{array}{c} F(fg) \\ \uparrow \\ Ff \end{array} \quad \begin{array}{c} Fh \\ \uparrow \\ F(gh) \end{array} = \begin{array}{c} F(fg) \\ \uparrow \\ Ff \end{array} \quad \begin{array}{c} Fh \\ \uparrow \\ F(gh) \end{array} = \begin{array}{c} F(fg) \\ \uparrow \\ Ff \end{array} \quad \begin{array}{c} Fh \\ \uparrow \\ F(gh) \end{array} \quad (2.88)$$

where associativity is used in the second equality, while strong separability (which holds by part (i)) is used in the first and third. Thus (2.70) is satisfied. The proof of (2.69) is analogous. ■

In particular, a strongly separable Frobenius functor is just the same as a pseudofunctor. Now recall the notion of F -conjugation; when restricting to the stronger class of functors consisting of rigid pseudofunctors, in Theorem 2.30 we can drop the restriction on the allowed corollas, so that it is sharpened to

2.33. COROLLARY. Let $F: \mathcal{B} \rightarrow \mathcal{B}'$ be a rigid pseudofunctor between strictly pivotal bicategories. The F -conjugation (defined in Theorem 2.30) canonically extends to a monoidal natural transformation

$$\begin{array}{ccc}
 \mathrm{Crl}_{\mathcal{B}} & \xrightarrow{\mathrm{GCal}_{\mathcal{B}}} & \mathrm{Vect}_{\mathbb{K}} \\
 F_* \downarrow & \swarrow (-)^F & \nearrow \mathrm{GCal}_{\mathcal{B}'} \\
 \mathrm{Crl}_{\mathcal{B}'} & &
 \end{array} \quad (2.89)$$

We end this section with an example of a rigid separable Frobenius functor that plays an important role in the application to conformal quantum field theories that we will discuss in Section 4.

2.34. EXAMPLE. Let \mathcal{C} be a strictly pivotal fusion category and $\mathcal{F}r(\mathcal{C})$ the strictly pivotal bicategory of simple special symmetric Frobenius algebras in \mathcal{C} (see Example 2.5). Recall that \mathcal{C} can be viewed as a strictly pivotal bicategory \mathcal{BC} (its delooping, see Example 2.4) with a single object $*$. Consider the functor

$$\mathcal{U}: \mathcal{F}r(\mathcal{C}) \longrightarrow \mathcal{BC} \quad (2.90)$$

that is defined by

$$\begin{array}{ccc}
 \begin{array}{c} A \\ \left(\begin{array}{c} \xRightarrow{\alpha} \end{array} \right) \\ X \quad Y \\ \left(\begin{array}{c} \downarrow \end{array} \right) \\ B \end{array} & \xrightarrow{\mathcal{U}} & \begin{array}{c} * \\ \left(\begin{array}{c} \xRightarrow{\dot{\alpha}} \end{array} \right) \\ \dot{X} \quad \dot{Y} \\ \left(\begin{array}{c} \downarrow \end{array} \right) \\ * \end{array}
 \end{array} \quad (2.91)$$

for Frobenius algebras $A, B \in \mathcal{F}r(\mathcal{C})$, A - B -bimodules X, Y and a bimodule morphism $\alpha: X \Rightarrow Y$, i.e. \mathcal{U} sends every object in $\mathcal{F}r(\mathcal{C})$ to the sole object $*$ in \mathcal{BC} and every bimodule and bimodule morphism to their underlying object and morphism in \mathcal{C} , respectively. In the sequel, we suppress the dot and just use the same symbol for bimodules and bimodule morphisms as for their underlying objects and morphisms.

The functor \mathcal{U} is canonically a rigid separable Frobenius functor, with the components

$$\mathcal{U}_{X,Y}^{(2)}: X \otimes Y \longrightarrow X \otimes_B Y \quad \text{and} \quad \mathcal{U}_{(2)X,Y}: X \otimes_B Y \longrightarrow X \otimes Y \quad (2.92)$$

of the lax and oplax functoriality constraints at $(X, Y) \in \mathcal{F}r(\mathcal{C})(A, B) \times \mathcal{F}r(\mathcal{C})(B, C)$ given by the splitting of the idempotent that realizes the tensor product \otimes_B as a retract

of \otimes , i.e. $\mathcal{U}_{X,Y}^{(2)} \circ \mathcal{U}_{(2)X,Y} = \text{id}_{X \otimes_B Y}$ and

$$\mathcal{U}_{(2)X,Y} \circ \mathcal{U}_{X,Y}^{(2)} = \begin{array}{c} \begin{array}{cc} X & Y \\ \uparrow & \uparrow \\ \text{---} \cup \text{---} \\ \uparrow \\ X \otimes_B Y \\ \downarrow \\ \text{---} \cup \text{---} \\ \uparrow & \uparrow \\ X & Y \end{array} \end{array} = \begin{array}{c} \begin{array}{cc} X & Y \\ \uparrow & \uparrow \\ \text{---} \text{---} \text{---} \text{---} \\ \uparrow & \downarrow \\ \text{---} \text{---} \text{---} \text{---} \\ \uparrow & \downarrow \\ X & Y \end{array} \end{array} \quad (2.93)$$

and with the lax and oplax unity constraints given by the units and counits of the Frobenius algebras. Note that to satisfy the axioms, associators and unitors must be inserted accordingly. Also, for the example to be non-trivial, a crucial requirement is that special symmetric Frobenius algebras other than the monoidal unit exist in \mathcal{C} .

3. String-net models based on pivotal bicategories

In Section 2 we have developed the graphical calculus on a canvas that is homeomorphic to a disk. It is natural to seek an extension of this calculus for which the canvas can have non-trivial topology. String-net models for surfaces provide such an extension. Throughout this section \mathcal{B} is a strictly pivotal bicategory. In addition we assume that \mathcal{B} is small and locally small, i.e. its objects form a set and all its hom-categories are small.

3.1. BICATEGORICAL STRING-NET SPACES. Recall from Definition 2.10 the notion of a \mathcal{B} -boundary datum on a compact oriented 1-manifold. Let Σ be a compact oriented surface and let \mathbf{b} be a \mathcal{B} -boundary datum on $\partial\Sigma$. A (fully) \mathcal{B} -colored graph Γ on Σ with \mathcal{B} -boundary datum \mathbf{b} on $\partial\Sigma$ is a partially \mathcal{B} -colored graph $\mathring{\Gamma}$ on Σ together with a coloring of its internal vertices, i.e. a choice of an element in the vector space $H_v^{\mathcal{B}}$ for each internal vertex $v \in V(\mathring{\Gamma})$, such that the canonical embedding $\partial\Sigma \hookrightarrow \Sigma$, when viewed as an outgoing parametrization of the boundary, pulls back the \mathcal{B} -boundary datum \mathbf{b} on $\partial\Sigma$.

Denote by $\mathbf{G}(\Sigma, \mathbf{b})$ the set of all \mathcal{B} -colored graphs on Σ with prescribed \mathcal{B} -boundary datum \mathbf{b} , and by $\mathbb{k}\mathbf{G}(\Sigma, \mathbf{b})$ the \mathbb{k} -vector space freely generated by it. Also recall from Definition 2.19 the notion of the *value* of a \mathcal{B} -colored embedded graph on the standard disk. We define the *string-net space* assigned to the pair (Σ, \mathbf{b}) as follows:

3.2. DEFINITION. Let \mathcal{B} be a strictly pivotal bicategory, Σ a compact oriented surface, and \mathbf{b} a \mathcal{B} -boundary datum on $\partial\Sigma$.

- (i) A null graph on Σ is an element $\sum_i \lambda_i \Gamma_i$ of $\mathbb{k}\mathbf{G}(\Sigma, \mathbf{b})$ such that there exists an embedding $\varphi: D \hookrightarrow \text{int}(\Sigma)$ of the standard disk D to the interior of Σ that satisfies the following requirements: the circle $\varphi(\partial D)$ does not contain any vertex of any of the graphs Γ_i ; any intersection of $\varphi(\partial D)$ and an edge of any of the graphs Γ_i is transversal; on the complement $\Sigma \setminus \varphi(D)$ all graphs Γ_i coincide; and the values of the graphs pulled back by φ sum up to zero, $\sum_i \lambda_i \langle \Gamma_i \cap \varphi(D) \rangle_D = 0$.

(ii) The (bare) string-net space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ is the quotient

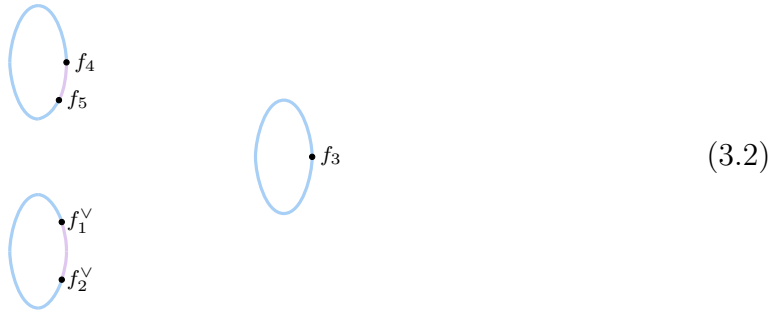
$$\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) := \mathbb{k}\text{G}(\Sigma, \mathbf{b}) / \text{N}(\Sigma, \mathbf{b}), \quad (3.1)$$

where $\text{N}(\Sigma, \mathbf{b})$ is the subspace of $\mathbb{k}\text{G}(\Sigma, \mathbf{b})$ spanned by all null graphs on Σ .

We call a vector in the quotient space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ that is the image of an element of the generating set $\text{G}(\Sigma, \mathbf{b})$ of $\mathbb{k}\text{G}(\Sigma, \mathbf{b})$ a *bare string net*, or also just a *string net*. The qualification “bare” and the notation $\text{SN}_{\mathcal{B}}^{\circ}$ are chosen because later on we will introduce a Karoubified version $\text{SN}_{\mathcal{B}}$ of string-net spaces for which the boundary data are enriched by certain idempotents. A string net that has a \mathcal{B} -colored graph Γ as a representative will be denoted by $[\Gamma]$. By abuse of language the term *string net* is also used when referring to an individual graph that represents an element $[\Gamma] \in \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$. The string-net space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ is linear in the color of each vertex of a graph Γ and additive with respect to taking direct sums of objects labeling an edge of Γ . By the nature of the graphical calculus for disks, isotopic graphs represent the same string net. Furthermore, all identities that are valid in the graphical calculus for \mathcal{B} also hold inside any disk embedded in Σ . In other words, for string nets the graphical calculus for \mathcal{B} applies locally on Σ .

Homeomorphisms of the surface Σ act naturally on embedded graphs. Isotopies can be localized on disks [EK, Cor. 1.3]; as a consequence, isotopic graphs are identified in the string-net space, and hence the action of homeomorphisms descends to an action of the mapping class group $\text{Map}(\Sigma)$ of the surface on the space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$. Moreover, string nets with matching boundary data can be *concatenated*.

3.3. EXAMPLE. Let $a, b, c \in \mathcal{B}$ be three objects in a pivotal bicategory (in the color version of the picture 3.3 below, they are indicated as green, blue, and purple, respectively), and let $f_1, f_5 \in \mathcal{B}(b, c)$, $f_2, f_4 \in \mathcal{B}(c, b)$, $f_3 \in \mathcal{B}(b, b)$ and $f_6 \in \mathcal{B}(a, c)$ be 1-morphisms in \mathcal{B} . Indicate by



a \mathcal{B} -boundary datum \mathbf{b} on the boundary of a genus-1 surface Σ with three boundary circles. Then

$$\Gamma = \text{[Diagram of a string net on a genus-1 surface with three boundary circles. The diagram shows a blue surface with three boundary circles (top, bottom, and right). A pink region is bounded by blue lines. The blue lines are labeled with functions: $f_1, f_2, f_3, f_4, f_5, f_6$. Two vertices are labeled c_1 and c_2 . A small green circle is also present near c_1 .]} \quad (3.3)$$

is an element of the set $\mathbf{G}(\Sigma, \mathbf{b})$ of \mathcal{B} -colored graphs on Σ , where c_1 and c_2 are elements of the appropriate spaces of vertex colors. Γ represents a string net $[\Gamma] \in \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$.

3.4. REMARK. Recall that a color $c \in H_v^{\mathcal{B}}$ for an internal vertex v of a partially \mathcal{B} -colored graph $\mathring{\Gamma}$ is completely determined by a choice of 2-morphism $c_k \in \widehat{h}_v^{\mathcal{B}}(e_k)$ for any choice of polarization on v , and the 2-morphisms for different choices of polarization are related by coherent isomorphisms. Since these coherent isomorphisms are devised in such a way that the string diagrams produced according to different choices of polarizations have the same value when restricted to embedded disks, it is equally admissible, if not more convenient for calculations, to represent string nets by string diagrams with rectangular coupons.

3.5. EXAMPLE. Let \mathbf{b} be a \mathcal{B} -boundary datum on $S^1 = \partial D$. Recall that according to the prescription (2.12) there is a unique partially colored corolla $K^{\mathbf{b}}$ associated to \mathbf{b} . We have a canonical isomorphism

$$\text{SN}_{\mathcal{B}}^{\circ}(D, \mathbf{b}) \xrightarrow{\cong} \mathbf{GCal}_{\mathcal{B}}(K^{\mathbf{b}}) \quad (3.4)$$

that is given by the evaluation $\Gamma \mapsto c_{\Gamma} \xrightarrow{\widehat{\mathbf{GCal}}_{\mathcal{B}}(\mathring{\Gamma})} \langle \Gamma \rangle$, where Γ is any graph representing the string net $[\Gamma] \in \text{SN}_{\mathcal{B}}^{\circ}(D, \mathbf{b})$, and where $c_{\Gamma} \in \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}}$ stands for the coloring of the internal vertex of the underlying partially colored graph $\mathring{\Gamma}$ that is given by Γ . The linear map obtained by this prescription is well defined because the graphical calculus is *local* in nature; it is injective since graphs that evaluate on D to the same value also must have the same value when evaluated on a slightly smaller disk embedded in D (after being replaced by isotopic graphs when necessary) and hence represent the same string net; and it is also surjective because, given any element $c \in \mathbf{GCal}_{\mathcal{B}}(K^{\mathbf{b}})$, coloring the center of $K^{\mathbf{b}}$ with c produces a fully colored corolla $K_c^{\mathbf{b}}$ whose value is c itself.

3.6. REMARK.

1. The idea to utilize pivotal bicategories for the construction of modular functors, and even of topological field theories, dates back at least to [MoW].
2. When the input bicategory \mathcal{B} is the delooping of a spherical fusion category \mathcal{C} , Definition 3.2 reduces to the definition of string-net spaces for \mathcal{C} that was used in

e.g. [LeW, Ki, FuSY, Bar]. We then write

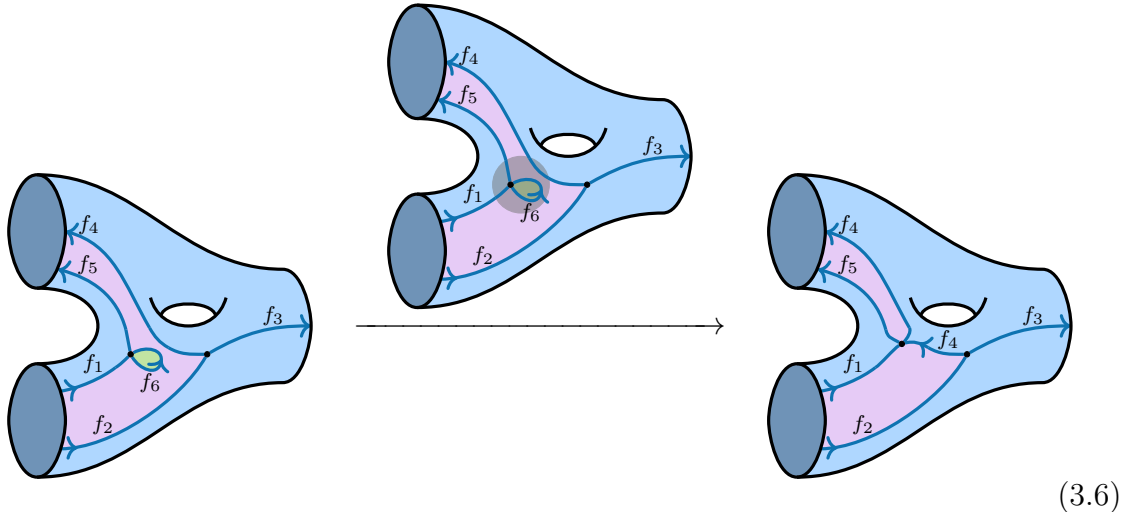
$$\mathrm{SN}_{\mathcal{C}}^{\circ}(\Sigma, \mathbf{b}) \equiv \mathrm{SN}_{\mathcal{BC}}^{\circ}(\Sigma, \mathbf{b}). \quad (3.5)$$

It should be appreciated that our definition of a string-net space $\mathrm{SN}_{\mathcal{B}}^{\circ}$ for a general strictly pivotal bicategory \mathcal{B} does not require any homological properties such as semisimplicity or finiteness. Also note that, while sphericity is needed for the string-net modular functor based on \mathcal{C} to be isomorphic to the Turaev-Viro modular functor, it is neither required for the string-net construction itself [Ru], nor for other constructions of modular functors [BroW, Sect. 6.4].

3.7. STRING-NET SPACES AS COLIMITS. As we will see now, the string-net spaces, which were defined as quotients of vector spaces, admit a natural description as colimits. Such a description will e.g. be instrumental when trying to generalize string-net constructions to a derived setting.

3.8. DEFINITION. *Let \mathcal{B} be a strictly pivotal bicategory, Σ a compact oriented surface, and \mathbf{b} a \mathcal{B} -boundary datum on $\partial\Sigma$.*

- (i) $\mathcal{G}\mathrm{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$ is the following category: The objects of $\mathcal{G}\mathrm{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$ are partially \mathcal{B} -colored graphs on Σ that have \mathbf{b} as their boundary datum. The morphisms of $\mathcal{G}\mathrm{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$ are generated under composition (and upon adjoining identities) by the morphisms exemplified in the following example:



In more detail, a generating morphism is given by an embedding φ of the standard disk D into the interior of Σ , such that $\varphi(\partial D)$ intersects the edges of the partially colored graph $\mathring{\Gamma}_1$ in the domain transversally and does not meet any vertices, while the codomain is the graph $\mathring{\Gamma}_2$ that is obtained by replacing $\mathring{\Gamma}_1 \cap \varphi(D)$ with the image of the unique partially colored corolla on D associated with the boundary datum on $S^1 = \partial D$ pulled back by φ .

(ii) The evaluation functor

$$\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}} : \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b}) \longrightarrow \text{Vect}_{\mathbb{k}} \quad (3.7)$$

is the following functor: $\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}$ sends an object, i.e. a partially \mathcal{B} -colored graph $\mathring{\Gamma}$ on Σ , to the vector space $\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) := \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}}$; and it sends a generating morphism $\gamma : \mathring{\Gamma}_1 \rightarrow \mathring{\Gamma}_2$, with associated embedding $\varphi : D \hookrightarrow \Sigma$, to the linear map

$$\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\gamma) : \bigotimes_{v \in V(\mathring{\Gamma}_1)} H_v^{\mathcal{B}} \xrightarrow{\text{id} \otimes_{\mathbb{k}} \widehat{\text{GCal}}_{\mathcal{B}}(\mathring{\Gamma}_1 \cap D_{\gamma})} \bigotimes_{v' \in V(\mathring{\Gamma}_2)} H_{v'}^{\mathcal{B}} \quad (3.8)$$

that is obtained by applying the graphical calculus on disks to the partially colored graph on D pulled back by the embedding φ .

3.9. THEOREM. Let \mathcal{B} be a strictly pivotal bicategory, Σ a compact oriented surface, and \mathbf{b} a \mathcal{B} -boundary datum on $\partial\Sigma$. The string-net space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ satisfies

$$\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) = \text{colim} \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}, \quad (3.9)$$

where the legs of the cocone are given by

$$\begin{aligned} \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) &= \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}} \longrightarrow \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}), \\ \mathbf{c} &= \bigotimes_{v \in V(\mathring{\Gamma})} \mathbf{c}_v \longmapsto [\mathring{\Gamma}_{\mathbf{c}}] \end{aligned} \quad (3.10)$$

for every $\mathring{\Gamma} \in \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$. Here $\mathring{\Gamma}_{\mathbf{c}}$ is the fully colored graph that is obtained by coloring $\mathring{\Gamma}$ with $\mathbf{c} \in \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}}$.

PROOF. Since graphs related by the local graphical calculus of \mathcal{B} represent the same string net, the prescription (3.10) indeed gives rise to a cocone. To show that the cocone is initial, consider an arbitrary cocone

$$\{f_{\mathring{\Gamma}} : \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) \rightarrow V\}_{\mathring{\Gamma} \in \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})} \quad (3.11)$$

to some \mathbb{k} -vector space V . We need to show that there is a unique linear map $f : \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) \rightarrow V$ that makes the diagram

$$\begin{array}{ccc} \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) & & \\ \downarrow & \searrow f_{\mathring{\Gamma}} & \\ \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) & \xrightarrow{\quad f \quad} & V \end{array} \quad (3.12)$$

commute for every object $\mathring{\Gamma} \in \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$. We claim that the desired linear map is given by

$$\Gamma \mapsto c_{\Gamma} \mapsto f_{\mathring{\Gamma}}(c_{\Gamma}), \quad (3.13)$$

where Γ is any graph representing $[\Gamma] \in \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ and $c_{\Gamma} \in \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma})$. To show that this map is well-defined, assume that the two colored graphs Γ and Γ' both represent the string-net $[\Gamma]$. By the definition of the string-net space, the underlying partially colored graphs $\mathring{\Gamma}$ and $\mathring{\Gamma}'$ are connected by a *zigzag*

$$\begin{array}{ccc} & & \dots \longrightarrow \mathring{\Gamma}' \\ & \vdots & \\ & \downarrow & \\ \mathring{\Gamma}^{(1)} & \longrightarrow & \mathring{\Gamma}^{(2)} \\ \downarrow & & \\ \mathring{\Gamma} & & \end{array} \quad (3.14)$$

in the category $\mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$, and the corresponding vertex colors c_{Γ} and $c_{\Gamma'}$ are related by the zigzag in $\text{Vect}_{\mathbb{k}}$ that is obtained by applying the functor $\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}$ to (3.14), and are therefore mapped to the same element in V by the legs of the cocone (3.11). By construction, f makes the relevant diagrams commute and is unique. ■

By recognizing the string-net space $\text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b})$ as a colimit we shed new light on the canonical mapping class group action: Let $\xi \in \text{Map}(\Sigma)$ be a mapping class group element and x a homeomorphism representing ξ . Then there is a canonical natural isomorphism

$$\begin{array}{ccc} \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b}) & & \\ \downarrow x_* & \searrow \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}} & \\ & \cong & \text{Vect}_{\mathbb{k}} \\ & \nearrow \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}} & \\ \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b}) & & \end{array} \quad (3.15)$$

whose component at an object $\mathring{\Gamma} \in \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$ is the canonical identification

$$\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) = \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}} \xrightarrow{\cong} \bigotimes_{v' \in V(x_* \mathring{\Gamma})} H_{v'}^{\mathcal{B}} = \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(x_* \mathring{\Gamma}), \quad (3.16)$$

where x_* is the endofunctor obtained by pushing forward the partially colored graphs.

Now consider the square

$$\begin{array}{ccc}
 \mathcal{E}_B^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) = \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}} & \longrightarrow & \mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b}) \\
 \downarrow \cong & & \downarrow \mathrm{SN}_B^{\circ}(\xi, \mathbf{b}) \\
 \mathcal{E}_B^{\Sigma, \mathbf{b}}(x_* \mathring{\Gamma}) = \bigotimes_{v' \in V(x_* \mathring{\Gamma})} H_{v'}^{\mathcal{B}} & \longrightarrow & \mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b})
 \end{array} \tag{3.17}$$

Since the composite of the left vertical isomorphism and the horizontal morphism in the bottom row is a component of a natural transformation followed by a leg of a cocone, it is the leg of a cocone under $\mathcal{E}_B^{\Sigma, \mathbf{b}}$. As a consequence, there is a unique endomorphism $\mathrm{SN}_B^{\circ}(\xi, \mathbf{b})$ of the string-net space $\mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b})$ which provides the dashed vertical arrow that makes the square (3.17) commute. By a straightforward diagram-chase we see that this endomorphism coincides with the action of $\xi = [x] \in \mathrm{Map}(\Sigma)$ on $\mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b})$.

Note that the vector space $\mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b})$ is, in general, not finite-dimensional. In particular, $\mathrm{SN}_B^{\circ}(D, \mathbf{b})$ is infinite-dimensional if the corresponding 2-hom space is infinite-dimensional.

3.10. FUNCTORIALITY UNDER RIGID PSEUDOFUNCTORS. As we have seen in Section 2.25, rigid pseudofunctors preserve the graphical calculus on disks for strictly pivotal bicategories. Since the string-net spaces are built with the help of this graphical calculus, one should expect that a rigid pseudofunctor between such bicategories induces canonical linear maps between the respective string-net spaces. Indeed we have

3.11. THEOREM. *Let Σ be a compact oriented surface, \mathcal{B} and \mathcal{B}' two strictly pivotal bicategories, $F: \mathcal{B} \rightarrow \mathcal{B}'$ a rigid pseudofunctor, and \mathbf{b} a \mathcal{B} -boundary datum on $\partial\Sigma$. There is a canonical $\mathrm{Map}(\Sigma)$ -intertwiner*

$$\mathrm{SN}_F^{\circ}(\Sigma, \mathbf{b}) : \quad \mathrm{SN}_B^{\circ}(\Sigma, \mathbf{b}) \longrightarrow \mathrm{SN}_{B'}^{\circ}(\Sigma, F_* \mathbf{b}), \tag{3.18}$$

where $F_* \mathbf{b}$ is the \mathcal{B}' -boundary datum on $\partial\Sigma$ obtained by changing the coloring according to the map of objects and the local functors entailed by F . The linear map (3.18) is defined by sending each representing \mathcal{B} -colored graph to the \mathcal{B}' -colored graph obtained by applying the F -conjugation (2.65). Moreover, the collection of such intertwiners corresponding to different surfaces and boundary data is compatible with the concatenation of string nets.

PROOF. Consider the natural transformation

$$\begin{array}{ccc}
 \mathrm{Graphs}_B(\Sigma, \mathbf{b}) & & \\
 \downarrow F_* & \searrow \mathcal{E}_B^{\Sigma, \mathbf{b}} & \\
 & (-)^F & \mathrm{Vect}_{\mathbb{K}} \\
 \mathrm{Graphs}_{B'}(\Sigma, F_* \mathbf{b}) & \nearrow \mathcal{E}_{B'}^{\Sigma, F_* \mathbf{b}} &
 \end{array} \tag{3.19}$$

whose component at $\mathring{\Gamma} \in \mathcal{G}\text{raphs}_{\mathcal{B}}(\Sigma, \mathbf{b})$ is the F -conjugation

$$(-)_{\mathring{\Gamma}}^F : \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) = \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}} \longrightarrow \bigotimes_{v' \in V(F_* \mathring{\Gamma})} H_{v'}^{\mathcal{B}'} = \mathcal{E}_{\mathcal{B}'}^{\Sigma, F_* \mathbf{b}}(\mathring{\Gamma}). \quad (3.20)$$

The naturality square for a generating morphism $\gamma : \mathring{\Gamma}_1 \rightarrow \mathring{\Gamma}_2$ reads

$$\begin{array}{ccc} \bigotimes_{v_1 \in V(\mathring{\Gamma}_1)} H_{v_1}^{\mathcal{B}} & \xrightarrow{\text{id} \otimes_{\mathbf{k}} \widehat{\text{GCal}}_B(\mathring{\Gamma}_1 \cap D_\gamma)} & \bigotimes_{v_2 \in V(\mathring{\Gamma}_2)} H_{v_2}^{\mathcal{B}} \\ \downarrow (-)_{\mathring{\Gamma}_1}^F & & \downarrow (-)_{\mathring{\Gamma}_2}^F \\ \bigotimes_{v'_1 \in V(F_* \mathring{\Gamma}_1)} H_{v'_1}^{\mathcal{B}} & \xrightarrow{\text{id} \otimes_{\mathbf{k}} \widehat{\text{GCal}}_{B'}(F_* \mathring{\Gamma}_1 \cap D_{F_* \gamma})} & \bigotimes_{v'_2 \in V(F_* \mathring{\Gamma}_2)} H_{v'_2}^{\mathcal{B}} \end{array} \quad (3.21)$$

Commutativity of this diagram follows from Corollary 2.33. Now consider the square

$$\begin{array}{ccc} \mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}(\mathring{\Gamma}) = \bigotimes_{v \in V(\mathring{\Gamma})} H_v^{\mathcal{B}} & \xrightarrow{(3.10)} & \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) \\ \downarrow (-)_{\mathring{\Gamma}}^F & & \downarrow \text{SN}_F^{\circ}(\Sigma, \mathbf{b}) \\ \mathcal{E}_{\mathcal{B}'}^{\Sigma, F_* \mathbf{b}}(F_* \mathring{\Gamma}) = \bigotimes_{v' \in V(F_* \mathring{\Gamma})} H_{v'}^{\mathcal{B}'} & \xrightarrow{(3.10)} & \text{SN}_{\mathcal{B}'}^{\circ}(\Sigma, F_* \mathbf{b}) \end{array} \quad (3.22)$$

Since the composite of the left vertical arrow and the horizontal arrow in the bottom row is a component of a natural transformation followed by a leg of a cocone, it is the leg of a cocone under $\mathcal{E}_{\mathcal{B}}^{\Sigma, \mathbf{b}}$. Therefore there is a unique linear map

$$\text{SN}_F^{\circ}(\Sigma, \mathbf{b}) : \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{b}) \longrightarrow \text{SN}_{\mathcal{B}'}^{\circ}(\Sigma, F_* \mathbf{b}) \quad (3.23)$$

that makes the square commute. A direct diagram-chase shows that this linear map has the asserted form. Equivariance and compatibility with concatenation are evident. \blacksquare

3.12. REMARK. Since a rigid separable Frobenius functor in general preserves horizontal products and whiskerings only up to idempotents of the form (2.85), the change of colors that is brought about by conjugation with respect to a rigid separable Frobenius functor *does not descend* to linear maps between string-net spaces. However, as we will see in Theorem 4.8 as a special case, every rigid separable Frobenius functor does induce linear maps (intertwining the mapping class group actions) between string-net spaces with the help of Frobenius graphs.

3.13. CYLINDER CATEGORIES OVER CIRCLES. The notion of string-net spaces for a strictly pivotal bicategory \mathcal{B} allows us to promote the set of \mathcal{B} -boundary data on a closed oriented 1-manifold to a (small) \mathbf{k} -linear category.

3.14. **DEFINITION.** Let \mathcal{B} be a strictly pivotal bicategory and ℓ a closed oriented 1-manifold. If ℓ is non-empty, the cylinder category $\text{Cyl}^\circ(\mathcal{B}, \ell)$ for \mathcal{B} over ℓ is the following category: An object of $\text{Cyl}^\circ(\mathcal{B}, \ell)$ is a \mathcal{B} -boundary datum on ℓ . A morphism of $\text{Cyl}^\circ(\mathcal{B}, \ell)$ between two boundary data is given by a string net on the cylinder $\ell \times I$ that matches the boundary data at $\ell \times \{0\}$ and $\ell \times \{1\}$. The composition of morphisms is given by the concatenation of string nets.

For the empty 1-manifold \emptyset we set $\text{Cyl}^\circ(\mathcal{B}, \emptyset) := \text{Vect}_k$.

3.15. **EXAMPLE.** For instance, for any choice of α and β ,

$$(3.24)$$

is a morphism in $\text{Cyl}^\circ(\mathcal{B}, S^1)$, with an appropriate \mathcal{B} -coloring. (The boundary component $\ell \times \{0\}$ is regarded as *in-coming* and supports the domain boundary datum, hence the opposite convention for the edge labels.)

3.16. **REMARK.** It follows in particular that for every compact oriented surface Σ the string-net construction provides a functor

$$\text{SN}_{\mathcal{B}}^\circ(\Sigma, -) : \text{Cyl}^\circ(\mathcal{B}, \partial\Sigma) \longrightarrow \text{Vect}_k \quad (3.25)$$

which maps a morphism of the cylinder category to the linear map that is obtained by sewing the cylinder to the boundary.

3.17. **POINTED PIVOTAL BICATEGORIES AND CYLINDER CATEGORIES OVER INTERVALS.** In Definition 3.14 we have introduced cylinder categories over closed oriented 1-manifolds, which uses a pivotal bicategory \mathcal{B} as an input. This is sufficient for obtaining a closed modular functor. In order to obtain instead an open-closed modular functor – a goal that we will achieve in Section 3.34 – we need to define cylinder categories for 1-manifolds with boundary as well. As we will explain in the next section, to this end we must make use of a *pointed* strictly pivotal bicategory $(\mathcal{B}, *_\mathcal{B})$, i.e. a strictly pivotal bicategory endowed with a distinguished object $*_\mathcal{B} \in \mathcal{B}$.

3.18. **DEFINITION.** Let $(\mathcal{B}, *_\mathcal{B})$ be a pointed strictly pivotal bicategory and ℓ a compact oriented 1-manifold with possibly non-empty boundary. The cylinder category $\text{Cyl}^\circ(\mathcal{B}, *_\mathcal{B}, \ell)$ for $(\mathcal{B}, *_\mathcal{B})$ over ℓ is the following category: The objects of $\text{Cyl}^\circ(\mathcal{B}, *_\mathcal{B}, \ell)$ are the $(\mathcal{B}, *_\mathcal{B})$ -boundary data on ℓ , i.e. the \mathcal{B} -boundary data whose 1-cells adjacent to a boundary point of ℓ are all colored with the distinguished object $*_\mathcal{B}$. The morphisms of $\text{Cyl}^\circ(\mathcal{B}, *_\mathcal{B}, \ell)$ are

string nets on the cylinder $\ell \times I$ over ℓ . Their composition is given by concatenating string nets.

We use a specific color (gray, in the color version) to indicate the distinguished object $*_{\mathcal{B}}$. As an illustration,

$$\mathbf{b} = \overset{f}{\bullet} \text{---} \overset{g}{\bullet} \text{---} \overset{h}{\bullet} \quad (3.26)$$

is an object in $\text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I)$. For a *closed* oriented 1-manifold ℓ , we just have $\text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \ell) \equiv \text{Cyl}^\circ(\mathcal{B}, \ell)$.

The following examples of pointed strictly pivotal bicategories are important to us:

3.19. EXAMPLE. The delooping \mathcal{BC} of a strictly pivotal tensor category \mathcal{C} has only a single object and is thus automatically pointed. The restriction on the objects in Definition 3.18 is vacuous in this case. Accordingly we write

$$\text{Cyl}^\circ(\mathcal{C}, \ell) \equiv \text{Cyl}^\circ(\mathcal{BC}, *, \ell) \quad (3.27)$$

and call an object in this bicategory a \mathcal{C} -boundary value on ℓ .

3.20. EXAMPLE. The strictly pivotal bicategory $\mathcal{Fr}(\mathcal{C})$ of simple special symmetric Frobenius algebras in a strictly pivotal tensor category \mathcal{C} is canonically pointed with

$$*_{\mathcal{Fr}(\mathcal{C})} := \mathbb{1} \in \mathcal{Fr}(\mathcal{C}), \quad (3.28)$$

i.e. the tensor unit of \mathcal{C} , canonically viewed as a Frobenius algebra.

3.21. FUNCTORIALITY UNDER EMBEDDINGS. The primary reason for taking as the categorical input for the definition of cylinder categories over compact oriented 1-manifolds with boundary a *pointed* strictly pivotal bicategory is that we want the prescription to be functorial with respect to the embedding of manifolds. This requirement arises, for instance, as a natural property when thinking of topological field theories and their modular functors in the spirit [BruFV] of general covariance in local quantum field theory. For us, functoriality under embeddings is a crucial ingredient for being able to associate a profunctor $\mathcal{SN}_{\mathcal{B}}^\circ(\Sigma; -, \sim)$ to a two-dimensional bordism Σ , which will be done in (3.42) below.

By setting $f \times I: (p, t) \mapsto (f(p), t)$, any continuous map $f: \ell_1 \rightarrow \ell_2$ between 1-manifolds ℓ_1 and ℓ_2 extends canonically to a map $f \times I: \ell_1 \times I \rightarrow \ell_2 \times I$ between the corresponding cylinders. Therefore an orientation preserving automorphism $x: S^1 \rightarrow S^1$ of the circle induces a functor

$$\text{Cyl}^\circ(\mathcal{B}, x): \text{Cyl}^\circ(\mathcal{B}, S^1) \longrightarrow \text{Cyl}^\circ(\mathcal{B}, S^1) \quad (3.29)$$

by pushing forward the objects via x and the morphisms via $x \times I$. In particular, the cylinder category $\text{Cyl}^\circ(\mathcal{B}, S^1)$ carries an action of the circle group $\text{U}(1) \subset \text{Aut}(S^1)$.

As we will see now, in order to achieve functoriality also under embeddings of general oriented 1-manifolds, it does not suffice to take a strictly pivotal bicategory \mathcal{B} as an input, but we also need to fix a suitable distinguished object $*_{\mathcal{B}}$ of \mathcal{B} . To understand this

requirement, consider first an orientation preserving embedding $f: \bigsqcup_{i=1}^n I \hookrightarrow I$ of n copies of the standard interval into itself. Once we impose the restriction on the \mathcal{B} -coloring of the 1-cells of the boundary data as stated in Definition 3.18, we can define a functor

$$\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, f) : \quad \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \bigsqcup_{i=1}^n I) = \prod_{i=1}^n \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I) \longrightarrow \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I) \quad (3.30)$$

by pushing forward the objects and morphisms via f and $f \times I$, respectively and then coloring the complement $I \setminus \mathrm{im} f$ (respectively, $I^2 \setminus \mathrm{im}(f \times I)$) with the object $*_{\mathcal{B}}$. In contrast, in the absence of a distinguished object of \mathcal{B} no consistent coloring of these complements is possible. As a special case, the embedding $\emptyset \hookrightarrow I$ of the empty 1-manifold induces a functor of the type $\mathrm{Cyl}^\circ(\mathcal{B}, \emptyset) = \mathrm{Vect}_{\mathbb{k}} \rightarrow \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I)$ by sending \mathbb{k} to the \mathcal{B} -boundary datum \mathbf{b}_I^* on I that has the entire interval colored with $*_{\mathcal{B}}$. (This works analogously for the embedding $\emptyset \hookrightarrow \ell$ for any oriented 1-manifold ℓ . In this sense, *all* cylinder categories are canonically pointed, by pointing the input bicategory.) Further, by fixing a binary embedding $I \sqcup I \rightarrow I$, the prescription (3.30) endows the category $\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I)$ with a monoidal structure, with the tensor unit given by \mathbf{b}_I^* . Moreover, by sending every object in $\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I)$ to the corresponding horizontal composite of the coloring 1-morphisms, we obtain an equivalence $\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I) \simeq \mathrm{End}_{\mathcal{B}}(*_{\mathcal{B}})$ of tensor categories, with the monoidal product for the endomorphism category given by the horizontal composition. Thus we have arrived at

3.22. PROPOSITION. *Let $(\mathcal{B}, *_{\mathcal{B}})$ be a pointed strictly pivotal bicategory. There is a canonical monoidal equivalence*

$$\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I) \simeq \mathrm{End}_{\mathcal{B}}(*_{\mathcal{B}}). \quad (3.31)$$

Next consider an embedding $f: I \hookrightarrow S^1$. Again by pushing forward the objects (respectively, morphisms) via the embedding (respectively, via $f \times I$) and coloring the complement of the image with the distinguished color, we obtain a functor

$$\mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, f) : \quad \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, I) \longrightarrow \mathrm{Cyl}^\circ(\mathcal{B}, S^1) \equiv \mathrm{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, S^1). \quad (3.32)$$

This is demonstrated in the following picture:

We now promote the assignment of cylinder categories to 1-manifolds to a symmetric monoidal functor between the following categories $\mathrm{Emb}_1^{\mathrm{or}}$ and $\mathrm{Cat}_{\mathbb{k}}$:

3.23. DEFINITION. *The symmetric monoidal category Emb_1^{or} has as objects compact oriented 1-manifolds and as morphisms orientation preserving embeddings; the monoidal product of Emb_1^{or} is disjoint union.*

The symmetric monoidal category $\text{Cat}_{\mathbb{k}}$ has as objects small \mathbb{k} -linear categories and as morphisms \mathbb{k} -linear functors; the monoidal product of $\text{Cat}_{\mathbb{k}}$ is the Cartesian product.

Combining the considerations above we have

3.24. PROPOSITION. *The assignment of cylinder categories for a pointed pivotal bicategory $(\mathcal{B}, *_{\mathcal{B}})$ to compact oriented 1-manifolds canonically extends to a symmetric monoidal functor*

$$\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, -) : \text{Emb}_1^{\text{or}} \longrightarrow \text{Cat}_{\mathbb{k}}. \quad (3.34)$$

3.25. REMARK. As already mentioned, every cylinder category is pointed by the embedding of the empty manifold \emptyset . As a consequence, the functor $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, -)$ factors through the forgetful functor $\text{Cat}_{\mathbb{k}}^{\text{pointed}} \rightarrow \text{Cat}_{\mathbb{k}}$. Even better, we actually obtain a symmetric monoidal 2-functor $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, -) : \mathcal{E}\text{mb}_1^{\text{or}} \rightarrow \text{Cat}_{\mathbb{k}}$, where $\mathcal{E}\text{mb}_1^{\text{or}}$ is the symmetric monoidal (2,1)-category – that is, a symmetric monoidal bicategory with only invertible 2-morphisms – of compact oriented 1-manifolds, orientation preserving embeddings and isotopy classes of isotopies between embeddings.

Finally, let ℓ be a compact oriented 1-manifold and $\bar{\ell}$ the same underlying 1-manifold but with opposite orientation. Due to the strict pivotality of \mathcal{B} we have a canonical isomorphism

$$\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell)^{\text{op}} \xrightarrow{\cong} \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \bar{\ell}) \quad (3.35)$$

which sends an object $\bar{\mathbf{b}} \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell)^{\text{op}}$ to the boundary datum \mathbf{b}^{\vee} on $\bar{\ell}$ that is obtained from \mathbf{b} by taking the duals of the coloring 1-morphisms. Through this identification an orientation reversing embedding $f : \ell_1 \rightarrow \ell_2$ induces a functor

$$\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, f) : \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell_1)^{\text{op}} \longrightarrow \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell_2). \quad (3.36)$$

3.26. REMARK. The string net-construction can be adapted to bicategories \mathcal{B} that do not possess a pivotal structure [KST]: one endows both 1- and 2-manifolds with 2-framings and restricts to string nets that are *progressive* with respect to the 2-framing. Concretely, as before we regard the standard sphere S^1 as a the unit circle in the complex plane \mathbb{C} and the standard interval I as the subset $[0, 1]$ of the real axis. A 2-framing on the cylinder over I and over S^1 , respectively, is then obtained by using as a first vector field a non-zero vector field given by the orientation of I and S^1 , and as a second vector field in the case of I one that uniformly points towards the positive y -direction of \mathbb{C} and in the case of S^1 one that points from each point on S^1 towards the origin $0 \in \mathbb{C}$. Upon fixing a framed embedding $I \hookrightarrow S^1$, we then obtain a functor

$$\bigsqcup_{a \in \mathcal{B}} \mathcal{B}(a, a) \simeq \bigsqcup_{a \in \mathcal{B}} \text{Cyl}^{\circ}(\mathcal{B}, a, I) \xrightarrow{\bigsqcup_{a \in \mathcal{B}} \text{Cyl}^{\circ}(\mathcal{B}, a, f)} \text{Cyl}^{\circ}(\mathcal{B}, S^1). \quad (3.37)$$

This functor admits the structure of a *categorified trace* on \mathcal{B} : A categorified trace on a bicategory \mathcal{B} with values in a category \mathcal{A} is a functor $\llbracket - \rrbracket: \bigsqcup_{a \in \mathcal{B}} \mathcal{B}(a, a) \rightarrow \mathcal{A}$ from the hom-categories of \mathcal{B} to the category \mathcal{A} equipped with a natural isomorphism $\theta: \llbracket fg \rrbracket \xrightarrow{\cong} \llbracket gf \rrbracket$ for any pair $f: a \rightarrow b$ and $g: b \rightarrow a$ of cyclically composable 1-morphisms, satisfying two hexagon and two triangle identities [Po, PS] (compare also [FuSS]). If the natural isomorphism θ is an involution, then the categorified trace is called *symmetric*, or also a *shadow*. An important example of a symmetric categorified trace is provided by the *topological Hochschild homology* (THH) of spectral categories [CaP, Thm. 2.17]. Recently [Ber, HeR], the notion of topological Hochschild homology was extended to that of a bicategory. The THH of a bicategory \mathcal{B} is a category $\mathrm{THH}(\mathcal{B})$ given by a pseudocolimit of a certain 2-truncated cyclic bar construction. As shown in [HeR, Thm. 3.19], $\mathrm{THH}(\mathcal{B})$ is canonically endowed with the structure of a *universal shadow* on \mathcal{B} , in the sense that for every category \mathcal{D} there is a canonical equivalence

$$\mathrm{Fun}(\mathrm{THH}(\mathcal{B}), \mathcal{D}) \longrightarrow \mathrm{Sha}(\mathcal{B}, \mathcal{D}) \quad (3.38)$$

of categories, where $\mathrm{Sha}(\mathcal{B}, \mathcal{D})$ is the category of shadows on \mathcal{B} with values in \mathcal{D} . We expect that the cylinder category $\mathrm{Cyl}^\circ(\mathcal{B}, S^1)$ provides a non-symmetric analogue of $\mathrm{THH}(\mathcal{B})$ – the *topological cyclic homology* of the bicategory \mathcal{B} . More specifically, we expect that the non-symmetric categorified traces obtained this way are universal, and that the cylinder categories over the circle are pseudocolimits of a variant of the cyclic bar construction in which the simplicial category Δ is replaced by the cyclic category Δ^C that was introduced in [Con].

3.27. IDEMPOTENT COMPLETION. The *idempotent completion*, or *Karoubi envelope*, of a category \mathcal{A} is the category $\mathrm{Kar}(\mathcal{A})$ whose objects are the idempotents in \mathcal{A} , while a morphism $f \in \mathrm{Kar}(\mathcal{A})(p_1, p_2)$ between the idempotents $p_1 \in \mathrm{End}_{\mathcal{A}}(a_1)$ and $p_2 \in \mathrm{End}_{\mathcal{A}}(a_2)$ is a morphism $f: a_1 \rightarrow a_2$ in \mathcal{A} satisfying $f \circ p_1 = f = p_2 \circ f$. The identity on an idempotent $p \in \mathrm{End}_{\mathcal{A}}(a)$ viewed as an object in $\mathrm{Kar}(\mathcal{A})$ is p itself, while id_a is not in $\mathrm{End}_{\mathrm{Kar}(\mathcal{A})}(p)$ unless $p = \mathrm{id}_a$. The Karoubi envelope $\mathrm{Kar}(\mathcal{A})$ comes with a canonical fully faithful functor $K_{\mathcal{A}}: \mathcal{A} \rightarrow \mathrm{Kar}(\mathcal{A})$ that sends an object $a \in \mathcal{A}$ to the idempotent id_a . The functor $K_{\mathcal{A}}$ is universal among the functors from \mathcal{A} that have a chosen *splitting* for every idempotent in their codomain. A crucial further property of $K_{\mathcal{A}}$ is that it is *cofinal* in the sense² that we can restrict functors from $\mathrm{Kar}(\mathcal{A})$ to functors from \mathcal{A} along $K_{\mathcal{A}}$ without changing their colimits (compare [Lu, Lemma 5.1.4.6] for the $(\infty, 1)$ -version of this result).

In the present subsection we introduce idempotent completions of the cylinder categories, to which we refer as *Karoubified cylinder categories*, and define string-net spaces whose boundary data are objects in these Karoubified cylinder categories. This is in line with the common practice in skein theoretic quantum topology to study various kinds of completions. For instance, one considers the free cocompletions of the skein categories – the higher-dimensional analogues of cylinder categories that are defined for surfaces – to

² We adopt the terminology of [Bo] and [Lu]. This differs from the one in [Ma] and [Jo], where such a functor is instead called a *final* functor.

recover the locally finitely presentable factorization homology [Coo] and provide geometric models for quantum character varieties when applied to quantum groups [BeBJ]. In dimension (2,1), for the case that \mathcal{B} is the delooping \mathcal{BC} of a spherical fusion category \mathcal{C} , the Karoubified cylinder category over the standard circle is canonically equivalent to the Drinfeld center $\mathcal{Z}(\mathcal{C})$ of the fusion category \mathcal{C} [Ki], and the corresponding Karoubified string-net construction extends to a (3,2,1)-dimensional topological field theory that is equivalent to the Turaev-Viro state-sum model for \mathcal{C} ([Bar]; see also [Go] for an approach based on the presentation of $\mathcal{Bord}_{3,2,1}^{\text{or}}$ that was conjectured in [BDSV1].) In the present paper, a major motivation for performing the idempotent completion of cylinder categories is that in the application of string-net models to the construction of correlators of rational conformal field theories, field objects are naturally realized as objects in the Karoubified cylinder categories (for details, see Chapter 3 of [FuSY] and Section 4 below).

3.28. DEFINITION. *Let ℓ be a compact oriented (not necessarily closed) 1-manifold, and let $(\mathcal{B}, *_{\mathcal{B}})$ a pointed strictly pivotal bicategory.*

- (i) *The Karoubified cylinder category for $(\mathcal{B}, *_{\mathcal{B}})$ over ℓ is the Karoubi envelope*

$$\text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \ell) := \text{Kar}(\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell)). \quad (3.39)$$

*Thus an object $\mathbf{B} \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \ell)$ is an idempotent $\mathbf{B}: \mathbf{B}^{\circ} \rightarrow \mathbf{B}^{\circ}$ in the ordinary cylinder category $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \ell)$; we call it a thickened $(\mathcal{B}, *_{\mathcal{B}})$ -boundary datum on ℓ .*

- (ii) *The Karoubified string-net space $\text{SN}_{\mathcal{B}}(\Sigma, \mathbf{B})$ for a compact oriented surface Σ and a thickened $(\mathcal{B}, *_{\mathcal{B}})$ -boundary datum $\mathbf{B} \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \partial\Sigma)$ is the subspace*

$$\text{SN}_{\mathcal{B}}(\Sigma, \mathbf{B}) := \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{B}^{\circ})^{\mathbf{B}} \subset \text{SN}_{\mathcal{B}}^{\circ}(\Sigma, \mathbf{B}^{\circ}) \quad (3.40)$$

consisting of string nets that are invariant under concatenation with the idempotent \mathbf{B} .

Note that, since $\partial\Sigma$ has empty boundary, the cylinder category over $\partial\Sigma$ actually does not depend on the distinguished object, i.e. $\text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \partial\Sigma) = \text{Cyl}(\mathcal{B}, \partial\Sigma)$. Thus with Definition 3.28 we obtain a functor

$$\text{SN}_{\mathcal{B}}(\Sigma, -) : \text{Cyl}(\mathcal{B}, \partial\Sigma) \longrightarrow \text{Vect}_{\mathbb{k}} \quad (3.41)$$

for every compact oriented surface Σ . Note that, just like $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, -)$, also the assignment $\text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, -)$ is functorial with respect to the embeddings of 1-manifolds.

3.29. FACTORIZATION. Recall, e.g. from [FuSY, Def. 2.1], the description of the symmetric monoidal bicategory $\mathcal{Bord}_{2,\text{o/c}}^{\text{or}}$ of *two-dimensional open-closed bordisms*: An object $\alpha \in \mathcal{Bord}_{2,\text{o/c}}^{\text{or}}$ is a finite disjoint union of copies of the standard interval $I = [0, 1] \subset \mathbb{R}$ (oriented from 1 to 0) and the standard circle $S^1 \subset \mathbb{C}$ (oriented counterclockwise). A 1-morphism between objects α and β – referred to as an open-closed bordism, or just *bordism*, for short, and to be denoted by $\Sigma: \alpha \rightarrow \beta$ – consists of an underlying compact oriented surface Σ , an in-going parametrization $\phi_-: \bar{\alpha} \hookrightarrow \partial\Sigma$ and an out-going parametrization $\phi_+: \beta \hookrightarrow \partial\Sigma$.

The functoriality of $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, -)$ and of $\text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, -)$ under embeddings of 1-manifolds that we established in Section 3.21 implies

3.30. LEMMA. *Let $(\mathcal{B}, *_{\mathcal{B}})$ be a pointed strictly pivotal bicategory. To any bordism $\Sigma: \alpha \rightarrow \beta$ there is naturally associated a \mathbb{k} -linear profunctor*

$$\mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; -, \sim) : \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \alpha) \multimap \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta), \quad (3.42)$$

as well as a Karoubified version

$$\mathcal{SN}_{\mathcal{B}}(\Sigma; -, \sim) : \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \alpha) \multimap \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \beta) \quad (3.43)$$

of $\mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; -, \sim)$.

PROOF. The profunctor $\mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; -, \sim)$ is given by the composite

$$\begin{aligned} \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \alpha)^{\text{op}} \times \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta) &\xrightarrow{(\phi_-)_* \sqcup (\phi_+)_*} \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \partial\Sigma) \\ &\xrightarrow{\mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma, -)} \text{Vect}_{\mathbb{k}} \end{aligned} \quad (3.44)$$

of the functor $\mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma, -)$ with the in- and out-going parametrizations of $\partial\Sigma$ (here we abbreviate $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \phi)$ by ϕ_*). The Karoubified version is defined analogously. ■

We are now ready to state the following result, which appears to be a variant of a folk theorem (see e.g. [Wa, GuJS, KiT]):

3.31. THEOREM. *Let $\Sigma: \alpha \sqcup \beta \rightarrow \beta \sqcup \gamma$ be a bordism, with $\alpha, \beta, \gamma \in \mathcal{Bord}_{2,o/c}^{\text{or}}$. Then the family*

$$\{s_{-, \mathbf{b}_0, \sim}^{\Sigma} : \mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; -, \mathbf{b}_0, \mathbf{b}_0, \sim) \Longrightarrow \mathcal{SN}_{\mathcal{B}}^{\circ}(\cup_{\beta}\Sigma; -, \sim)\}_{\mathbf{b}_0 \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \quad (3.45)$$

of natural transformations whose members are given by the sewing of string nets, is dinatural and exhibits the functor $\mathcal{SN}_{\mathcal{B}}^{\circ}(\cup_{\beta}\Sigma; -, \sim)$ as the coend

$$\int^{\mathbf{b} \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; -, \mathbf{b}, \mathbf{b}, \sim) : \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \alpha) \multimap \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \gamma). \quad (3.46)$$

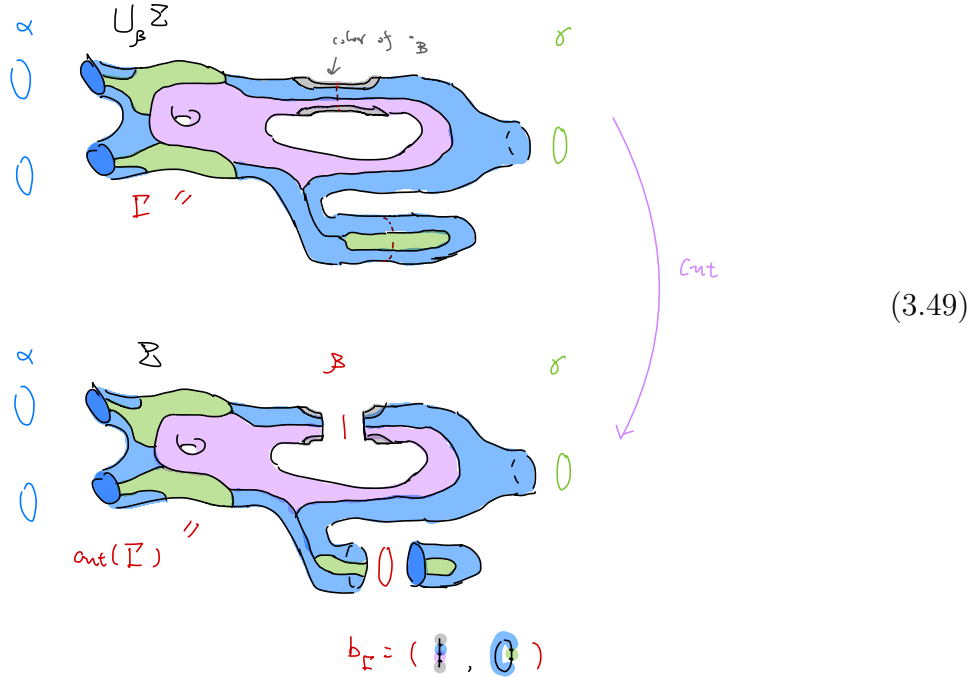
PROOF. We need to show that the family 3.45 is dinatural, and that it is universal among dinatural families of the same type. Dinaturality holds by the fact that sewing a morphism F in the cylinder category $\text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)$ to the first slot (along the appropriate 1-manifold) of the profunctor and then sewing along β yields the same result as sewing F to the the second slot and then doing the final sewing. To show universality, consider a dinatural family

$$\{g_{\mathbf{b}_0} : \mathcal{SN}_{\mathcal{B}}^{\circ}(\Sigma; \mathbf{a}_0, \mathbf{b}_0, \mathbf{b}_0, \mathbf{c}_0) \longrightarrow V\}_{\mathbf{b}_0 \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \quad (3.47)$$

for arbitrary boundary data $\mathbf{a}_0 \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \alpha)$ and $\mathbf{c}_0 \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \gamma)$ and an arbitrary \mathbb{k} -vector space $V \in \text{Vect}_{\mathbb{k}}$. Define the linear map

$$\begin{aligned} g : \mathcal{SN}_{\mathcal{B}}^{\circ}(\cup_{\beta}\Sigma; \mathbf{a}_0, \mathbf{c}_0) &\longrightarrow V, \\ [\Gamma] &\longmapsto [\text{cut}(\Gamma)] \longmapsto g_{\mathbf{b}_r}([\text{cut}(\Gamma)]), \end{aligned} \quad (3.48)$$

where Γ is a representative graph that intersects β only at edges and any such intersection is transversal (such a representative can be chosen without loss of generality), $\text{cut}(\Gamma)$ is the fully colored graph on Σ obtained by cutting the representative graph Γ along β , and $\mathbf{b}_\Gamma \in \text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \beta)$ is the boundary datum arising from the cut. The scenario is illustrated by the following schematic example:



The linear map (3.48) is well-defined because the family $\{g_{\mathbf{b}_0}\}_{\mathbf{b}_0 \in \text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \beta)}$ is dinatural by assumption and because all relevant isotopies as well as the generating local relations provided by the graphical calculus on disks are contained within embedded disks, which implies that a different choice of representative for the string-net $[\Gamma] \in \mathcal{SN}_\mathcal{B}^\circ(\cup_\beta \Sigma; \mathbf{a}_0, \mathbf{c}_0)$ only differs by the action of some morphism in the cylinder category $\text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \beta)$.

Uniqueness of the map (3.48) is guaranteed by design. Hence the dinatural family at each component does exhibit a coend at the level of vector spaces and linear maps. By a standard argument the so defined component-wise coends automatically combine to a coend at the level of linear functors and linear natural transformations. ■

An analogous statement holds for the Karoubified string-net functors. This is achieved by the following observation. Recall that the *twisted arrow category* $\text{Tw}(\mathcal{A})$ associated with a category \mathcal{A} is the category whose objects are the morphisms of \mathcal{A} and for which a morphism from $f: a \rightarrow b$ to $f': a' \rightarrow b'$ is a pair $(g, h) \in \mathcal{A}^{\text{op}}(a, a') \times \mathcal{A}(b, b')$ such that the square

$$\begin{array}{ccc} a & \xrightarrow{f} & b \\ g \uparrow & & \downarrow h \\ a' & \xrightarrow{f'} & b' \end{array} \quad (3.50)$$

commutes. The category $\mathrm{Tw}(\mathcal{A})$ comes with a canonical projection functor

$$\pi_{\mathcal{A}} : \mathrm{Tw}(\mathcal{A}) \rightarrow \mathcal{A}^{\mathrm{op}} \times \mathcal{A}, \quad (3.51)$$

which sends an object $f : a \rightarrow b$ to $(a, b) \in \mathcal{A}^{\mathrm{op}} \times \mathcal{A}$ and keeps the morphisms as they are. The relevance of this construction to us is that the coend of a functor $F : \mathcal{A}^{\mathrm{op}} \times \mathcal{A} \rightarrow \mathrm{Vect}_{\mathbb{k}}$ can be identified with a *colimit* over the twisted arrow category, according to

$$\int^{a \in \mathcal{A}} F(a, a) = \mathrm{colim}(\mathrm{Tw}(\mathcal{A}^{\mathrm{op}})^{\mathrm{op}} \xrightarrow{\pi_{\mathcal{A}^{\mathrm{op}}}^{\mathrm{op}}} \mathcal{A}^{\mathrm{op}} \times \mathcal{A} \xrightarrow{F} \mathrm{Vect}_{\mathbb{k}}), \quad (3.52)$$

see e.g. [Lo, Sect. 1.2].

3.32. LEMMA. *For \mathcal{A} a category, let $F : \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \rightarrow \mathrm{Vect}_{\mathbb{k}}$ be a functor whose coend exists. Then*

$$\int^{a \in \mathcal{A}} F|_{\mathcal{A}}(a, a) = \int^{A \in \mathrm{Kar}(\mathcal{A})} F(A, A), \quad (3.53)$$

where $F|_{\mathcal{A}} : \mathcal{A}^{\mathrm{op}} \times \mathcal{A} \rightarrow \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \xrightarrow{F} \mathrm{Vect}_{\mathbb{k}}$ is the restriction of F along the canonical embedding.

PROOF. Consider the functor

$$\begin{aligned} G : \mathrm{Tw}(\mathcal{A}^{\mathrm{op}})^{\mathrm{op}} &\longrightarrow \mathrm{Tw}(\mathrm{Kar}(\mathcal{A})^{\mathrm{op}})^{\mathrm{op}}, \\ (a \xleftarrow{f} b) &\longmapsto (\mathrm{id}_a \xleftarrow{f} \mathrm{id}_b). \end{aligned} \quad (3.54)$$

Owing to the commutativity of the square

$$\begin{array}{ccc} \mathrm{Tw}(\mathcal{A}^{\mathrm{op}})^{\mathrm{op}} & \xrightarrow{G} & \mathrm{Tw}(\mathrm{Kar}(\mathcal{A})^{\mathrm{op}})^{\mathrm{op}} \\ \pi_{\mathcal{A}^{\mathrm{op}}}^{\mathrm{op}} \downarrow & & \downarrow \pi_{\mathrm{Kar}(\mathcal{A})^{\mathrm{op}}}^{\mathrm{op}} \\ \mathcal{A}^{\mathrm{op}} \times \mathcal{A} & \longrightarrow & \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \end{array} \quad (3.55)$$

it suffices to show that G is cofinal. Indeed, if that is the case, we have

$$\begin{aligned} \int^{a \in \mathcal{A}} F|_{\mathcal{A}}(a, a) &= \mathrm{colim}(\mathrm{Tw}(\mathcal{A}^{\mathrm{op}})^{\mathrm{op}} \xrightarrow{\pi_{\mathcal{A}^{\mathrm{op}}}^{\mathrm{op}}} \mathcal{A}^{\mathrm{op}} \times \mathcal{A} \longrightarrow \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \xrightarrow{F} \mathrm{Vect}_{\mathbb{k}}) \\ &= \mathrm{colim}(\mathrm{Tw}(\mathcal{A}^{\mathrm{op}})^{\mathrm{op}} \xrightarrow{G} \mathrm{Tw}(\mathrm{Kar}(\mathcal{A})^{\mathrm{op}})^{\mathrm{op}} \\ &\quad \xrightarrow{\pi_{\mathrm{Kar}(\mathcal{A})^{\mathrm{op}}}^{\mathrm{op}}} \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \xrightarrow{F} \mathrm{Vect}_{\mathbb{k}}) \\ &= \mathrm{colim}(\mathrm{Tw}(\mathrm{Kar}(\mathcal{A})^{\mathrm{op}})^{\mathrm{op}} \xrightarrow{\pi_{\mathrm{Kar}(\mathcal{A})^{\mathrm{op}}}^{\mathrm{op}}} \mathrm{Kar}(\mathcal{A})^{\mathrm{op}} \times \mathrm{Kar}(\mathcal{A}) \xrightarrow{F} \mathrm{Vect}_{\mathbb{k}}) \\ &= \int^{A \in \mathrm{Kar}(\mathcal{A})} F(A, A), \end{aligned} \quad (3.56)$$

where cofinality of G is used in the third equality.

To show that the functor $G: \text{Tw}(\mathcal{A}^{\text{op}})^{\text{op}} \rightarrow \text{Tw}(\text{Kar}(\mathcal{A})^{\text{op}})^{\text{op}}$ is cofinal, we use the fact that this is the case iff for every object $(p \xrightarrow{g} q) \in \text{Tw}(\text{Kar}(\mathcal{A})^{\text{op}})^{\text{op}}$, with $p \in \text{End}_{\mathcal{A}}(a)$ and $q \in \text{End}_{\mathcal{A}}(b)$ idempotents, the comma category $g \downarrow G$ is connected, i.e. it is non-empty and for each pair of objects in $g \downarrow G$ there exists a zigzag connecting them. To see that the latter conditions are satisfied, first note that the square

$$\begin{array}{ccc} p & \xleftarrow{g} & q \\ p \uparrow & & \downarrow q \\ \text{id}_a & \xleftarrow{g} & \text{id}_b \end{array} \quad (3.57)$$

commutes, due to the defining condition for g to be a morphism of type $q \rightarrow p$. This commutative square provides us with an object $(g \rightarrow Gg) \in g \downarrow G$, hence $g \downarrow G$ is non-empty. Next assume that we have a pair of objects $g \rightarrow Gf$ and $g \rightarrow Gh$ in the comma category $g \downarrow G$, given by the two commutative squares

$$\begin{array}{ccc} p & \xleftarrow{g} & q \\ r \uparrow & & \downarrow s \\ \text{id}_c & \xleftarrow{f} & \text{id}_d \end{array} \quad \text{and} \quad \begin{array}{ccc} p & \xleftarrow{g} & q \\ r' \uparrow & & \downarrow s' \\ \text{id}_{c'} & \xleftarrow{h} & \text{id}_{d'} \end{array} \quad (3.58)$$

Observe that the diagram

$$\begin{array}{ccccc} & \text{id}_c & \xleftarrow{f} & \text{id}_d & \\ & \swarrow r & & \nwarrow s & \\ & & \text{id}_a & \xleftarrow{g} & \text{id}_b \\ p & \xleftarrow{g} & q & \xleftarrow{q} & \\ & \nwarrow r' & & \swarrow s' & \\ & \text{id}_{c'} & \xleftarrow{h} & \text{id}_{d'} & \end{array} \quad (3.59)$$

commutes and hence the two squares that do not involve the objects p and q (and in the color version are drawn in violet) give rise to a span

$$(g \rightarrow Gf) \longleftarrow (g \rightarrow Gg) \longrightarrow (g \rightarrow Gh) \quad (3.60)$$

in the comma category $g \downarrow G$. It follows that for every $g \in \text{Tw}(\text{Kar}(\mathcal{A})^{\text{op}})^{\text{op}}$ the comma category $g \downarrow G$ is connected. Hence the functor $G: \text{Tw}(\mathcal{A}^{\text{op}})^{\text{op}} \rightarrow \text{Tw}(\text{Kar}(\mathcal{A})^{\text{op}})^{\text{op}}$ is cofinal, as claimed. \blacksquare

Combining Theorem 3.31 with Lemma 3.32 we arrive at

3.33. COROLLARY. *Let $\Sigma: \alpha \sqcup \beta \rightarrow \beta \sqcup \gamma$ be a bordism, with $\alpha, \beta, \gamma \in \mathcal{B}\text{ord}_{2,0/c}^{\text{or}}$. Then the family*

$$\{\widehat{s}_{-,b_0,\sim}^\Sigma: \mathcal{SN}_{\mathcal{B}}(\Sigma; -, B_0, B_0, \sim) \Rightarrow \mathcal{SN}_{\mathcal{B}}(\cup_\beta \Sigma; -, \sim)\}_{B_0 \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \quad (3.61)$$

whose members are given by the sewing of string nets, is dinatural and exhibits the functor $\mathcal{SN}_{\mathcal{B}}(\cup_\beta \Sigma; -, \sim)$ as the coend

$$\int^{\mathcal{B} \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{SN}_{\mathcal{B}}(\Sigma; -, B, B, \sim) : \quad \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \alpha) \rightarrow \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \gamma). \quad (3.62)$$

PROOF. By direct calculation we have

$$\begin{aligned} \mathcal{SN}_{\mathcal{B}}(\cup_\beta \Sigma; , A_0, C_0) &= \mathcal{SN}_{\mathcal{B}}^\circ(\cup_\beta \Sigma; , A_0^\circ, C_0^\circ)^{(A_0, C_0)} \\ &= \int^{\mathcal{B} \in \text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{SN}_{\mathcal{B}}^\circ(\Sigma; A_0^\circ, b, b, C_0^\circ)^{(A_0, C_0)} \\ &= \int^{\mathcal{B} \in \text{Cyl}^\circ(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{SN}_{\mathcal{B}}(\Sigma; A_0, \text{id}_b, \text{id}_b, C_0) \\ &\stackrel{(3.53)}{=} \int^{\mathcal{B} \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{SN}_{\mathcal{B}}(\Sigma; A_0, B, B, C_0) \end{aligned} \quad (3.63)$$

for every $A_0 \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \alpha)$ and $C_0 \in \text{Cyl}(\mathcal{B}, *_{\mathcal{B}}, \gamma)$. ■

3.34. OPEN-CLOSED MODULAR FUNCTORS FROM BICATEGORICAL STRING NETS. We are now going to show that bicategorical string nets provide us with modular functors. With the application to conformal field theory in mind (see Section 4, and in particular Remark 4.11), we are interested in *open-closed* modular functors, for which the 1-manifolds that are objects of the domain bicategory may have a non-empty boundary. In Section 3.17 we have already prepared the ground for dealing with this case by defining cylinder categories not only over circles, but also over intervals.

3.35. DEFINITION. *An open-closed modular functor is a symmetric monoidal pseudo-functor*

$$\mathcal{B}\text{ord}_{2,0/c}^{\text{or}} \longrightarrow \mathcal{P}\text{rof}_{\mathbb{k}} \quad (3.64)$$

from the symmetric monoidal bicategory $\mathcal{B}\text{ord}_{2,0/c}^{\text{or}}$ of two-dimensional open-closed bordisms to the symmetric monoidal bicategory of $\mathcal{P}\text{rof}_{\mathbb{k}}$ of \mathbb{k} -linear profunctors.

The objects of the bicategory $\mathcal{B}\text{ord}_{2,0/c}^{\text{or}}$ are finite disjoint unions of copies of the standard circle $S^1 = \{z \mid |z| = 1\} \subset \mathbb{C}$ and of the standard interval $I = [0, 1] \subset \mathbb{R} \subset \mathbb{C}$, the 1-morphisms are bordisms between such 1-manifolds, and the 2-morphisms are isotopy classes of diffeomorphisms; for more details on $\mathcal{B}\text{ord}_{2,0/c}^{\text{or}}$ see e.g. Definition 2.1 of [FuSY]. The bicategory $\mathcal{P}\text{rof}_{\mathbb{k}}$ is defined as follows:³

³ The definition used here deviates somewhat from the one in [FuSY].

- Objects of $\mathcal{P}\text{rof}_{\mathbb{k}}$ are small categories enriched in the cocomplete category $\text{Vect}_{\mathbb{k}}$ of (not necessarily finite-dimensional) \mathbb{k} -vector spaces.
- For objects A and B , a 1-morphism $P: A \multimap B$ is a \mathbb{k} -linear profunctor from A to B , that is, a \mathbb{k} -linear functor $P: A^{\text{op}} \times B \rightarrow \text{Vect}_{\mathbb{k}}$.
- For 1-morphisms $P, Q: A \multimap B$, a 2-morphism $\varphi: P \Rightarrow Q$ is a natural transformation of the underlying functors.
- The composite of a composable pair $A \xrightarrow{P} B \xrightarrow{Q} C$ of 1-morphisms is the coend

$$P \cdot Q := \int^{b \in B} P(-, b) \otimes_{\mathbb{k}} Q(b, \sim) : A \multimap C. \quad (3.65)$$

The horizontal composition of 2-morphisms is induced by the composition of 1-morphisms. (The respective coends exist because all the domain categories are small and the target categories are cocomplete, see e.g. [Ri, Prop. 4.5.3].)

- Vertical composition is given by the vertical composition of natural transformations.
- The monoidal structure given by the Cartesian product of \mathbb{k} -linear categories,⁴ with the obvious symmetric braiding.

A specific implication of Theorem 3.31 is that for any pair of composable bordisms $\Sigma: \alpha \multimap \beta$ and $\Sigma': \beta \multimap \gamma$ the dinatural family of sewing maps exhibits the structure of a coend on

$$\begin{aligned} \mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ}(\Sigma \cup_{\beta} \Sigma'; -, \sim) &= \int^{\mathbf{b} \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ}(\Sigma \sqcup \Sigma'; -, \mathbf{b}, \mathbf{b}, \sim) \\ &= \int^{\mathbf{b} \in \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \beta)} \mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ}(\Sigma; -, \mathbf{b}) \otimes_{\mathbb{k}} \mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ}(\Sigma'; \mathbf{b}, \sim). \end{aligned} \quad (3.66)$$

An analogous statement follows from Corollary 3.33 for the Karoubified string-net spaces. This immediately implies

3.36. THEOREM. *Let $(\mathcal{B}, *_{\mathcal{B}})$ be a pointed strictly pivotal bicategory. Then the assignments*

$$\alpha \mapsto \text{Cyl}^{\circ}(\mathcal{B}, *_{\mathcal{B}}, \alpha) \quad \text{and} \quad \Sigma \mapsto \mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ}(\Sigma; -, \sim) \quad (3.67)$$

extend to an open-closed modular functor, i.e. a symmetric monoidal pseudofunctor

$$\mathcal{S}\mathcal{N}_{\mathcal{B}}^{\circ} : \text{Bord}_{2, \text{o/c}}^{\text{or}} \longrightarrow \mathcal{P}\text{rof}_{\mathbb{k}} \quad (3.68)$$

⁴ The appropriate notion of Cartesian product for \mathbb{k} -linear categories is the one provided by the framework of enriched categories, i.e. for small \mathbb{k} -linear categories A and B , $A \times B$ is the small \mathbb{k} -linear category whose objects are ordered pairs and whose morphism spaces are obtained as tensor products over \mathbb{k} , i.e. $(A \times B)(a_1 \times b_1, a_2 \times b_2) := A(a_1, a_2) \otimes_{\mathbb{k}} B(b_1, b_2)$ (see e.g. [Ke, Sect. 1.4], where this product is instead referred to as *tensor product*).

from the symmetric monoidal bicategory of open-closed bordisms to the symmetric monoidal bicategory of \mathbb{k} -linear profunctors. Similarly, the Karoubified cylinder categories and string-net spaces give rise to another open-closed modular functor

$$\mathcal{SN}_{\mathcal{B}} : \mathcal{Bord}_{2,o/c}^{\text{or}} \longrightarrow \mathcal{Prof}_{\mathbb{k}}. \quad (3.69)$$

Now let \mathcal{C} be a modular fusion category that encodes the chiral data for a rational conformal field theory. The results of Section 3.27 then amount to the statement that the open-closed modular functor $\mathcal{SN}_{\mathcal{C}} = \mathcal{SN}_{\mathcal{BC}}$ models the conformal blocks of the conformal field theory. In other words, by setting

$$\text{Bl}_{\mathcal{C}}(\alpha) := \text{Cyl}(\mathcal{C}, \alpha) \quad (3.70)$$

for every $\alpha \in \mathcal{Bord}_{2,o/c}^{\text{or}}$, and

$$\text{Bl}_{\mathcal{C}}(\Sigma; -, \sim) := \mathcal{SN}_{\mathcal{C}}(\Sigma; -, \sim) \quad (3.71)$$

for every bordism Σ , we obtain an open-closed modular functor that when restricted to the closed sector is isomorphic to the modular functor that is provided by the Turaev-Viro state sum construction and that furnishes canonical equivalences $\text{Bl}_{\mathcal{C}}(I) \simeq \mathcal{C}$ and $\text{Bl}_{\mathcal{C}}(S^1) \simeq \mathcal{Z}(\mathcal{C})$. (For more details, see Chapter 3.3 of [FuSY].)

3.37. REMARK. For \mathcal{C} a modular fusion category, the categories $\text{Bl}_{\mathcal{C}}(I) = \text{Cyl}(\mathcal{C}, I) \simeq \mathcal{C}$ and $\text{Bl}_{\mathcal{C}}(S^1) = \text{Cyl}(\mathcal{C}, S^1) \simeq \mathcal{Z}(\mathcal{C})$ are finite and semisimple. As a consequence, for any bordism $\Sigma: \alpha \rightarrow \beta$ the functor $\text{Bl}_{\mathcal{C}}(\Sigma; -, \sim) = \mathcal{SN}_{\mathcal{C}}(\Sigma; -, \sim)$ is exact in each of its variables. It can therefore be substituted by an exact functor $\widehat{\text{Bl}}_{\mathcal{C}}(\Sigma; - \boxtimes \sim): \text{Bl}_{\mathcal{C}}(\alpha)^{\text{op}} \boxtimes \text{Bl}_{\mathcal{C}}(\beta) \rightarrow \text{Vect}_{\mathbb{k}}$. Thus the open-closed modular functor $\text{Bl}_{\mathcal{C}}$ factors through the forgetful functor $\mathcal{Prof}_{\mathbb{k}}^{\mathcal{L}^{\text{ex}}} \rightarrow \mathcal{Prof}_{\mathbb{k}}$, where $\mathcal{Prof}_{\mathbb{k}}^{\mathcal{L}^{\text{ex}}}$ is the symmetric monoidal bicategory of finite \mathbb{k} -linear abelian categories, left exact profunctors and natural transformations, the horizontal composition of which is given by left exact coends (in the sense of [FuS, Sect. 3.2]), and with the monoidal product being the Deligne product. We thus also have a modular functor in the sense of Definition 2.1 of [FuSY], in which $\mathcal{Prof}_{\mathbb{k}}^{\mathcal{L}^{\text{ex}}}$ is taken as the target category. (The motivation for making this choice in [FuSY] is that it suits a potential extension to non-semisimple modular tensor categories. In the case of a modular fusion category \mathcal{C} considered here, the profunctors are in fact both left and right exact.)

We finally mention that not only $\mathcal{SN}_{\mathcal{C}}$, but also the open-closed modular functor

$$\mathcal{SN}_{\mathcal{Fr}(\mathcal{C})}^{\circ} : \mathcal{Bord}_{2,o/c}^{\text{or}} \longrightarrow \mathcal{Prof}_{\mathbb{k}} \quad (3.72)$$

that is obtained by taking $(\mathcal{Fr}(\mathcal{C}), \mathbb{1})$ as the decorating pointed pivotal bicategory is relevant to rational conformal field theory. This will be explained in Section 4.

4. Application: Universal correlators in RCFT

The string-net construction of correlators for a two-dimensional rational conformal field theory (RCFT) that was developed in [FuSY] provides a natural motivation for the bicategorical string-net construction that we presented in Section 3: Given a modular fusion category \mathcal{C} describing the chiral data for the RCFT, a topological world sheet \mathcal{S} with physical boundaries and topological defect lines – a *world sheet*, for short – can be regarded as a $\mathcal{F}r(\mathcal{C})$ -colored string diagram $\tilde{\mathcal{S}}$ on its underlying surface $\Sigma_{\mathcal{S}}$. In [FuSY], we assigned to each world sheet \mathcal{S} an element $\text{Cor}_{\mathcal{C}}(\mathcal{S}) \in \text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}}))$ in a (Karoubified) string-net space colored by the modular fusion category \mathcal{C} , where $\mathbf{b}_{\mathcal{S}}$ is the boundary datum of the $\mathcal{F}r(\mathcal{C})$ -colored string diagram $\tilde{\mathcal{S}}$, while $\mathbb{F}_{\partial\Sigma_{\mathcal{S}}}: \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, \partial\Sigma_{\mathcal{S}}) \rightarrow \text{Cyl}(\mathcal{C}, \partial\Sigma_{\mathcal{S}})$ is a canonical functor sending an $\mathcal{F}r(\mathcal{C})$ -colored boundary datum to an idempotent in the \mathcal{C} -colored cylinder category over the boundary $\partial\Sigma_{\mathcal{S}}$ of $\Sigma_{\mathcal{S}}$. The \mathcal{C} -colored string-net space $\text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}}))$ can be regarded as the *space of conformal blocks* for the world sheet \mathcal{S} (which only depends on the underlying surface $\Sigma_{\mathcal{S}}$). As shown in [FuSY], the specific element $\text{Cor}_{\mathcal{C}}(\mathcal{S})$ in this space of conformal blocks is invariant under the action of a subgroup $\text{Map}(\mathcal{S})$ of the mapping class group $\text{Map}(\Sigma_{\mathcal{S}})$ that is determined by the world sheet \mathcal{S} , and the assignment $\mathcal{S} \mapsto \text{Cor}_{\mathcal{C}}(\mathcal{S})$ is compatible with sewing. Accordingly, $\text{Cor}_{\mathcal{C}}(\mathcal{S})$ is naturally interpreted as the *correlator* for the world sheet \mathcal{S} .

An evident question to ask is whether the correlator $\text{Cor}_{\mathcal{C}}(\mathcal{S})$ only depends on the equivalence class $[\tilde{\mathcal{S}}]$ in the $\mathcal{F}r(\mathcal{C})$ -colored string-net space $\text{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}(\Sigma_{\mathcal{S}}, \mathbf{b}_{\mathcal{S}})$, i.e. whether two world sheets with the same underlying surface that are related by the local graphical calculus for the pivotal bicategory $\mathcal{F}r(\mathcal{C})$ have the same correlator. In [FuSY] an affirmative answer to this question was given without proof. Making use of the tools developed in the previous section, below we give a complete proof of our claim.

4.1. THE STRING-NET CONSTRUCTION OF RCFT CORRELATORS. Let us briefly review the string-net construction of RCFT correlators. For details we refer to [FuSY] and [Ya].

In the categorical approach to conformal field theory, one first develops the representation theory of the pertinent chiral vertex operator algebra \mathfrak{V} . If the chiral symmetries are sufficiently nice, the representation category of \mathfrak{V} has the structure of a *modular fusion category*; conformal field theories with such chiral symmetries are called *rational CFTs*, or RCFTs, for short. The modular fusion category \mathcal{C} of chiral data controls the monodromy behavior of the conformal blocks and gives rise to an open-closed modular functor

$$\text{Bl}_{\mathcal{C}} : \mathcal{B}\text{ord}_{2,o/c}^{\text{or}} \longrightarrow \mathcal{P}\text{rof}_{\mathbf{k}}. \quad (4.1)$$

Conjecturally [BakK], this modular functor is equivalent to the modular functor furnished by the Reshetikhin-Turaev surgery topological field theory for the Drinfeld center $\mathcal{Z}(\mathcal{C})$ or, equivalently, [TV, Bal], by the Turaev-Viro state-sum TFT for \mathcal{C} . The following result [Ki] allows us to adopt the string-net modular functor $\text{SN}_{\mathcal{C}}: \mathcal{B}\text{ord}_{2,o/c}^{\text{or}} \rightarrow \mathcal{P}\text{rof}_{\mathbf{k}}$ as an alternative realization of the modular functor $\text{Bl}_{\mathcal{C}}$:

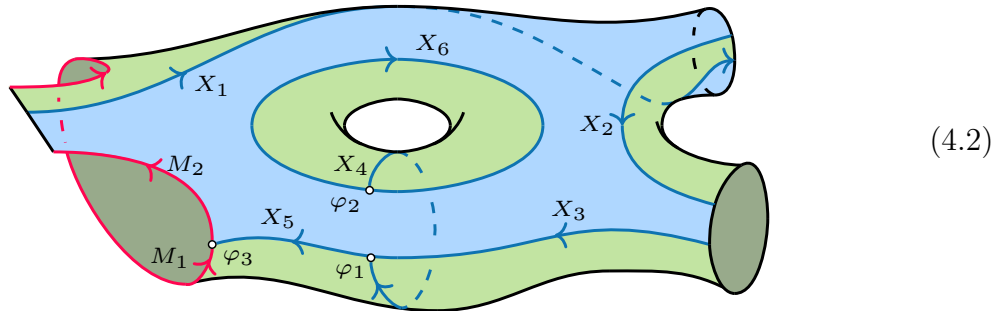
4.2. THEOREM. *Let \mathcal{C} be a spherical fusion category. The Karoubified string-net modular functor $\text{SN}_{\mathcal{C}}$ is equivalent to the Turaev-Viro modular functor $\text{TV}_{\mathcal{C}}$.*

4.3. **REMARK.** In fact, the modular functor $\text{SN}_{\mathcal{C}}$ canonically extends to a 3-2-1 topological field theory (in the sense of [BDSV2]) that is isomorphic to the once-extended Turaev-Viro TFT for \mathcal{C} [Go]. For these topological field theories to exist it is crucial that the modular tensor category \mathcal{C} is a fusion category, i.e. semisimple. The absence of any semisimplicity condition in the constructions in the previous sections may be taken as an indication that, in contrast, Theorem 4.2 admits a generalization to non-semisimple modular tensor categories.

A world sheet \mathcal{S} is a stratified surface. The 1-cells in the interior of \mathcal{S} are topological *defect lines*, while the 1-cells on the geometric boundary $\partial\mathcal{S}$ are either *sewing boundaries*, along which world sheets can be sewn, or *physical boundaries*. Further, \mathcal{S} comes with a decoration by the pointed strictly pivotal bicategory $\mathcal{Fr}(\mathcal{C})$ of simple special symmetric Frobenius algebras (as described in Examples 2.5 and 3.20) that is associated to the modular fusion category \mathcal{C} : each 2-cell is colored by a *phase* of the RCFT, which is an object of $\mathcal{Fr}(\mathcal{C})$; defect lines and physical boundaries are colored by 1-morphisms in $\mathcal{Fr}(\mathcal{C})$, and junctions of defect lines as well as junctions of defect lines and physical boundaries by 2-morphisms. We illustrate this decoration by the following example:

4.4. **EXAMPLE.**

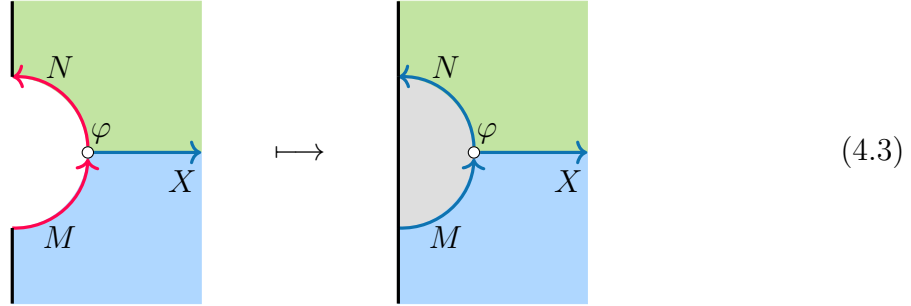
The world sheet



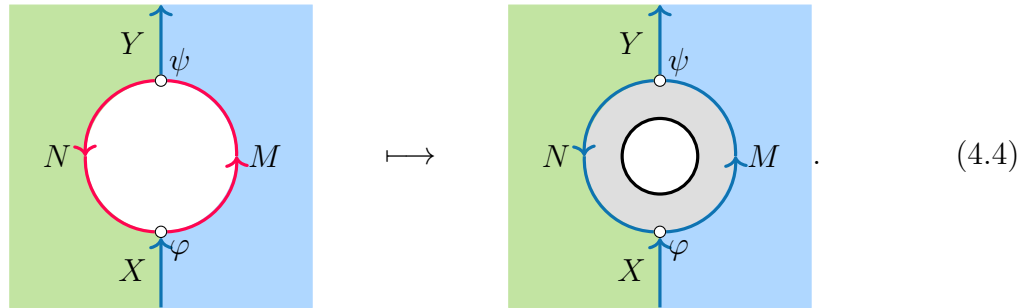
is decorated as follows: Its two 2-cells are colored with phases $A, B \in \mathcal{Fr}(\mathcal{C})$, which we indicate by different shadings of the 2-cells (green and blue, respectively, in the color version); its six line defects are colored by *defect conditions* $X_1, X_2, X_3, X_5 \in A\text{-mod}^{\mathcal{C}}\text{-}B$, $X_4 \in A\text{-mod}^{\mathcal{C}}\text{-}A$ and $X_6 \in B\text{-mod}^{\mathcal{C}}\text{-}A$; its two physical boundary segments are colored with *boundary conditions* $M_1 \in \text{mod}^{\mathcal{C}}\text{-}A = \mathbb{1}\text{-mod}^{\mathcal{C}}\text{-}A$ and $M_2 \in \text{mod}^{\mathcal{C}}\text{-}B = \mathbb{1}\text{-mod}^{\mathcal{C}}\text{-}B$; finally, its three point defects are colored by φ_1, φ_2 and φ_3 , which (upon making the auxiliary choice of a polarization for each point defect viewed as an internal vertex) are bimodule morphisms of appropriate type.

Owing to the presence of physical boundary segments, world sheets like the one shown in the picture (4.2) are not quite $\mathcal{Fr}(\mathcal{C})$ -colored graphs on their underlying surface in the sense defined at the beginning of Section 3.1: These segments are colored edges contained in the boundary of the surface, and an $\mathcal{Fr}(\mathcal{C})$ -colored graph cannot have such edges. To remedy this problem, we replace the world sheet \mathcal{S} by another stratified surface $\tilde{\mathcal{S}}$ that does constitute an $\mathcal{Fr}(\mathcal{C})$ -colored graph. In performing this replacement we are guided

by the requirement to keep the correct space of conformal blocks, which means that the topology of the world sheet should not be altered, together with the observation that a physical boundary segment adjacent to a 2-cell with phase label A can equivalently be viewed as a line defect between the phase A and the *trivial phase* – the phase labeled with the trivial Frobenius algebra $\mathbb{1} \in \mathcal{F}r(\mathcal{C})$. This leads us to the following procedure for turning a world sheet \mathcal{S} into an $\mathcal{F}r(\mathcal{C})$ -colored string diagram on a surface $\tilde{\mathcal{S}}$ that is homeomorphic to \mathcal{S} (for more details see [FuSY, Sect.3.4] and [Ya, Sect.2.4]): If a geometric boundary circle c of \mathcal{S} contains both physical boundary segments and sewing boundaries, we attach to each connected component of the union of physical boundaries in c a 2-cell that is homeomorphic to a disk and is *transparent* in the sense that it is colored by $\mathbb{1}$, as illustrated in the following picture:



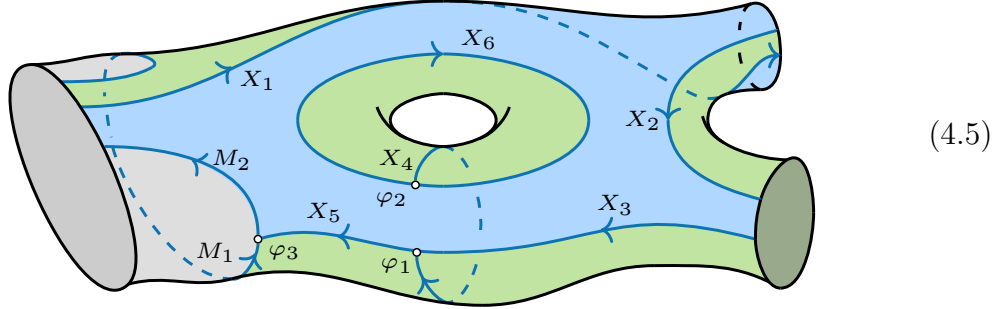
If, on the other hand, c is a *pure physical boundary*, i.e. does not contain any sewing boundaries, then we attach to it a transparent 2-cell that is homeomorphic to a cylinder, for instance



For any world sheet \mathcal{S} this prescription gives us a new stratified surface $\tilde{\mathcal{S}}$, to be referred to as the *complemented world sheet* of \mathcal{S} . We denote the underlying surface of $\tilde{\mathcal{S}}$, obtained by forgetting all strata along with their labels, by $\Sigma_{\mathcal{S}}$ and call it the *ambient surface* of \mathcal{S} . Note that there is a canonical embedding $\mathcal{S} \hookrightarrow \Sigma_{\mathcal{S}}$. Strictly speaking, neither the complemented world sheet $\tilde{\mathcal{S}}$ nor the ambient surface $\Sigma_{\mathcal{S}}$ are uniquely determined by our prescription. However, given any two such constructions, there exists a unique homeomorphism (up to isotopies) between the ambient surfaces that is compatible with the embeddings.

4.5. EXAMPLE.

For the world sheet of Example 4.4 the complemented world sheet looks as follows:



The complemented world sheet $\tilde{\mathcal{S}}$ of a world sheet \mathcal{S} can be regarded as an $\mathcal{F}r(\mathcal{C})$ -colored graph on the ambient surface $\Sigma_{\mathcal{S}}$. Denote by $\mathbf{b}_{\mathcal{S}} \in \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \partial\Sigma)$ the boundary datum of $\tilde{\mathcal{S}}$. The *field functor* [Ya, Sect. 5.2] $\mathbb{F}_{\partial\Sigma_{\mathcal{S}}}: \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, \partial\Sigma_{\mathcal{S}}) \rightarrow \text{Cyl}(\mathcal{C}, \partial\Sigma_{\mathcal{S}})$ for the 1-manifold $\partial\Sigma$ sends $\mathbf{b}_{\mathcal{S}}$ to an object $\mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}})$ in the Karoubified cylinder category $\text{Cyl}(\mathcal{C}, \partial\Sigma_{\mathcal{S}})$. We define the *space of conformal blocks for the world sheet \mathcal{S}* to be the Karoubified \mathcal{C} -colored string-net space for the pair $(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}}))$:

$$\text{Bl}_{\mathcal{C}}(\mathcal{S}) := \text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}})). \quad (4.6)$$

Now recall the rigid separable Frobenius functor $\mathcal{U}: \mathcal{F}r(\mathcal{C}) \rightarrow \mathcal{BC}$ from Example 2.34. The correlator

$$\text{Cor}_{\mathcal{C}}(\mathcal{S}) \in \text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma_{\mathcal{S}}}(\mathbf{b}_{\mathcal{S}})) \quad (4.7)$$

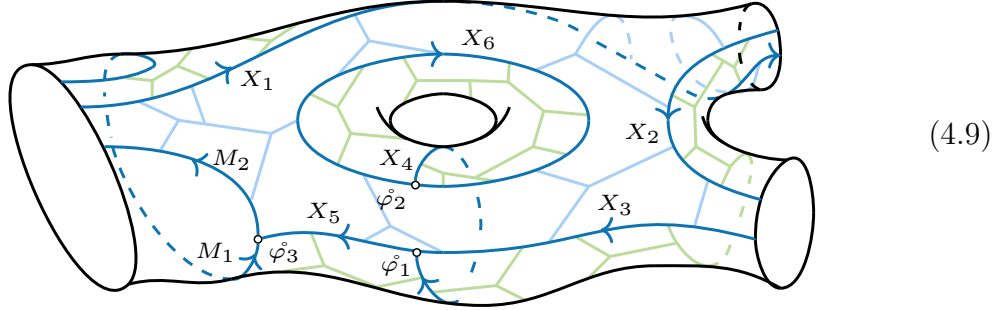
for the world sheet \mathcal{S} is defined as the \mathcal{C} -colored string-net that is represented by a \mathcal{C} -colored graph on $\Sigma_{\mathcal{S}}$ which is obtained by the following two-step procedure: First we perform \mathcal{U} -conjugation, as given in Definition 2.26, on $\tilde{\mathcal{S}}$, that is, relabel the edges of $\tilde{\mathcal{S}}$ with the underlying \mathcal{C} -objects of the bimodules and relabel the internal vertices with the \mathcal{U} -conjugates of the bimodule morphisms; the 2-cells are relabeled with the unique object $*$ of \mathcal{BC} of the delooping of \mathcal{C} . This first step results in a \mathcal{C} -colored graph $\Gamma_{\mathcal{S}}$ on $\Sigma_{\mathcal{S}}$, which we call the *partial defect network for \mathcal{S}* . In the second step we add a *full Frobenius graph* [FuSY, Def. 3.20] Γ_{ϑ} to each 2-cell of $\tilde{\mathcal{S}}$. The correlator is then the string-net equivalence class of the so obtained graph [FuSY, Def. 3.26]:

$$\text{Cor}_{\mathcal{C}}(\mathcal{S}) := [\Gamma_{\mathcal{S}} \cup \bigcup_{\vartheta \in \tilde{\mathcal{S}}} \Gamma_{\vartheta}] \in \text{SN}_{\mathcal{C}}(\Sigma_{\mathcal{S}}, \mathbb{F}_{\partial\Sigma}(\mathbf{b}_{\mathcal{S}})). \quad (4.8)$$

The properties of special Frobenius algebras and their bimodules ensure that the string-net correlator $\text{Cor}_{\mathcal{C}}(\mathcal{S})$ is well defined and invariant under the action of $\text{Map}(\mathcal{S})$. As shown in Theorem 3.28 of [FuSY], the assignment $\mathcal{S} \mapsto \text{Cor}_{\mathcal{C}}(\mathcal{S})$ gives a consistent system of correlators.

4.6. EXAMPLE.

As an illustration of the prescription, consider again the world sheet (4.2)). A representative \mathcal{C} -colored graph for the string net that gives the correlator $\text{Cor}_{\mathcal{C}}(\mathcal{S})$ is shown in the following picture:



Here the unlabeled trivalent graphs (drawn in light green and light blue in the color version) are implicitly colored with the corresponding Frobenius algebra A and B and by their structure morphisms (using simplified graphical calculus, as explained in Appendix A.8 of [FuSY]).

4.7. UNIVERSAL CORRELATORS. Let Σ be a compact oriented surface and $\mathbf{b} \in \text{Cyl}^\circ(\mathcal{F}r(\mathcal{C}), \partial\Sigma)$ an $\mathcal{F}r(\mathcal{C})$ -boundary datum over $\partial\Sigma$. Denote by $\mathbb{k}\mathbf{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b})$ the vector space freely generated by the set of $\mathcal{F}r(\mathcal{C})$ -colored graphs on Σ with boundary datum \mathbf{b} (every such graph can be viewed as a complemented world sheet). The assignment $\tilde{\mathcal{S}} \mapsto \text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b})$ defines a linear map

$$\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b}) : \mathbb{k}\mathbf{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b}) \longrightarrow \text{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b})). \quad (4.10)$$

Evaluating the map $\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b})$ on any vector of the distinguished basis $\mathbf{G}_{\mathcal{F}r(\mathcal{C})}$ of $\mathbb{k}\mathbf{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b})$ yields a correlator. In this sense, $\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b})$ collects correlators for different world sheet structures on the surface Σ with boundary datum \mathbf{b} .

The following result, which was conjectured in [FuSY, Sect. 6.2], shows that in fact some of these correlators coincide.

4.8. THEOREM. *For every compact oriented surface Σ and every $\mathcal{F}r(\mathcal{C})$ -boundary datum $\mathbf{b} \in \text{Cyl}^\circ(\mathcal{F}r(\mathcal{C}), \partial\Sigma)$ there exists a unique $\text{Map}(\Sigma)$ -intertwiner*

$$\text{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b}) : \text{SN}_{\mathcal{F}r(\mathcal{C})}^\circ(\Sigma, \mathbf{b}) \longrightarrow \text{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b})) \quad (4.11)$$

such that the diagram

$$\begin{array}{ccc} \mathbb{k}\mathbf{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b}) & \xrightarrow{\text{Cor}_{\mathcal{C}}(\Sigma, \mathbf{b})} & \text{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b})) \\ \downarrow & \nearrow \text{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b}) & \\ \text{SN}_{\mathcal{F}r(\mathcal{C})}^\circ(\Sigma, \mathbf{b}) & & \end{array} \quad (4.12)$$

commutes, where $\mathbb{k}\mathbf{G}_{\mathcal{F}r(\mathcal{C})}(\Sigma, \mathbf{b}) \twoheadrightarrow \mathrm{SN}_{\mathcal{F}r(\mathcal{C})}^\circ(\Sigma, \mathbf{b})$ is the canonical quotient map.

We call this unique map $\mathrm{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b})$ the *universal correlator* for the pair (Σ, \mathbf{b}) .

PROOF. (i) Every representative of a string net in $\mathrm{SN}_{\mathcal{F}r(\mathcal{C})}^\circ(\Sigma, \mathbf{b})$ can be regarded as a complemented world sheet. Therefore commutativity of the triangle (4.12) forces the universal correlator to be given by

$$\begin{aligned} \mathrm{UCor}_{\mathcal{C}}(\Sigma, \mathbf{b}) : \quad \mathrm{SN}_{\mathcal{F}r(\mathcal{C})}^\circ(\Sigma, \mathbf{b}) &\longrightarrow \mathrm{SN}_{\mathcal{C}}(\Sigma, \mathbb{F}_{\partial\Sigma}(\mathbf{b})), \\ [\tilde{\mathcal{S}}] &\longmapsto \mathrm{Cor}_{\mathcal{C}}(\mathcal{S}). \end{aligned} \quad (4.13)$$

Thus we first need to show that this is indeed well defined as a linear map, i.e. that any two world sheets related by the graphical calculus on disks for the pivotal bicategory $\mathcal{F}r(\mathcal{C})$ have the same correlator. To see this we assume, without loss of generality, that $\tilde{\mathcal{S}}_1$ and $\tilde{\mathcal{S}}_2$ are two complemented world sheets which have the same ambient surface Σ , are identical outside an embedded disk $D \hookrightarrow \Sigma$, and yield the same value, in the sense of (2.45), on the disk, i.e. $\langle \tilde{\mathcal{S}}_1 \cap D \rangle_{\mathcal{F}r(\mathcal{C})} = \langle \tilde{\mathcal{S}}_2 \cap D \rangle_{\mathcal{F}r(\mathcal{C})}$. Let $\mathrm{Cor}_{\mathcal{C}}(\mathcal{S}_1) = [\Gamma_1]$ and $\mathrm{Cor}_{\mathcal{C}}(\mathcal{S}_2) = [\Gamma_2]$ be the correlators for the corresponding world sheets, with the representative graphs Γ_1 and Γ_2 chosen in a way such that they coincide outside D , and such that for each of them there are no other Frobenius lines within D apart from, for each pair of distinct connected components (within D) of the partial defect network, a single Frobenius line connecting the two components. Such a choice is possible because every \mathcal{U} -conjugation of bimodule morphisms commutes with the action of the relevant Frobenius algebras and because all the Frobenius graphs involved are *full*. It then remains to be shown that the equality $\langle \Gamma_1 \cap D \rangle_{\mathcal{C}} = \langle \Gamma_2 \cap D \rangle_{\mathcal{C}}$ holds. This is indeed the case, because the functor $\mathcal{U} : \mathcal{F}r(\mathcal{C}) \rightarrow \mathcal{BC}$ introduced in Example 2.34 (with the canonical lax and oplax structures) is rigid separable Frobenius, so that \mathcal{U} -conjugation preserves *operadic compositions* and *partial trace maps* (see Theorem 2.30), while the presence of the Frobenius lines which connect the connected components (on the embedded disk D) of the partial defect networks compensates for the fact that \mathcal{U} -conjugation preserves *horizontal products* and *whiskerings* only up to idempotents of the type (2.93). It is also readily clear that the linear map (4.13) intertwines the mapping class group actions.

(ii) That the collection of universal correlators is compatible with sewing translates exactly to the statement that the prescription of string-net correlators is compatible with sewing. \blacksquare

4.9. REMARK. In [FuSY] we called the string net $[\tilde{\mathcal{S}}] \in \mathrm{SN}_{\mathcal{B}}^\circ(\Sigma_{\mathcal{S}}, \mathbf{b}_{\mathcal{S}})$ a *quantum world sheet*. That the correlator $\mathrm{Cor}_{\mathcal{C}}(\mathcal{S})$ depends only on the quantum world sheet $[\tilde{\mathcal{S}}]$ implies the validity of the *calculus of defects* which is a crucial feature of conformal field theory: one is allowed to modify a world sheet \mathcal{S} locally according to the graphical calculus for the pivotal bicategory $\mathcal{F}r(\mathcal{C})$ of defects without changing the value of its correlator $\mathrm{Cor}_{\mathcal{C}}(\mathcal{S})$. This feature is e.g. implicitly assumed in [FrFRS], where a less conceptual justification is given. In [FreMT] it is demonstrated that the calculus of defects, in a higher-dimensional

setting, is a powerful tool for implementing categorical symmetries and should therefore be postulated for *every* reasonable quantum field theory with topological defects.

Consider now a rigid separable Frobenius functor $F: \mathcal{B} \rightarrow \mathcal{B}'$ between two arbitrary pivotal bicategories. Recall from Remark 2.29 that the lax and oplax structures of F canonically provide for every object a in the domain bicategory \mathcal{B} a Δ -separable symmetric Frobenius algebra $F(\text{id}_a)$ in the pivotal tensor category $\mathcal{B}'(Fa, Fa)$, for every 1-morphism $f: a \rightarrow b$ in \mathcal{B} an $F(\text{id}_a)$ - $F(\text{id}_b)$ -bimodule $F(f) \in \mathcal{B}'(Fa, Fb)$, and for every 2-morphism α in \mathcal{B} a bimodule morphism given by the F -conjugate α^F of α . (As an illustration, for 1-morphisms $f: a \rightarrow b$, $g: b \rightarrow c$, $h: a \rightarrow b'$ and $k: b' \rightarrow c$ in \mathcal{B} , the F -conjugate of a 2-morphism $\alpha \in \text{Hom}_{\mathcal{B}(a,c)}(f \star g, h \star k)$ is a bimodule morphism $\alpha^F \in \text{Hom}_{\mathcal{B}'(Fa,Fc)}(Ff \otimes_{F\text{id}_b} Fg, Fh \otimes_{F\text{id}_{b'}} Fk)$.)

Just like in the case of string-net correlators, we can associate to any \mathcal{B} -colored graph on any oriented surface Σ a \mathcal{B}' -colored graph (and thus a string net) on Σ by performing a change of color prescribed by the 2-functor F accompanied by adding a full Frobenius graph in each 2-cell of the embedded graph. This is demonstrated schematically by the following figure:

$$(4.14)$$

in which the Frobenius graphs are labeled with the images of identity 1-morphisms and lax/oplax functoriality and unit constraints. Moreover, by an argument analogous to the one that yields Theorem 4.8, such a transformation of colored graphs descends to the level of string nets. We then obtain the following result which is parallel to Theorem 3.11:

4.10. THEOREM. *Let Σ be a compact oriented surface, \mathcal{B} and \mathcal{B}' two strictly pivotal bicategories, $F: \mathcal{B} \rightarrow \mathcal{B}'$ a rigid separable Frobenius functor, and \mathbf{b} a \mathcal{B} -boundary datum on $\partial\Sigma$. There is a canonical $\text{Map}(\Sigma)$ -intertwiner*

$$\text{SN}_F^\circ(\Sigma, \mathbf{b}) : \text{SN}_\mathcal{B}^\circ(\Sigma, \mathbf{b}) \longrightarrow \text{SN}_{\mathcal{B}'}(\Sigma, \mathbb{F}^F \mathbf{b}), \quad (4.15)$$

where $\mathbb{F}^F \mathbf{b}$ is an object in the Karoubified cylinder category $\text{Cyl}(\mathcal{B}', \partial\Sigma)$. The linear map (4.15) is defined by sending each representing \mathcal{B} -colored graph to the \mathcal{B}' -colored graph obtained via F -conjugation, with full Frobenius graphs added. Moreover, the collection of such intertwiners corresponding to different surfaces and boundary data is compatible with the concatenation of string nets.

4.11. **REMARK.** The ambient surface $\Sigma_{\mathcal{S}}$ of any world sheet \mathcal{S} can be equipped with the structure of an open-closed bordism; accordingly we refer to $\Sigma_{\mathcal{S}}$ as an *ambient bordism* of \mathcal{S} . The boundary parametrization of $\Sigma_{\mathcal{S}}$ is compatible with the sewing intervals and sewing circles that are contained in the boundary of \mathcal{S} (see [Ya, Sect. 2.4]). Moreover, the parametrization map for an interval I determines from the complemented world sheet $\tilde{\mathcal{S}}$ an object in the cylinder category $\text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, I)$. Conversely, the \mathbb{k} -linear profunctor

$$\mathcal{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}(\Sigma) : \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, \alpha_1) \rightarrow \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, \alpha_2) \quad (4.16)$$

that is the value of an open-closed bordism $\Sigma: \alpha_1 \rightarrow \alpha_2$ under the modular functor $\mathcal{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}$ classifies all world sheets with prescribed ambient bordism Σ . These relationships fit well with the following observation [Ya, Sect. 9], which we will further explore elsewhere: First, the modular functors $\mathcal{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}$ and $\mathcal{SN}_{\mathcal{C}}$ can be promoted to symmetric monoidal double functors

$$\mathcal{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ}, \mathcal{SN}_{\mathcal{C}} : \text{Bord}_{2, \text{o/c}}^{\text{or}} \longrightarrow \text{Prof}_{\mathbb{k}} \quad (4.17)$$

between the symmetric monoidal double categories $\text{Bord}_{2, \text{o/c}}^{\text{or}}$ of open-closed bordisms (which has orientation preserving embeddings as vertical 1-morphisms) and $\text{Prof}_{\mathbb{k}}$ of \mathbb{k} -linear profunctors (having linear functors as vertical 1-morphisms). And second, the universal correlators together with the field functors $\{\mathbb{F}_{\alpha} : \text{Cyl}^{\circ}(\mathcal{F}r(\mathcal{C}), \mathbb{1}, \alpha) \rightarrow \text{Cyl}(\mathcal{C}, \alpha)\}_{\alpha \in \text{Bord}_{2, \text{o/c}}^{\text{or}}}$ fit into a monoidal vertical transformation $\mathcal{SN}_{\mathcal{F}r(\mathcal{C})}^{\circ} \Rightarrow \mathcal{SN}_{\mathcal{C}}$ between these symmetric monoidal double functors.

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